

Online Social Networks and Media

Link Analysis and Web Search

How to Organize the Web

First try: Human curated Web directories

Yahoo, DMOZ, LookSmart

YAHOO! DIRECTORY Yahoo! | Help

Search: the Web | the Directory

Yahoo! Directory Advanced Search | Suggest a Site | Email This Page

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Become an Editor Help build the largest human-edited directory of the web



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5,114,642 sites - 96,895 editors - over 1,014,858 categories

How to organize the web

- **Second try: Web Search**

- Information Retrieval investigates:

- Find relevant docs in a small and trusted set e.g., Newspaper articles, Patents, etc. (“needle-in-a-haystack”)
 - Limitation of keywords (synonyms, polysemy, etc)

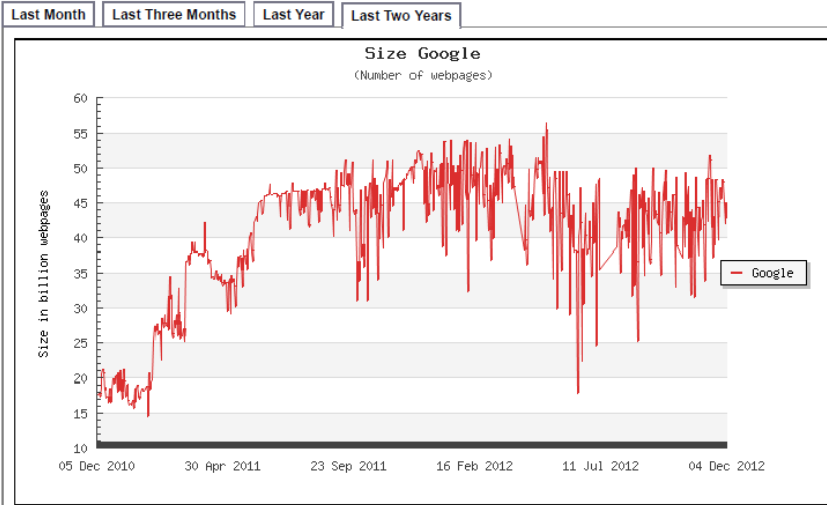
But: Web is huge, full of untrusted documents, random things, web spam, etc.

- Everyone can create a web page of high production value
- Rich diversity of people issuing queries
- Dynamic and constantly-changing nature of web content

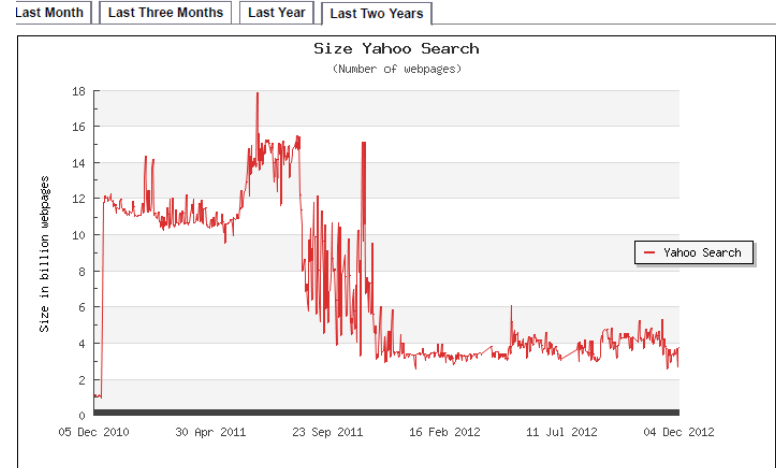
Size of the Search Index



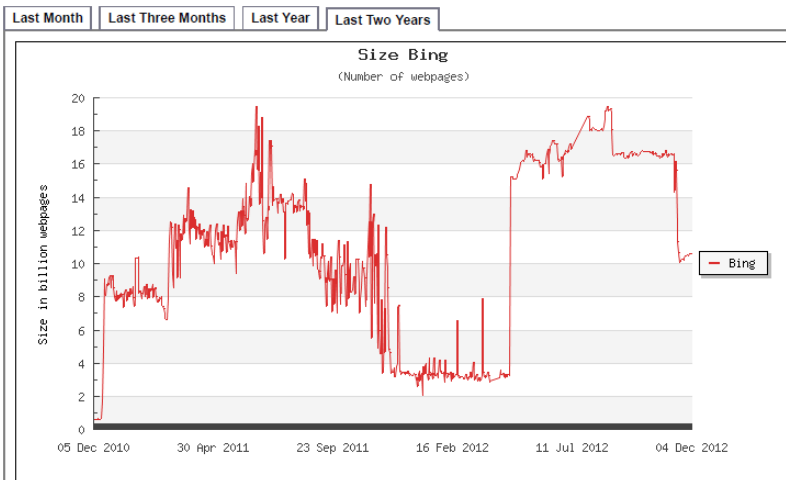
The size of the World Wide Web:
Estimated size of Google's index



The size of the World Wide Web:
Estimated size of Yahoo Search index



The size of the World Wide Web:
Estimated size of Bing index



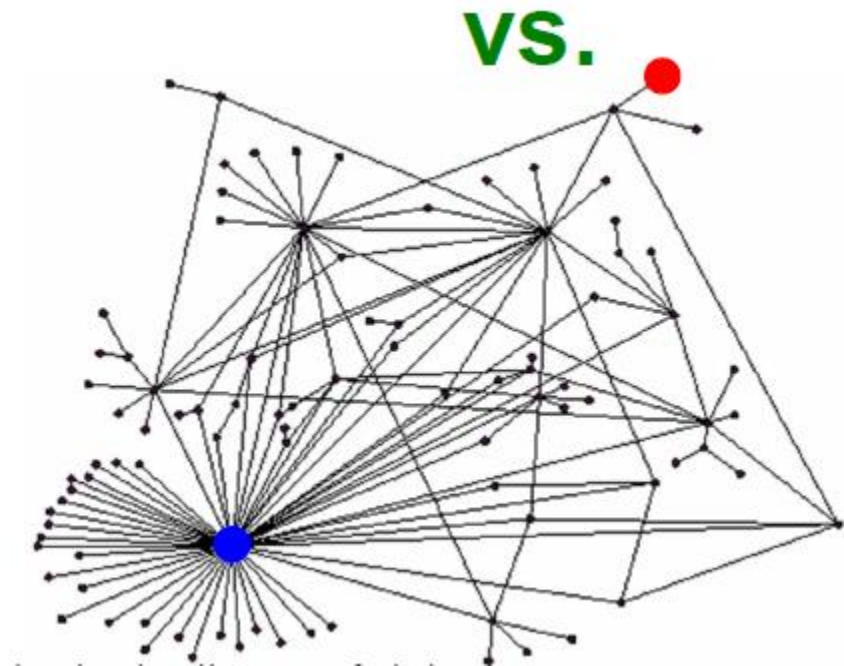
How to organize the web

- Third try (the Google era): using the web graph
 - Swift from relevance to *authoritativeness*
 - It is not only important that a page is relevant, but that it is also important on the web
- For example, what kind of results would we like to get for the query “greek newspapers”?

Link Analysis

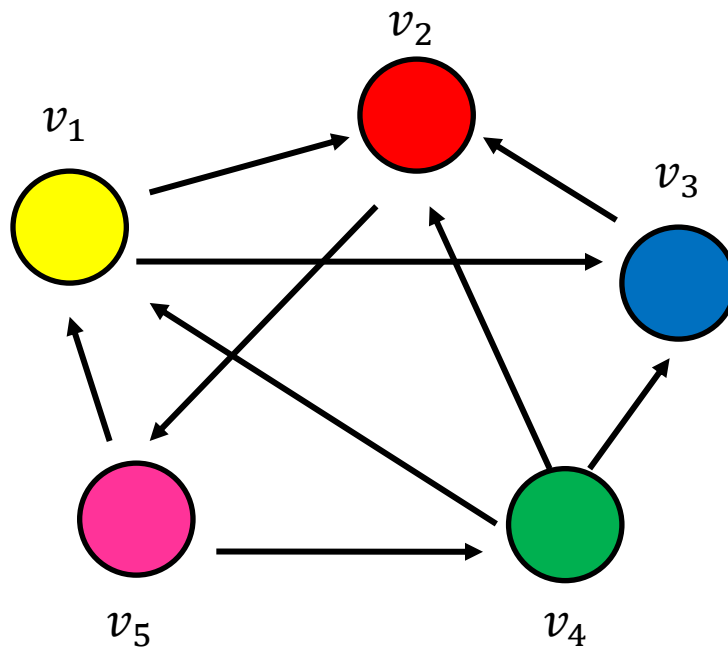
- Not all web pages are equal on the web
- The links act as **endorsements**:
 - When page p **links** to q it **endorses** the content of the content of q

What is the simplest way to measure importance of a page on the web?



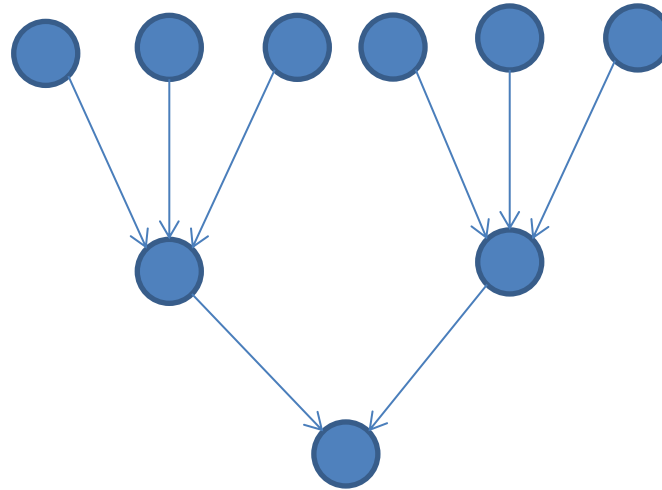
Rank by Popularity

- Rank pages according to the number of incoming edges (**in-degree**, **degree centrality**)



- 1. Red Page**
- 2. Yellow Page**
- 3. Blue Page**
- 4. Purple Page**
- 5. Green Page**

Popularity



- It is not important only **how many** link to you, but also **how important** are the people that link to you.
- **Good** authorities are pointed by **good** authorities
 - Recursive definition of importance

THE PAGERANK ALGORITHM

PageRank

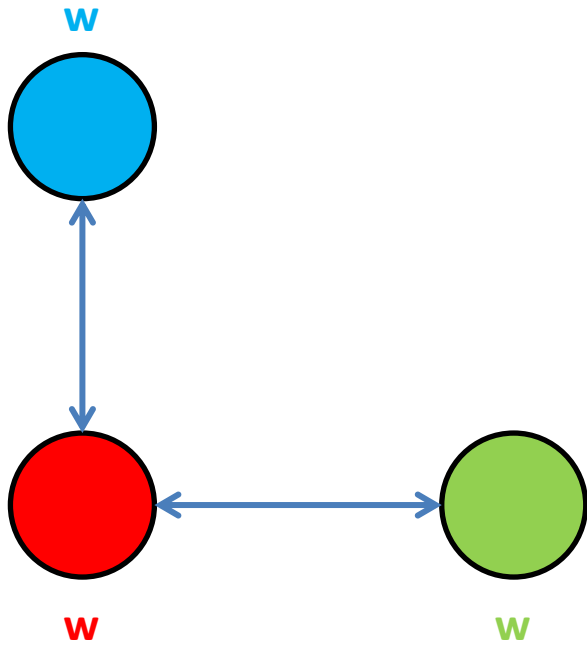
- **Good** authorities should be pointed by **good** authorities
 - The value of a node is the value of the nodes that point to it.
- How do we implement that?
 - Assume that we have **a unit of authority** to distribute to all nodes.
 - Initially each node gets $\frac{1}{n}$ amount of authority
 - Each node **distributes** the authority value they have **to their neighbors**
 - The authority value of each node is the sum of the **authority fractions** it collects from its neighbors.

$$w_v = \sum_{u \rightarrow v} \frac{1}{d_{out}(u)} w_u$$

w_v : the **PageRank value** of node v

Recursive definition

A simple example



$$w + w + w = 1$$

$$w = w + w$$

$$w = \frac{1}{2} w$$

$$w = \frac{1}{2} w$$

- Solving the system of equations we get the authority values for the nodes

$$- w = \frac{1}{2} \quad w = \frac{1}{4} \quad w = \frac{1}{4}$$

A more complex example

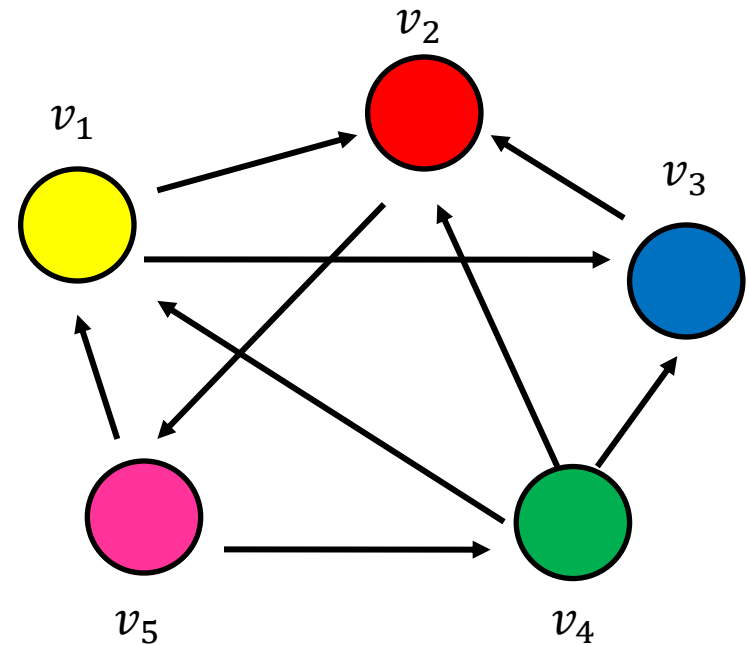
$$w_1 = \frac{1}{3} w_4 + \frac{1}{2} w_5$$

$$w_2 = \frac{1}{2} w_1 + w_3 + \frac{1}{3} w_4$$

$$w_3 = \frac{1}{2} w_1 + \frac{1}{3} w_4$$

$$w_4 = \frac{1}{2} w_5$$

$$w_5 = w_2$$



$$w_v = \sum_{u \rightarrow v} \frac{1}{d_{out}(u)} w_u$$

Computing PageRank weights

- A simple way to compute the weights is by iteratively updating the weights
- PageRank Algorithm

Initialize all PageRank weights to $\frac{1}{n}$

Repeat:

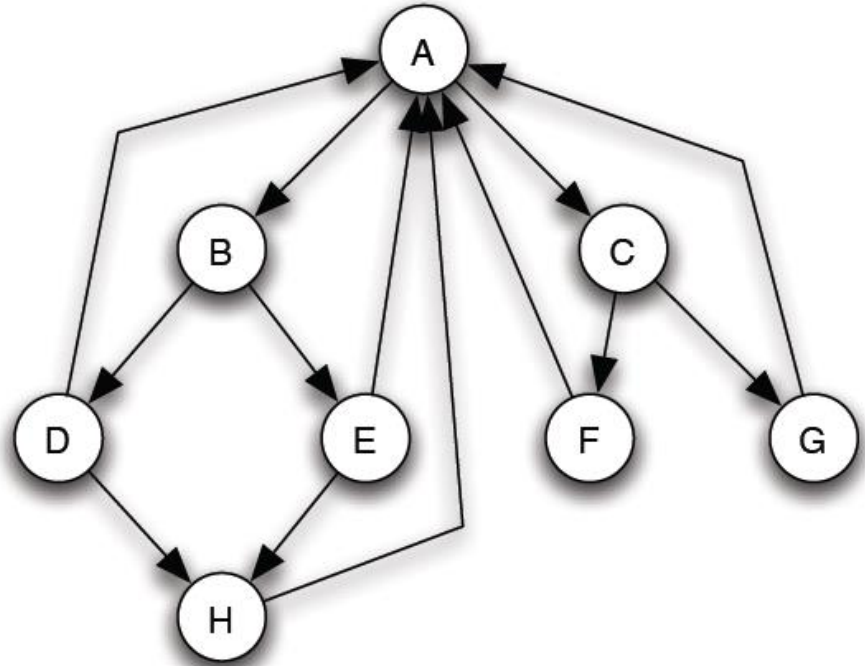
$$w_v = \sum_{u \rightarrow v} \frac{1}{d_{out}(u)} w_u$$

Until the weights do not change

- This process converges

PageRank

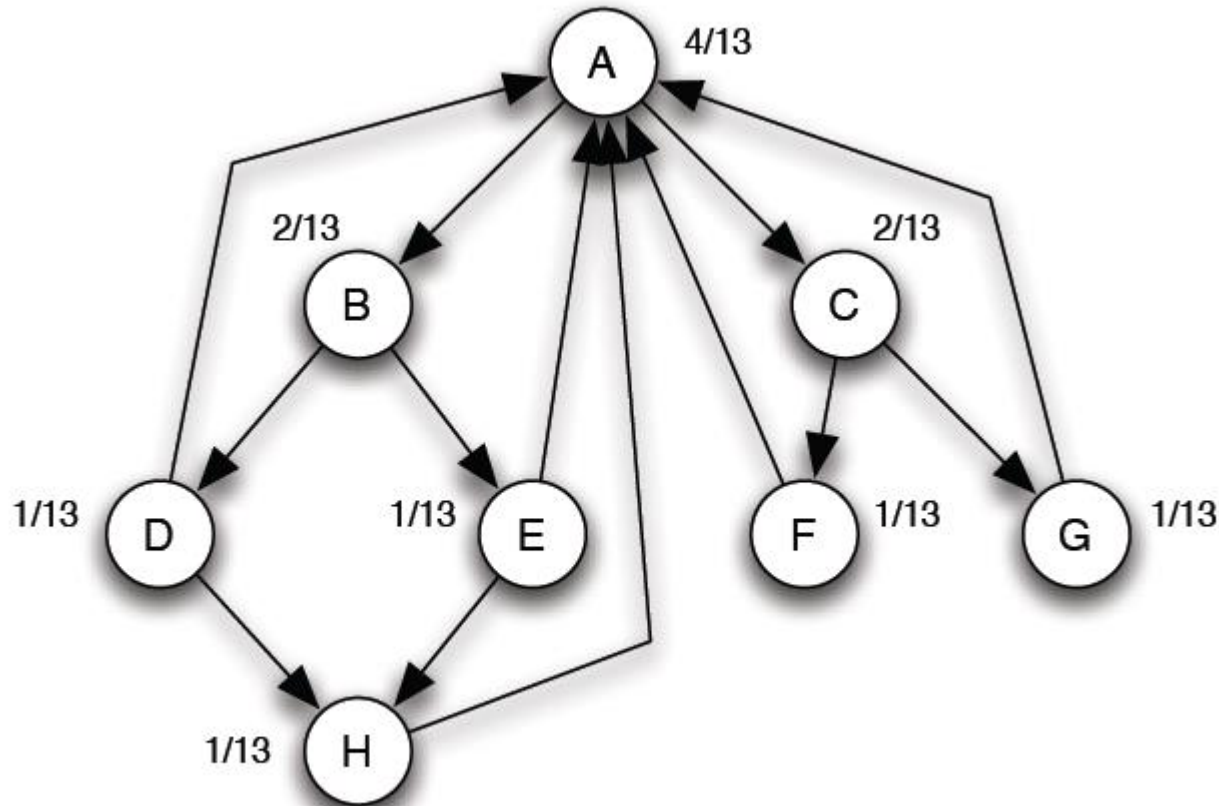
Initially, all nodes PageRank $1/8$



Step	A	B	C	D	E	F	G	H
1	$1/2$	$1/16$	$1/16$	$1/16$	$1/16$	$1/16$	$1/16$	$1/8$
2	$3/16$	$1/4$	$1/4$	$1/32$	$1/32$	$1/32$	$1/32$	$1/16$

- ✓ As a kind of “fluid” that circulates through the network
- ✓ The total PageRank in the network remains constant (no need to normalize)

PageRank: equilibrium



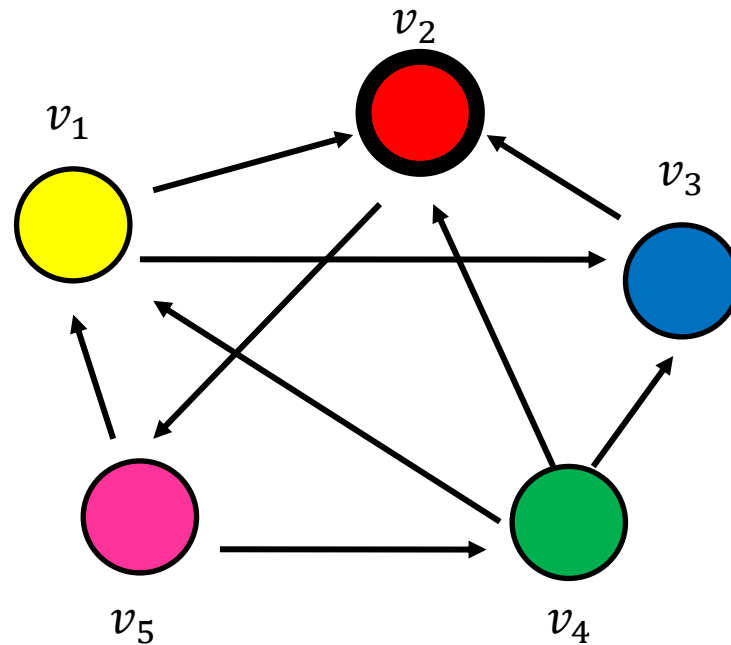
- A simple way to *check* whether an assignment of numbers forms an equilibrium set of PageRank values: check that they sum to 1, and that when apply the Basic PageRank Update Rule, we get the same values back.
- If the network is *strongly connected*, then there is a unique set of equilibrium values.

Random Walks on Graphs

- The algorithm defines a **random walk** on the graph
- Random walk:
 - **Start** from a node chosen **uniformly at random** with probability $\frac{1}{n}$.
 - **Pick** one of the **outgoing edges** **uniformly at random**
 - **Move** to the destination of the edge
 - Repeat.
- The PageRank of node v is the probability that the random walk is at node v after a very large number of steps.
- The **Random Surfer** model
 - Users wander on the web, following links.

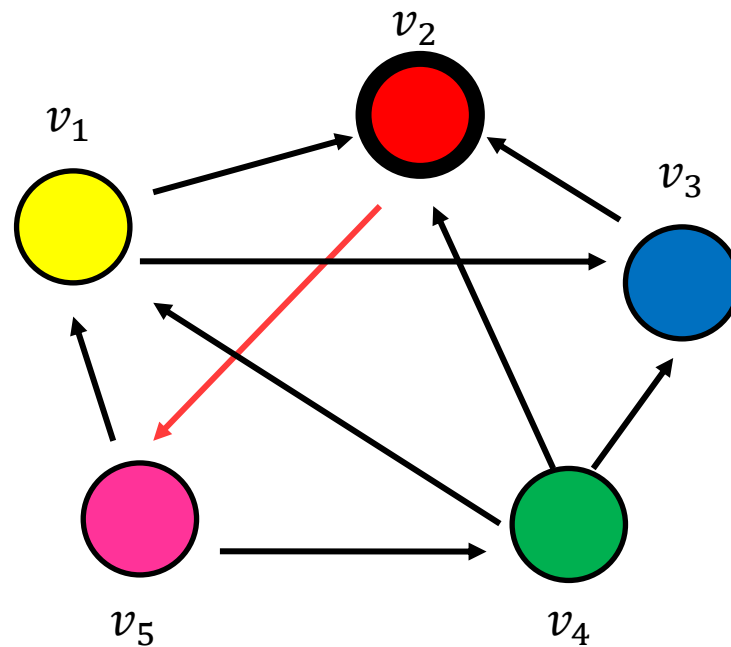
Example

- Step 0



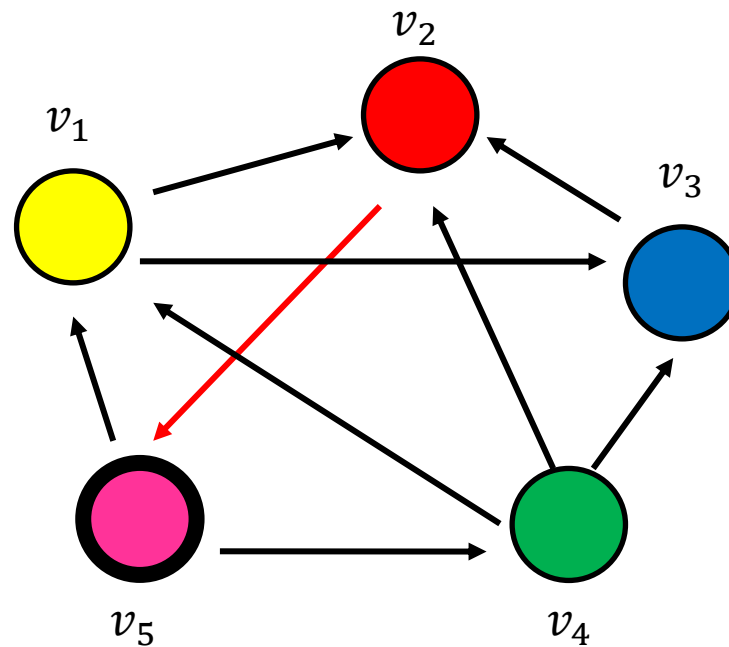
Example

- Step 0



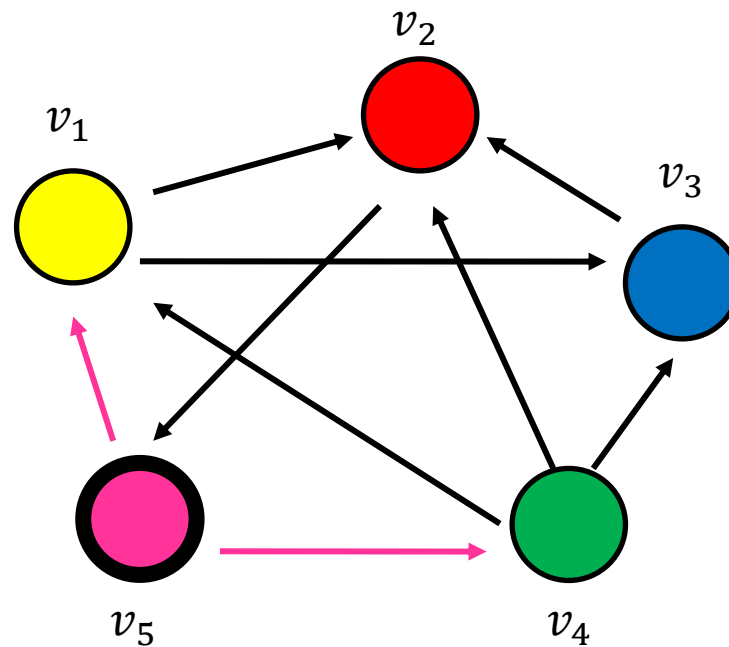
Example

- Step 1



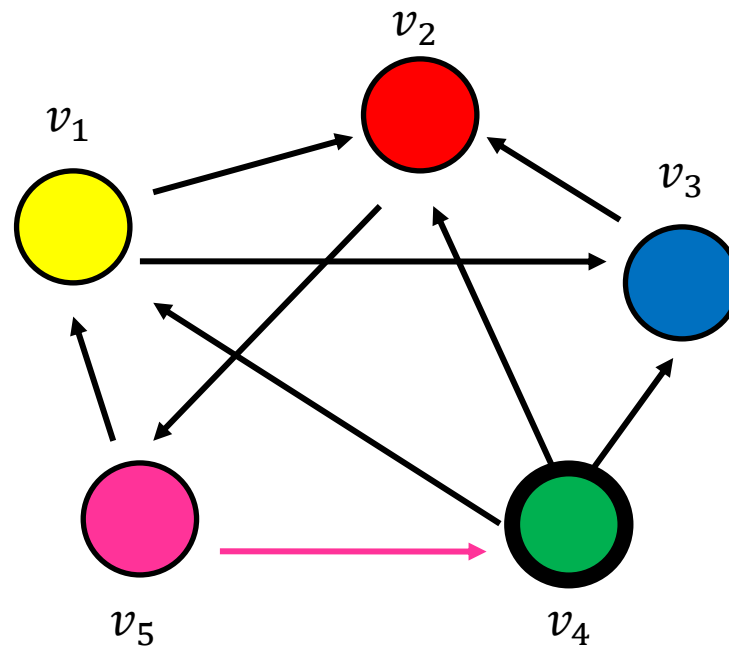
Example

- Step 1



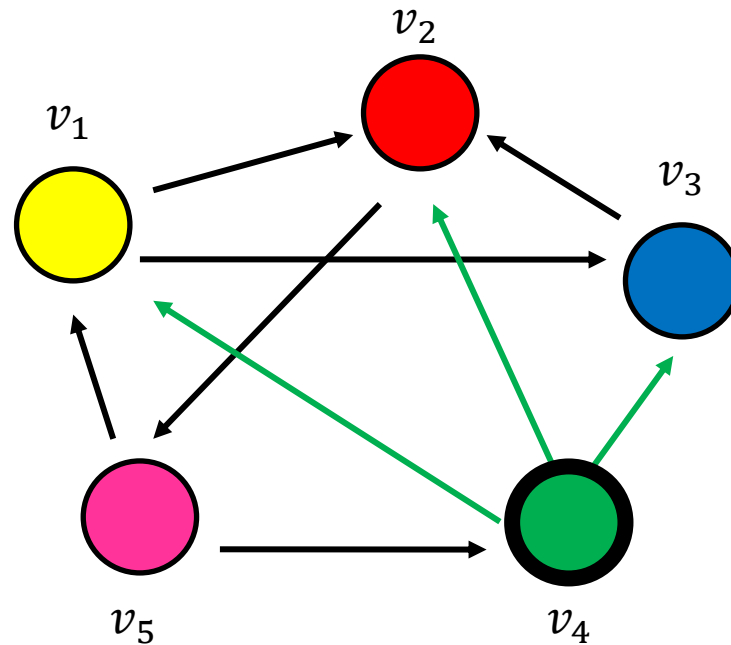
Example

- Step 2



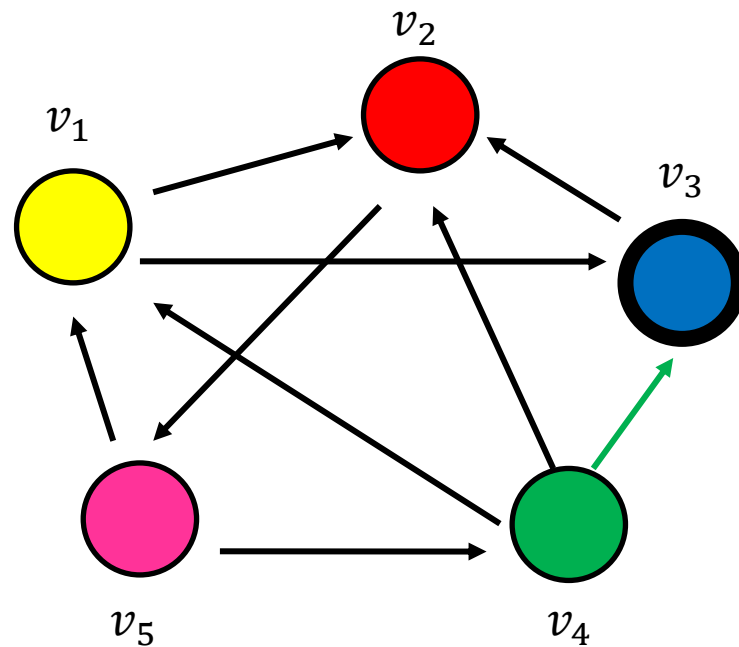
Example

- Step 2



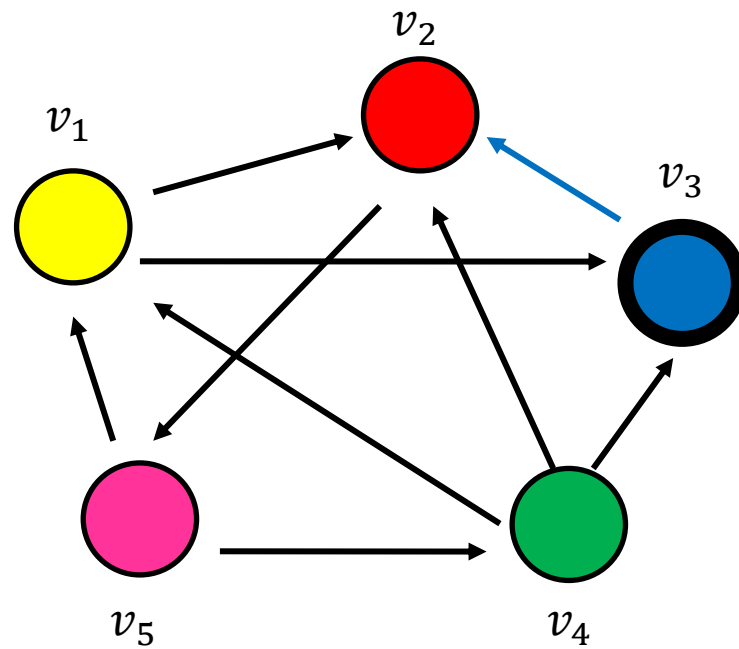
Example

- Step 3



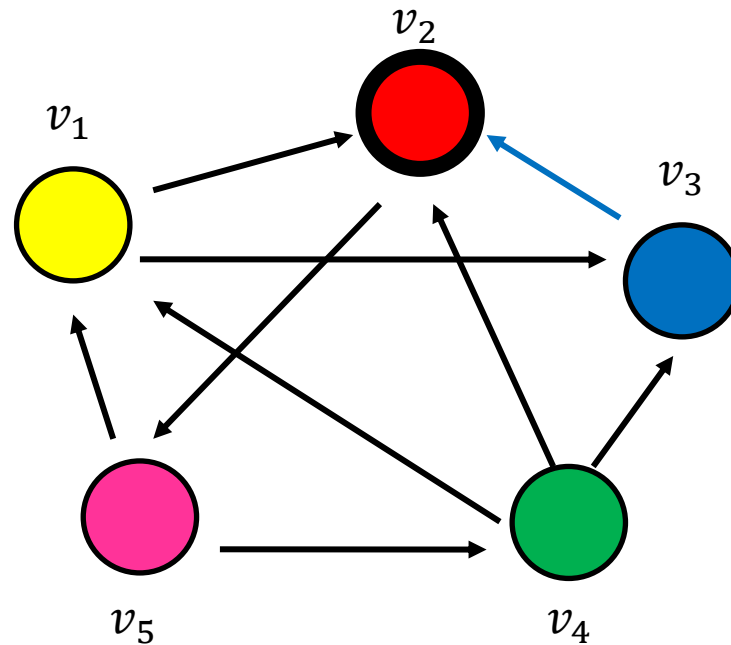
Example

- Step 3



Example

- Step 4...



Random walk

- Question: what is the probability p_i^t of being at node i after t steps?

$$p_1^0 = \frac{1}{5}$$

$$p_2^0 = \frac{1}{5}$$

$$p_3^0 = \frac{1}{5}$$

$$p_4^0 = \frac{1}{5}$$

$$p_5^0 = \frac{1}{5}$$

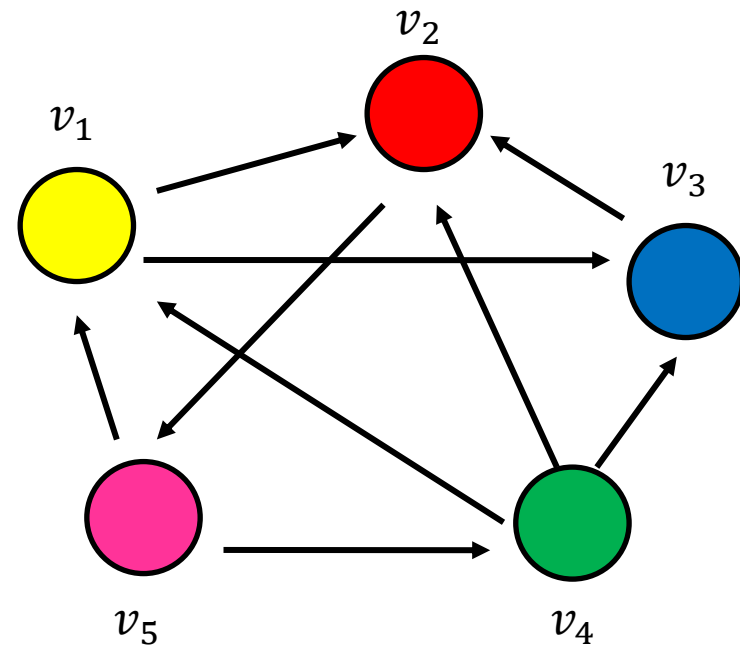
$$p_1^t = \frac{1}{3}p_4^{t-1} + \frac{1}{2}p_5^{t-1}$$

$$p_2^t = \frac{1}{2}p_1^{t-1} + p_3^{t-1} + \frac{1}{3}p_4^{t-1}$$

$$p_3^t = \frac{1}{2}p_1^{t-1} + \frac{1}{3}p_4^{t-1}$$

$$p_4^t = \frac{1}{2}p_5^{t-1}$$

$$p_5^t = p_2^{t-1}$$



Markov chains

- A Markov chain describes a **discrete time stochastic process** over a set of states

$$S = \{s_1, s_2, \dots, s_n\}$$

according to a transition probability matrix $P = \{P_{ij}\}$

- P_{ij} = probability of moving to state j when at state i

- Matrix P has the property that the entries of all **rows sum to 1**

$$\sum_j P[i, j] = 1$$

A matrix with this property is called **stochastic**

- **State probability distribution**: The vector $p^t = (p_1^t, p_2^t, \dots, p_n^t)$ that stores the probability of being at state s_i after t steps
- **Memorylessness property**: The **next state** of the chain **depends only at the current state** and not on the past of the process (**first order MC**)
 - **Higher order** MCs are also possible
- **Markov Chain Theory**: After infinite steps the **state probability vector converges** to a **unique** distribution if the chain is **irreducible** (possible to get from any state to any other state) and **aperiodic**
 - These are the PageRank values

Random walks

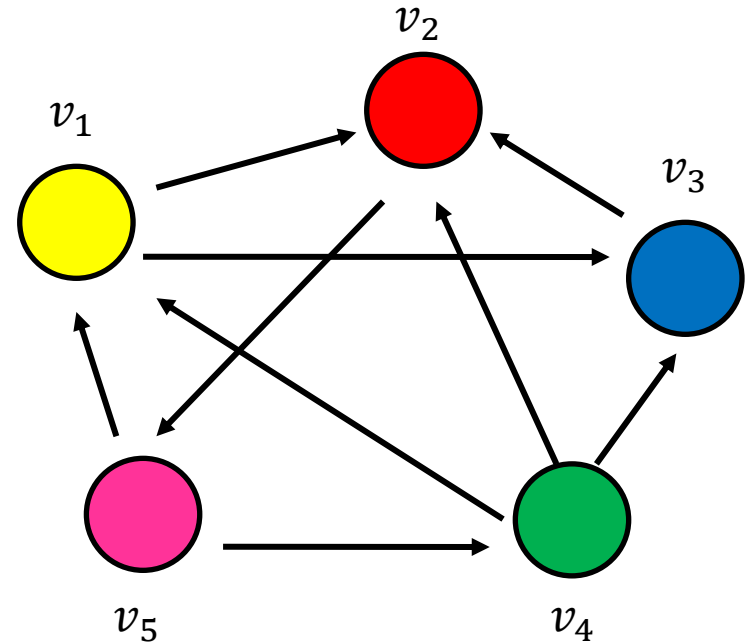
- Random walks on graphs correspond to Markov Chains
 - The set of states S is the set of nodes of the graph G
 - The **transition probability matrix** is the probability that we follow an edge from one node to another

$$P[i, j] = 1 / \text{deg}_{out}(i)$$

An example

$$A = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$



Node Probability vector

- The vector $p^t = (p_1^t, p_2^t, \dots, p_n^t)$ that stores the probability of being at node v_i at step t
- p_i^0 = the probability of starting from state i (usually) set to **uniform**
- We can compute the vector p^t at step t using a vector-matrix multiplication

$$p^t = p^{t-1} P$$

An example

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

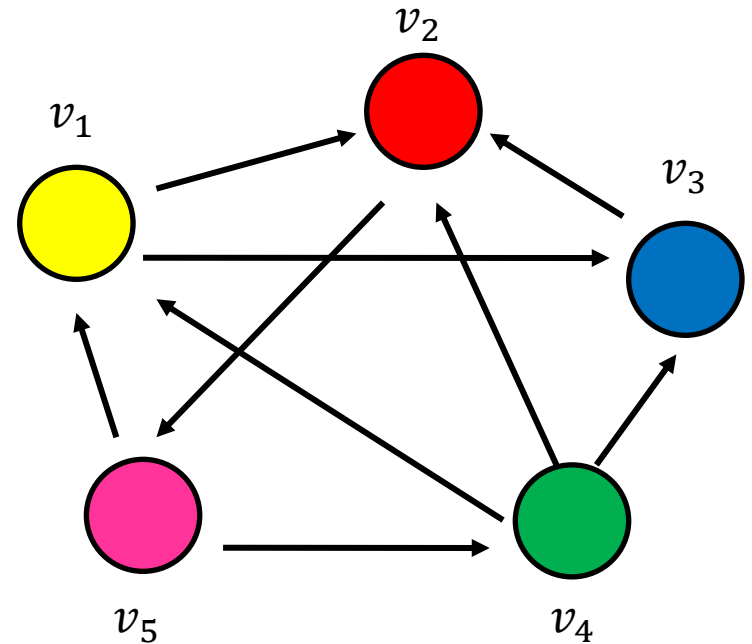
$$p_1^t = \frac{1}{3}p_4^{t-1} + \frac{1}{2}p_5^{t-1}$$

$$p_2^t = \frac{1}{2}p_1^{t-1} + p_3^{t-1} + \frac{1}{3}p_4^{t-1}$$

$$p_3^t = \frac{1}{2}p_1^{t-1} + \frac{1}{3}p_4^{t-1}$$

$$p_4^t = \frac{1}{2}p_5^{t-1}$$

$$p_5^t = p_2^{t-1}$$



Stationary distribution

- The **stationary distribution** of a random walk with transition matrix P , is a probability distribution π , such that $\pi = \pi P$
- The stationary distribution is an **eigenvector** of matrix P
 - the **principal left eigenvector** of P – stochastic matrices have maximum eigenvalue 1
- The probability π_i is the fraction of times that we visited state i as $t \rightarrow \infty$
- **Markov Chain Theory**: The random walk converges to a **unique stationary distribution independent of the initial vector** if the graph is **strongly connected**, and **not bipartite**.
 - In our case these are the PageRank values.

Computing the stationary distribution

- The Power Method

Initialize p^0 to some distribution

Repeat

$$p^t = p^{t-1}P$$

Until convergence

- After many iterations $p^t \rightarrow \pi$ regardless of the initial vector p^0
- Power method because it computes $p^t = p^0 P^t$
- Rate of convergence
 - determined by the second eigenvalue $\frac{|\lambda_2|}{|\lambda_1|}$

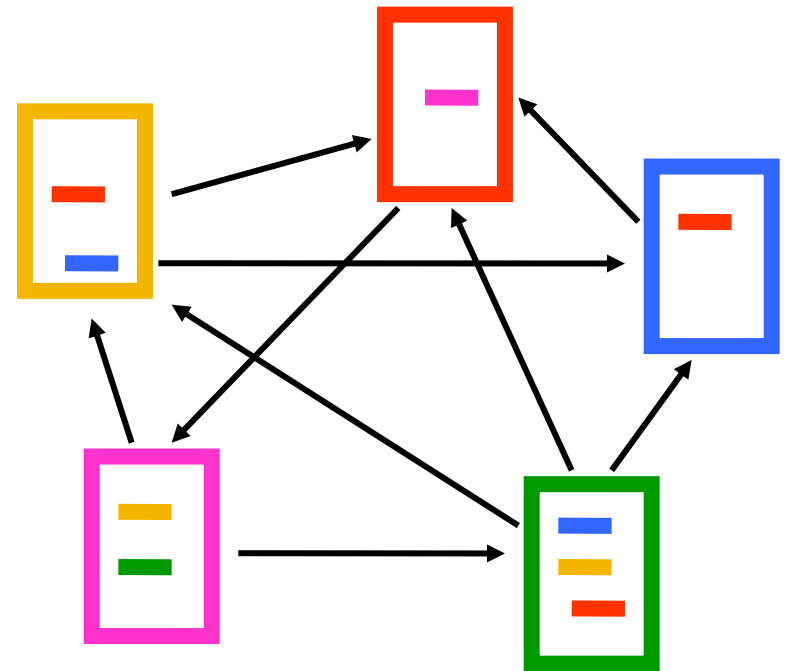
The stationary distribution

- What is the meaning of the stationary distribution π of a random walk?
- $\pi(i)$: the probability of being at node i after very large (infinite) number of steps
- $\pi = p_0 P^\infty$, where P is the transition matrix, p_0 the original vector
 - $P(i, j)$: probability of going from i to j in one step
 - $P^2(i, j)$: probability of going from i to j in two steps (probability of all paths of length 2)
 - $P^\infty(i, j) = \pi(j)$: probability of going from i to j in infinite steps – starting point does not matter.

The PageRank random walk

- Vanilla random walk
 - make the adjacency matrix stochastic and run a random walk

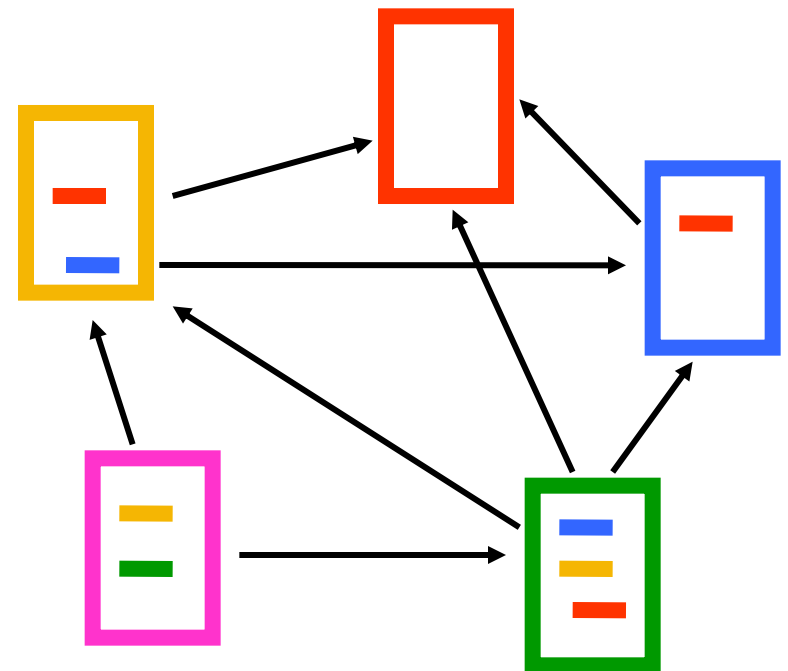
$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$



The PageRank random walk

- What about **sink** nodes?
 - what happens when the random walk moves to a node without any outgoing links?

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

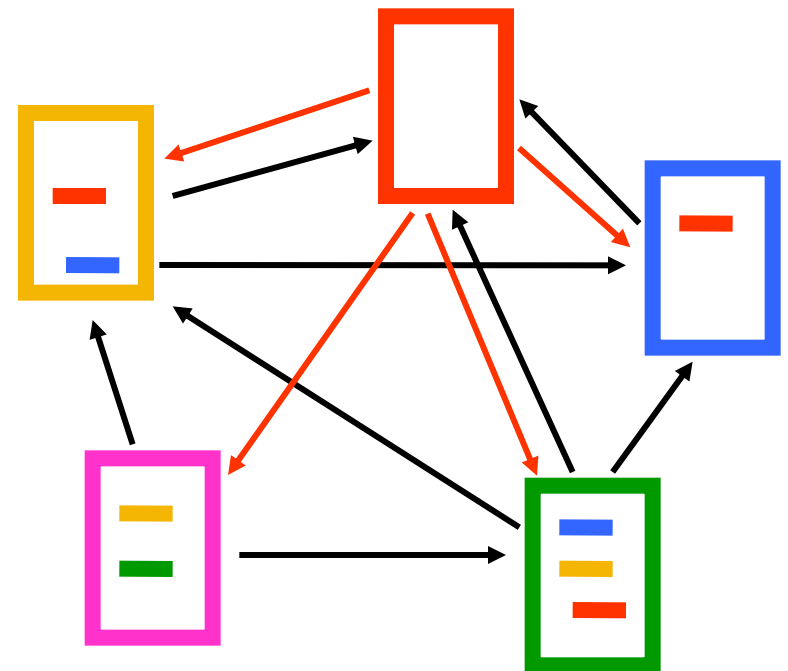


The PageRank random walk

- Replace these row vectors with a vector \mathbf{v}
 - typically, the uniform vector

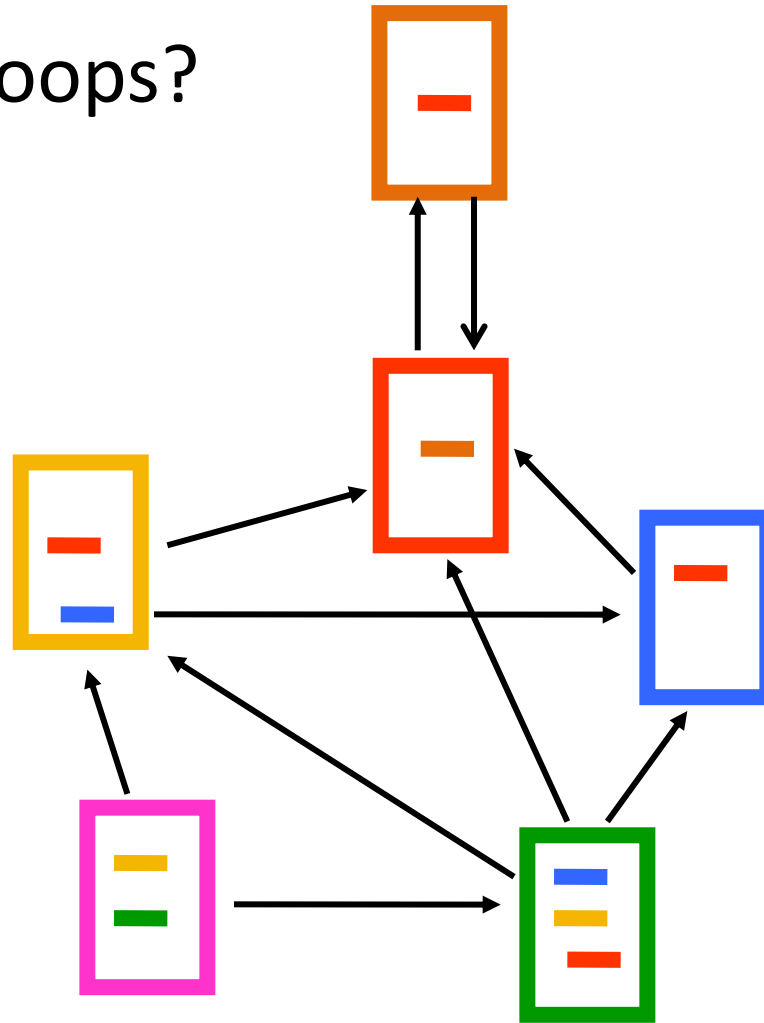
$$P' = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

$$P' = P + d\mathbf{v}^T \quad d = \begin{cases} 1 & \text{if } i \text{ is sink} \\ 0 & \text{otherwise} \end{cases}$$



The PageRank random walk

- What about loops?
 - Spider traps



The PageRank random walk

- Add a **random jump** to vector v with prob $1-\alpha$
 - typically, to a uniform vector
- Restarts after $1/(1-\alpha)$ steps in expectation
 - Guarantees irreducibility, convergence

$$P'' = \alpha \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 0 & 1/2 \end{bmatrix} + (1-\alpha) \begin{bmatrix} 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \end{bmatrix}$$

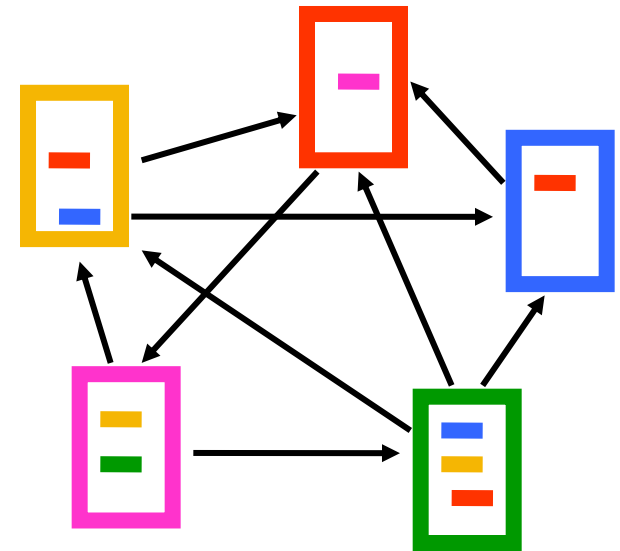
$P'' = \alpha P' + (1-\alpha)uv^T$, where u is the vector of all 1s

Random walk with restarts

PageRank algorithm [BP98]

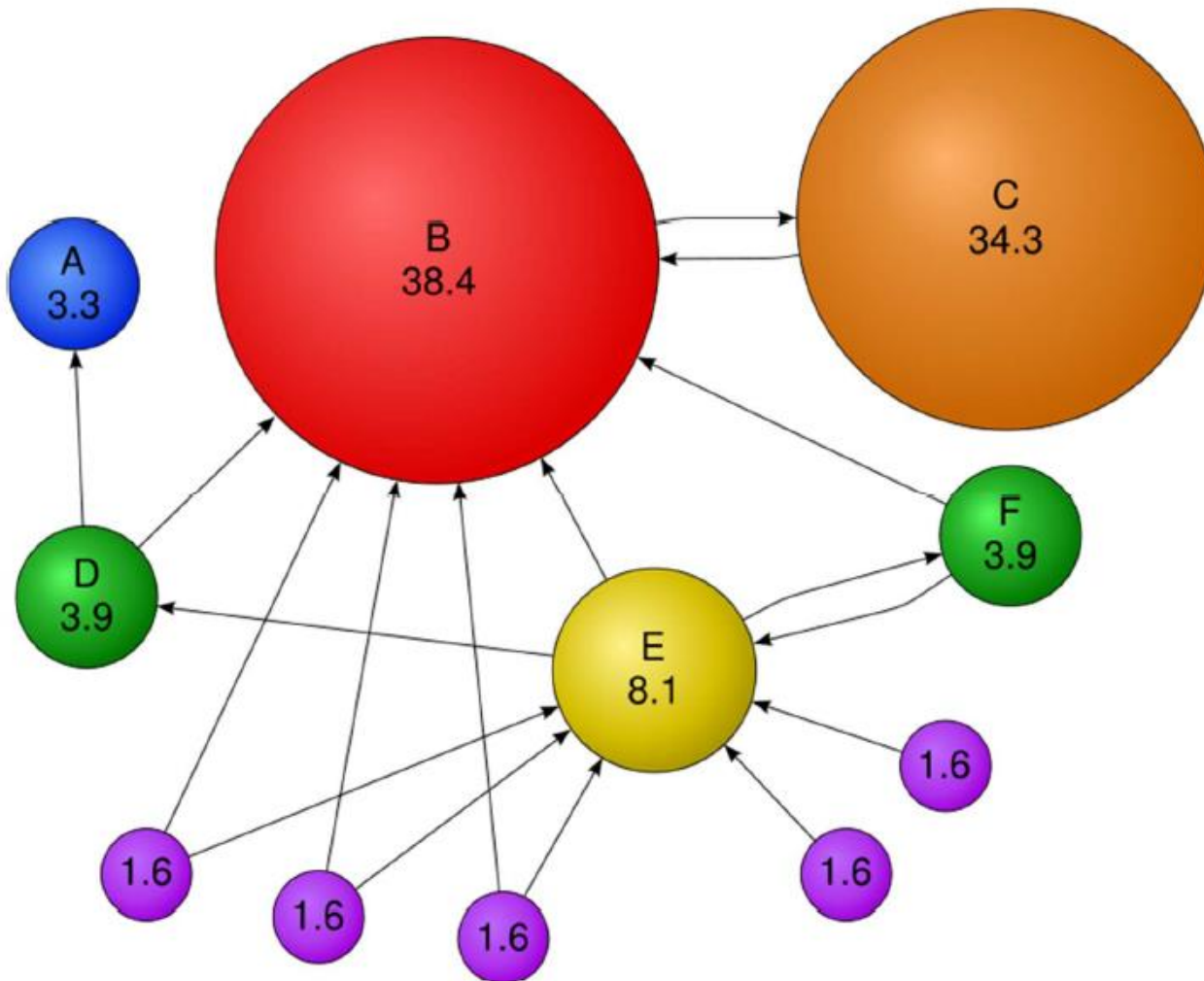
- The Random Surfer model
 - pick a page at random
 - with probability $1 - \alpha$ jump to a random page
 - with probability α follow a random outgoing link
- Rank according to the stationary distribution
- $$w_v = \alpha \sum_{u \rightarrow v} \frac{1}{d_{out}(u)} w_u + (1 - \alpha) \frac{1}{n}$$

$\alpha = 0.85$ in most cases
- We repeat this computation until converge



1. Red Page
2. Purple Page
3. Yellow Page
4. Blue Page
5. Green Page

PageRank: Example



Stationary distribution with random jump

- If v is the jump vector

$$p^0 = v$$

$$p^1 = \alpha p^0 P' + (1 - \alpha)v = \alpha v P' + (1 - \alpha)v$$

$$p^2 = \alpha p^1 P' + (1 - \alpha)v$$

$$= \alpha^2 v P'^2 + (1 - \alpha)v \alpha P' + (1 - \alpha)v$$

⋮

$$p^\infty = (1 - \alpha)v + (1 - \alpha)v \alpha P' + (1 - \alpha)v \alpha^2 P'^2 + \dots$$
$$= (1 - \alpha)(I - \alpha P')^{-1}$$

- With the random jump the shorter paths are more important, since the weight decreases exponentially
 - makes sense when thought of as a restart

Random walks with restarts

- If v is **not uniform**, we can bias the random walk towards the nodes that are **close** to v
- **Personalized PageRank**:
 - Restart the random walk from a specific node x
 - All nodes are ranked according to their closeness to x
- **Topic-Specific Pagerank**.
 - Restart the random walk from a specific set of nodes (e.g., nodes about a topic)
 - All nodes are ranked according to their closeness to the topic.
- **Random Walks with restarts** is a general technique for measuring closeness on graphs.

Effects of random jump

- Guarantees **convergence** to unique distribution
- Motivated by the concept of **random surfer**
- Offers additional flexibility
 - **personalization**
 - **anti-spam**
- Controls the **rate of convergence**
 - the second eigenvalue of matrix P'' is α

Random walks on undirected graphs

- For **undirected** graphs, the stationary distribution of a **random walk** is proportional to the degrees of the nodes
 - Thus in this case a random walk is the same as **degree popularity**
- This is **not longer true** if we do **random jumps**
 - Now the short paths play a greater role, and the previous distribution does not hold.
 - Random walks with restarts to a single node are commonly used on undirected graphs for measuring similarity between nodes

PageRank implementation

- Store the graph as a list of edges
- Keep current pagerank values and new pagerank values
- Go through edges and update the values of the destination nodes.
- Repeat until the difference between the pagerank vectors (L_1 or L_∞ difference) is below some small value ε .

A (Matlab-friendly) PageRank algorithm

- Performing vanilla power method is now too expensive – the matrix is not sparse

$$q^0 = v$$

$$t = 1$$

repeat

$$q^t = (P'')^T q^{t-1}$$

$$\delta = \|q^t - q^{t-1}\|$$

$$t = t + 1$$

until $\delta < \epsilon$

Efficient computation of $y = (P'')^T x$

$$y = \alpha P^T x$$

$$\beta = \|x\|_1 - \|y\|_1$$

$$y = y + \beta v$$

P = normalized adjacency matrix

$P' = P + dv^T$, where d_i is 1 if i is sink and 0 o.w.

$P'' = \alpha P' + (1-\alpha)uv^T$, where u is the vector of all 1s

PageRank history

- Huge advantage for Google in the early days
 - It gave a way to get an idea for the **value of a page**, which was useful in many different ways
 - Put an **order to the web**.
 - After a while it became clear that the **anchor text** was probably more important for ranking
 - Also, **link spam** became a new (dark) art
- Flood of research
 - Numerical analysis got rejuvenated
 - Huge number of variations
 - **Efficiency** became a great issue.
 - Huge number of applications in different fields
 - Random walk is often referred to as PageRank.

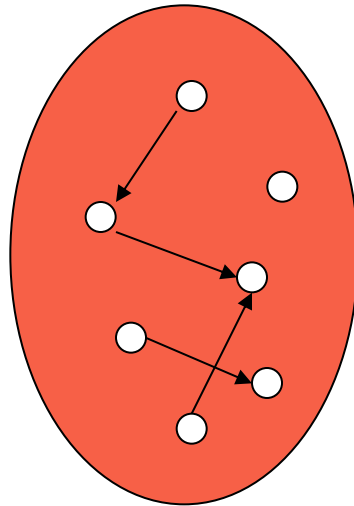
THE HITS ALGORITHM

The HITS algorithm

- Another algorithm proposed around the same time as PageRank for using the hyperlinks to rank pages
 - Kleinberg: then an intern at IBM Almaden
 - IBM never made anything out of it

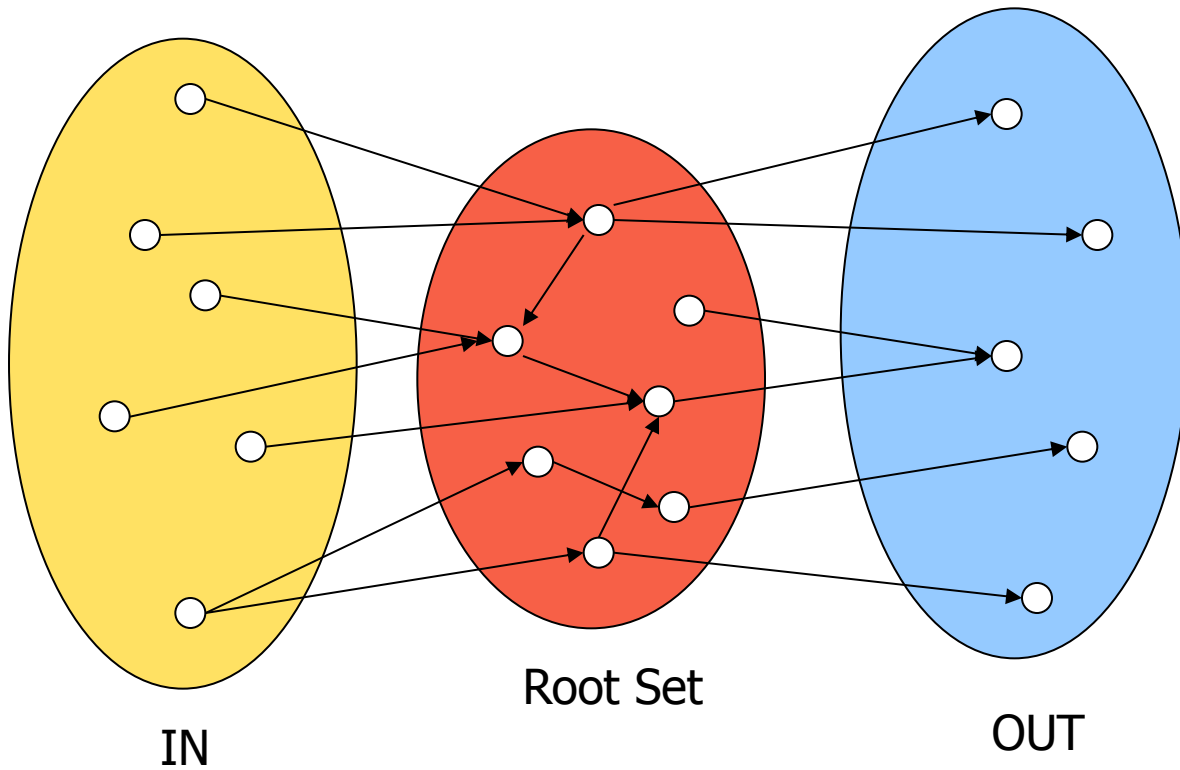
Query dependent input

Root set obtained from a text-only search engine

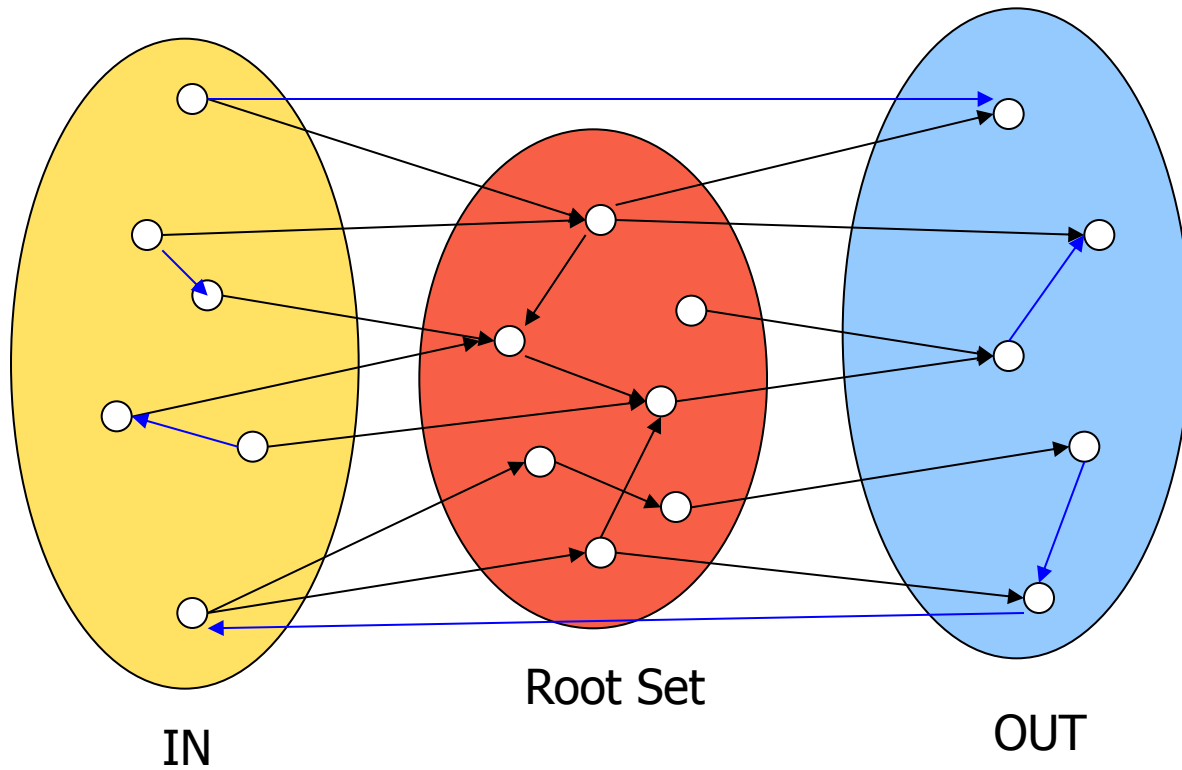


Root Set

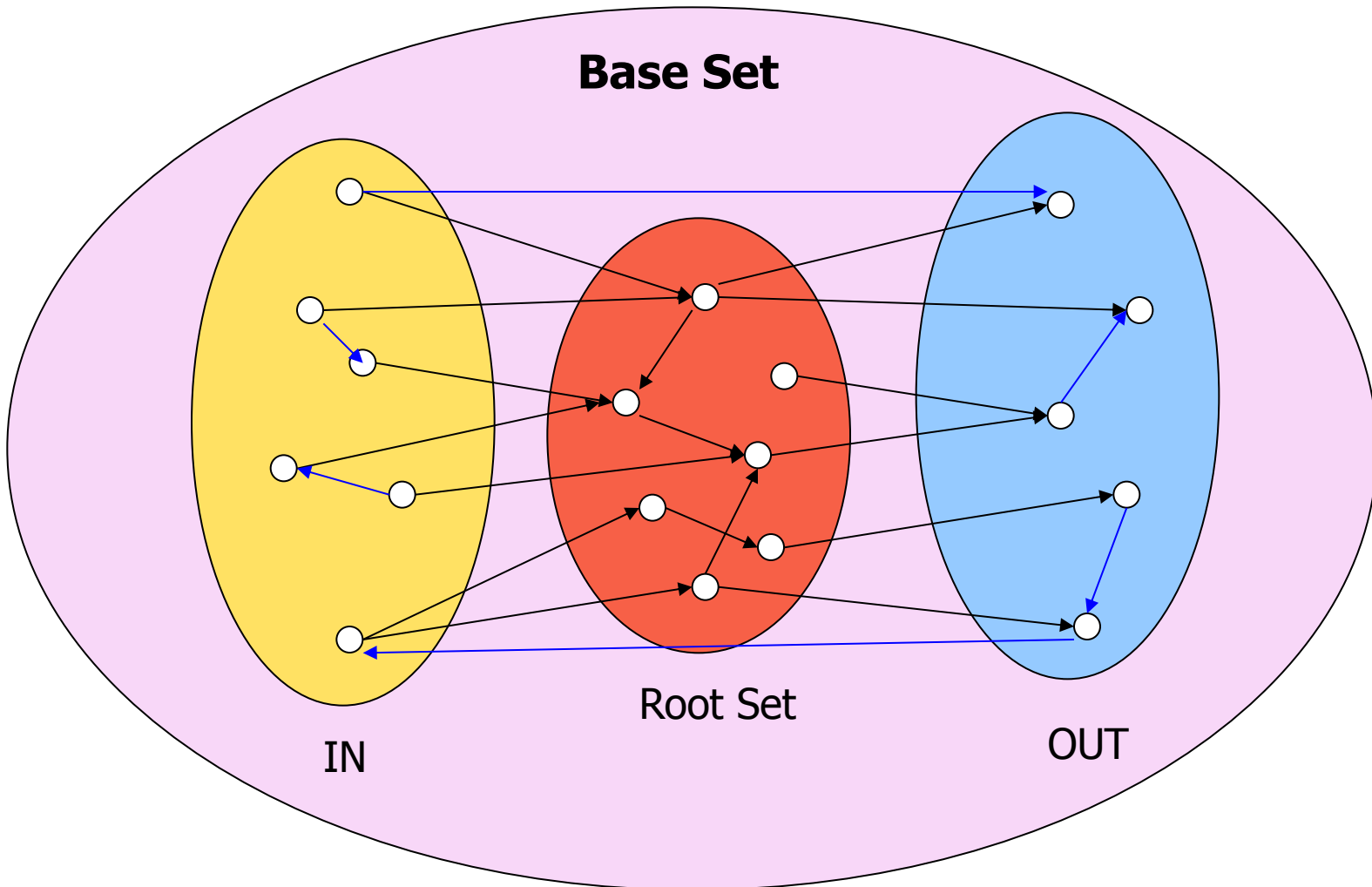
Query dependent input



Query dependent input

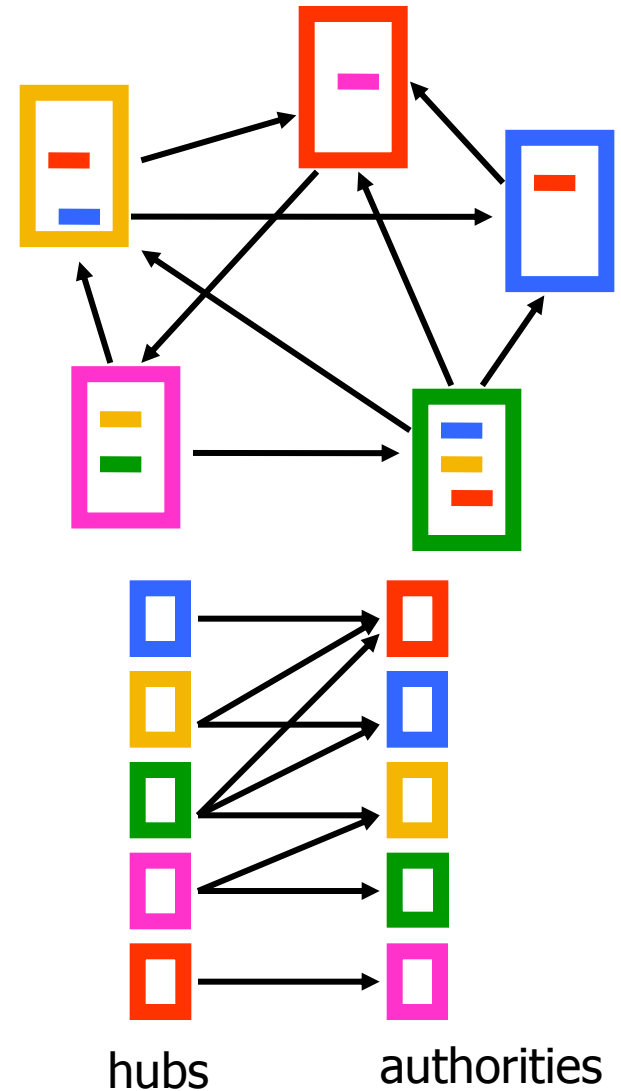


Query dependent input



Hubs and Authorities [K98]

- Authority is not necessarily transferred directly between authorities
- Pages have double identity
 - hub identity
 - authority identity
- Good hubs point to good authorities
- Good authorities are pointed by good hubs



Hubs and Authorities

- Two kind of weights:
 - Hub weight
 - Authority weight
- The hub weight is the sum of the authority weights of the authorities pointed to by the hub
- The authority weight is the sum of the hub weights that point to this authority.

HITS Algorithm

- Initialize all weights to 1.
- Repeat until convergence
 - *O* operation : hubs collect the weight of the authorities

$$h_i = \sum_{j:i \rightarrow j} a_j$$

- *I* operation: authorities collect the weight of the hubs

$$a_i = \sum_{j:j \rightarrow i} h_j$$

- Normalize weights under some norm

HITS and eigenvectors

- The HITS algorithm is a **power-method** eigenvector computation
- In vector terms
 - $a^t = A^T h^{t-1}$ and $h^t = A a^{t-1}$
 - $a^t = A^T A a^{t-1}$ and $h^t = A A^T h^{t-1}$
 - Repeated iterations will converge to the eigenvectors
- The **authority** weight vector a is the **eigenvector** of $A^T A$ and the **hub** weight vector h is the **eigenvector** of $A A^T$
- The vectors a and h are called the **singular vectors** of the matrix A

Singular Value Decomposition

$$A = U \Sigma V^T = \begin{bmatrix} \vec{u}_1 & \vec{u}_2 & \cdots & \vec{u}_r \end{bmatrix} \begin{bmatrix} \sigma_1 & & & \\ & \sigma_2 & & \\ & & \ddots & \\ & & & \sigma_r \end{bmatrix} \begin{bmatrix} \vec{v}_1 \\ \vec{v}_2 \\ \vdots \\ \vec{v}_r \end{bmatrix}$$

$[n \times r] \quad [r \times r] \quad [r \times n]$

- r : rank of matrix A
- $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_r$: singular values (square roots of eig-vals $AA^T, A^T A$)
- $\vec{u}_1, \vec{u}_2, \dots, \vec{u}_r$: left singular vectors (eig-vectors of AA^T)
- $\vec{v}_1, \vec{v}_2, \dots, \vec{v}_r$: right singular vectors (eig-vectors of $A^T A$)

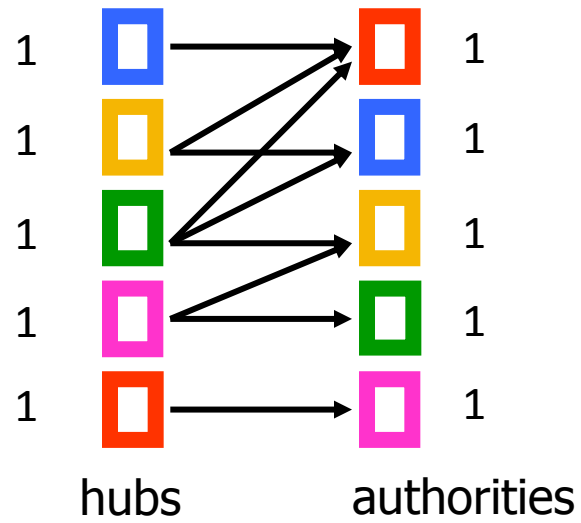
$$A = \sigma_1 \vec{u}_1 \vec{v}_1^T + \sigma_2 \vec{u}_2 \vec{v}_2^T + \cdots + \sigma_r \vec{u}_r \vec{v}_r^T$$

Why does the Power Method work?

- If a matrix R is real and symmetric, it has real eigenvalues and eigenvectors: $(\lambda_1, w_1), (\lambda_2, w_2), \dots, (\lambda_r, w_r)$
 - r is the rank of the matrix
 - $|\lambda_1| \geq |\lambda_2| \geq \dots \geq |\lambda_r|$
- For any matrix R , the eigenvectors w_1, w_2, \dots, w_r of R define a **basis** of the vector space
 - For any vector x , $Rx = \alpha_1 w_1 + \alpha_2 w_2 + \dots + \alpha_r w_r$
- After t multiplications we have:
 - $R^t x = \lambda_1^{t-1} \alpha_1 w_1 + \lambda_2^{t-1} \alpha_2 w_2 + \dots + \lambda_r^{t-1} \alpha_r w_r$
- Normalizing (divide by λ_1^{t-1}) leaves only the term w_1 .

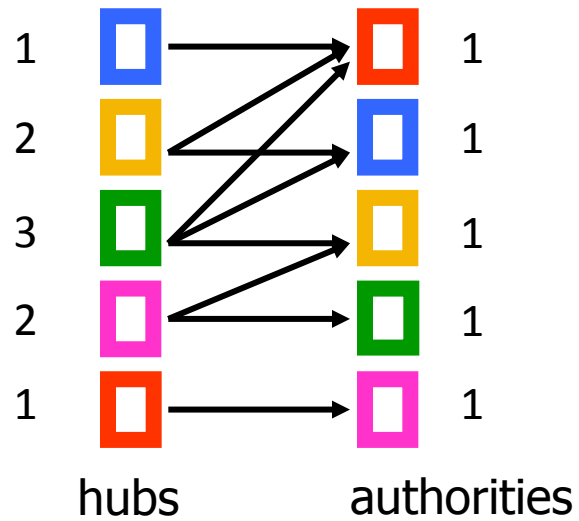
Example

Initialize



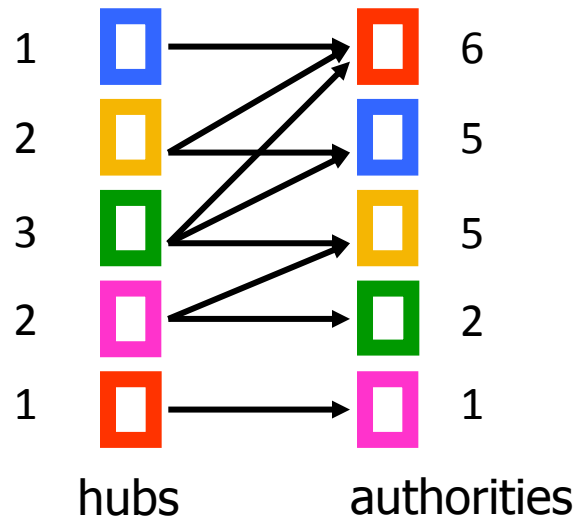
Example

Step 1: O operation



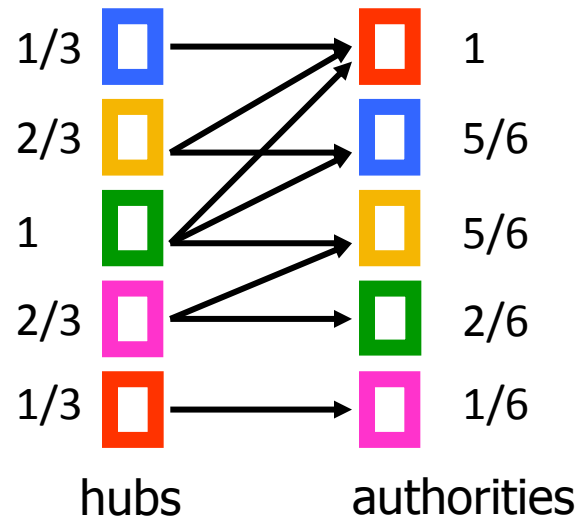
Example

Step 1: I operation



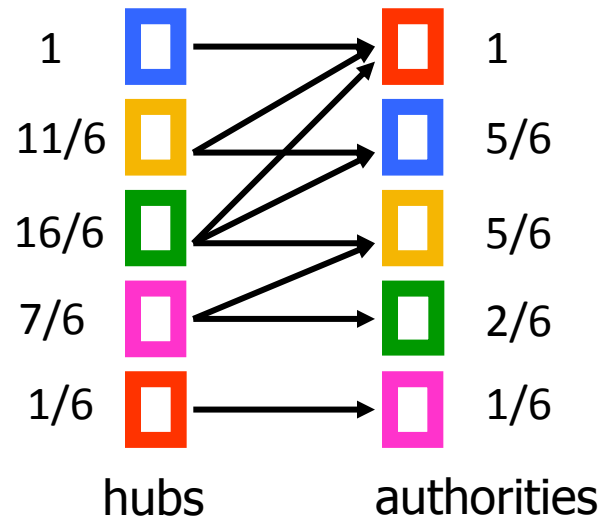
Example

Step 1: Normalization (Max norm)



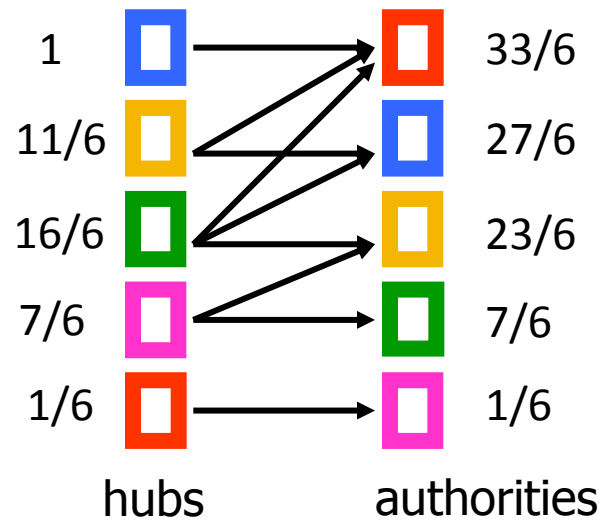
Example

Step 2: 0 step



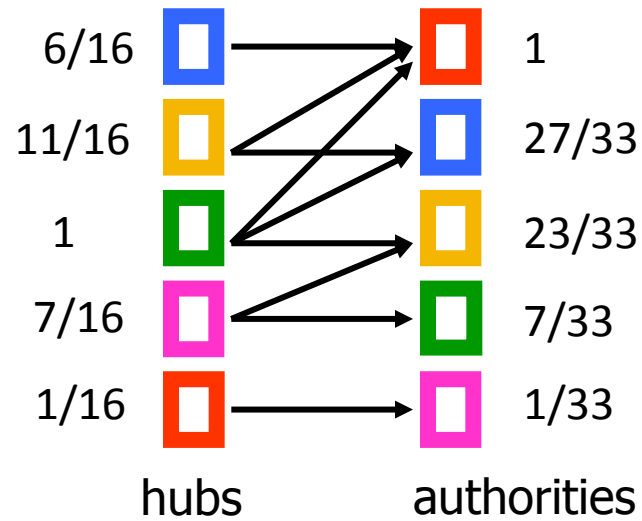
Example

Step 2: 1 step



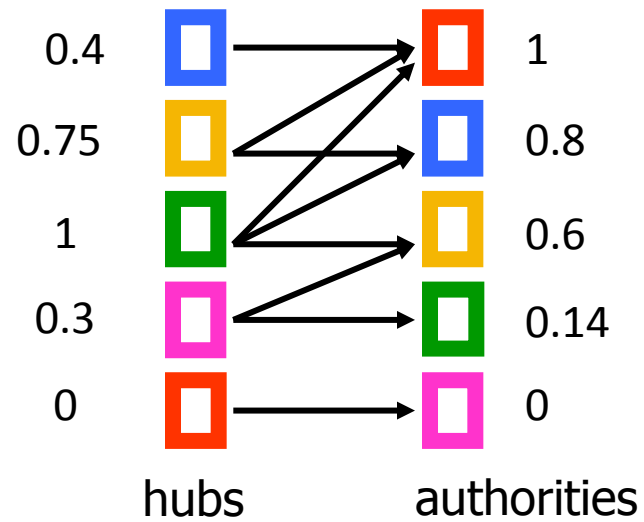
Example

Step 2: Normalization



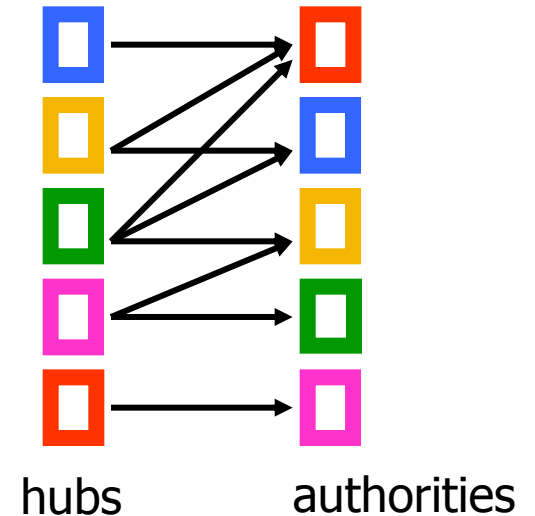
Example

Convergence



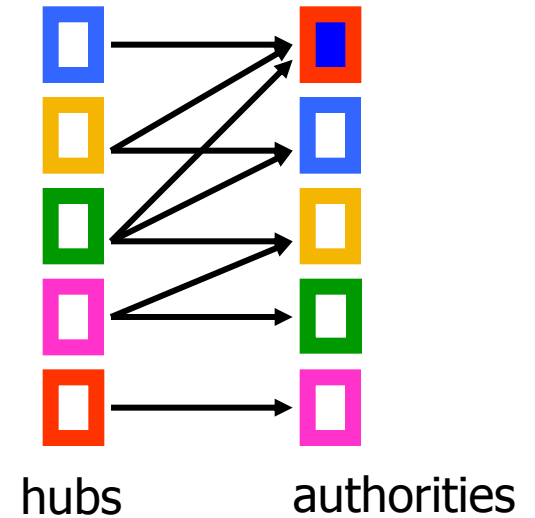
The SALSA algorithm

- Perform a random walk on the bipartite graph of hubs and authorities alternating between the two



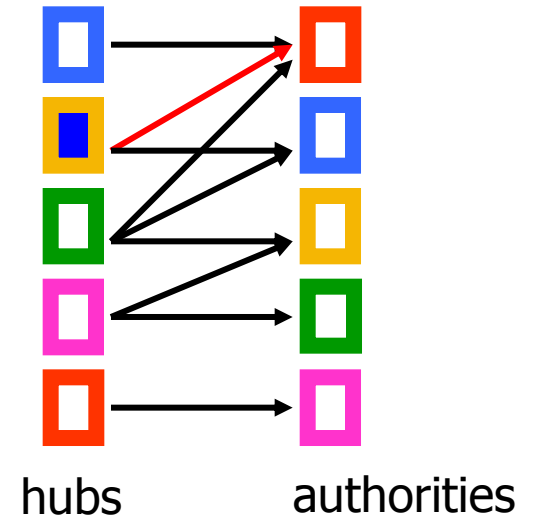
The SALSA algorithm

- Start from an authority chosen uniformly at random
 - e.g. the red authority



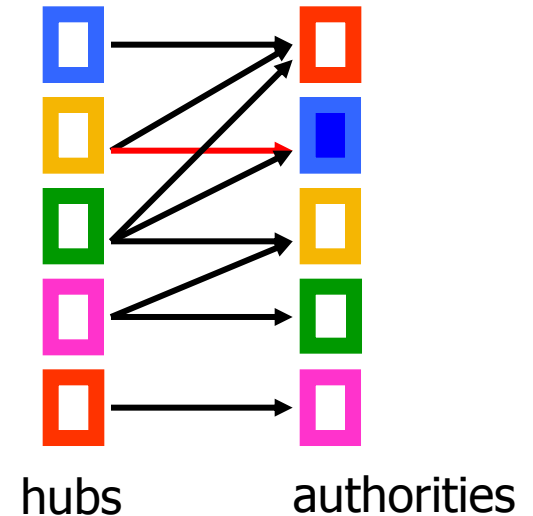
The SALSA algorithm

- Start from an authority chosen uniformly at random
 - e.g. the red authority
- Choose one of the in-coming links uniformly at random and move to a hub
 - e.g. move to the yellow authority with probability $1/3$



The SALSA algorithm

- Start from an authority chosen uniformly at random
 - e.g. the red authority
- Choose one of the in-coming links uniformly at random and move to a hub
 - e.g. move to the yellow authority with probability $1/3$
- Choose one of the out-going links uniformly at random and move to an authority
 - e.g. move to the blue authority with probability $1/2$



The SALSA algorithm

- Formally we have probabilities:
 - a_i : probability of being at authority i
 - h_j : probability of being at hub j
- The probability of being at authority i is computed as:

$$a_i = \sum_{j \in N_{in}(i)} \frac{1}{d_{out}(j)} h_j$$

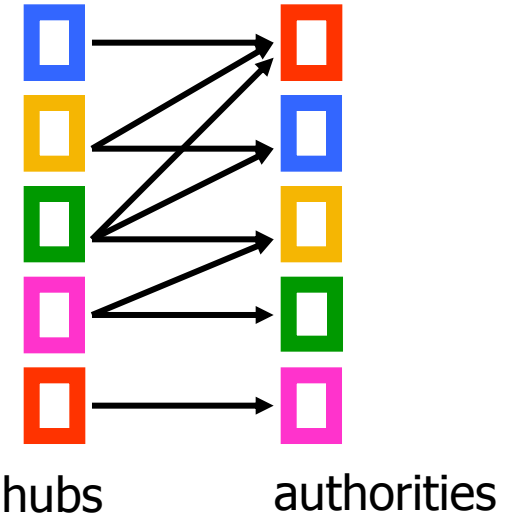
- The probability of being at hub j is computed as

$$h_j = \sum_{i \in N_{out}(j)} \frac{1}{d_{in}(i)} a_i$$

- Repeated computation converges

The SALSA algorithm [LM00]

- In matrix terms
 - A_c = the matrix A where **columns** are normalized to sum to 1
 - A_r = the matrix A where **rows** are normalized to sum to 1
- The hub computation
 - $h = A_c a$
- The authority computation
 - $a = A_r^T h = A_r^T A_c a$
- In MC terms the transition matrix
 - $P = A_r A_c^T$



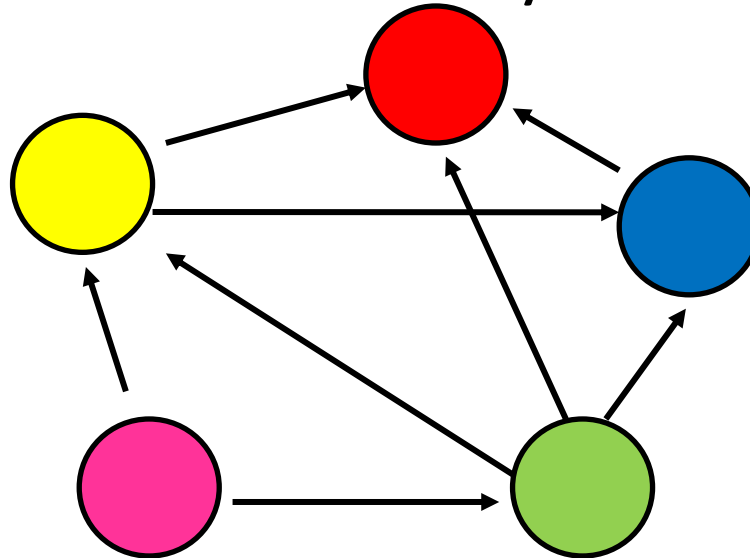
$$h_2 = 1/3 a_1 + 1/2 a_2$$

$$a_1 = h_1 + 1/2 h_2 + 1/3 h_3$$

ABSORBING RANDOM WALKS
LABEL PROPAGATION
OPINION FORMATION ON SOCIAL
NETWORKS

Random walk with absorbing nodes

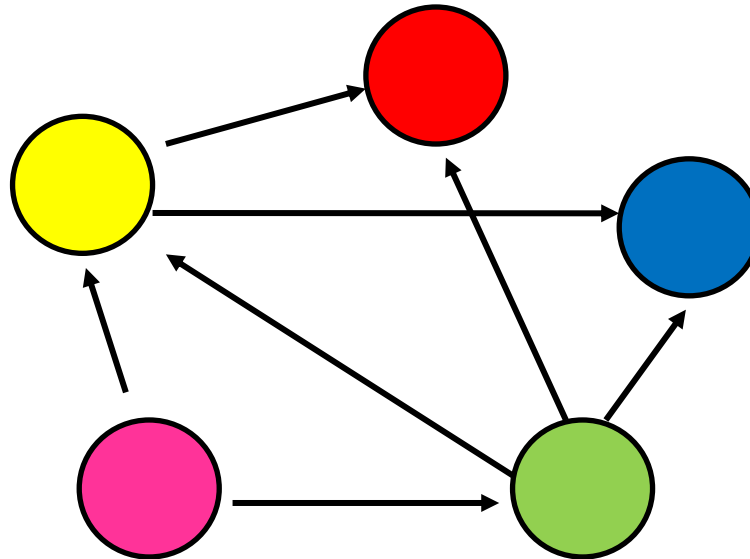
- What happens if we do a random walk on this graph? What is the stationary distribution?



- All the probability mass on the red **sink** node:
 - The red node is an **absorbing node**

Random walk with absorbing nodes

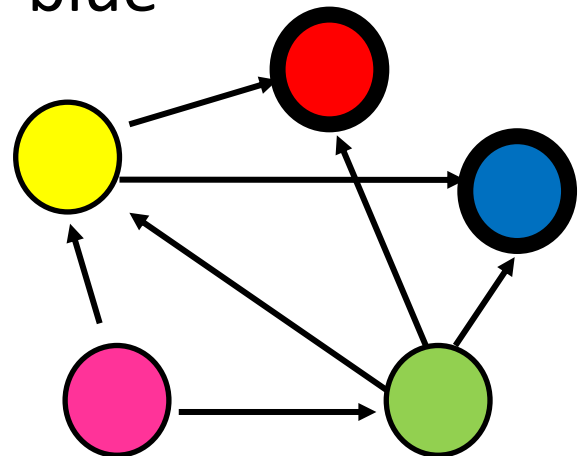
- What happens if we do a random walk on this graph? What is the stationary distribution?



- There are two absorbing nodes: the red and the blue.
- The probability mass will be **divided** between the two

Absorption probability

- If there are more than one **absorbing nodes** in the graph a random walk that starts from a **non-absorbing** node will be absorbed in one of them with some probability
 - The **probability of absorption** gives an estimate of how **close** the node is to red or blue



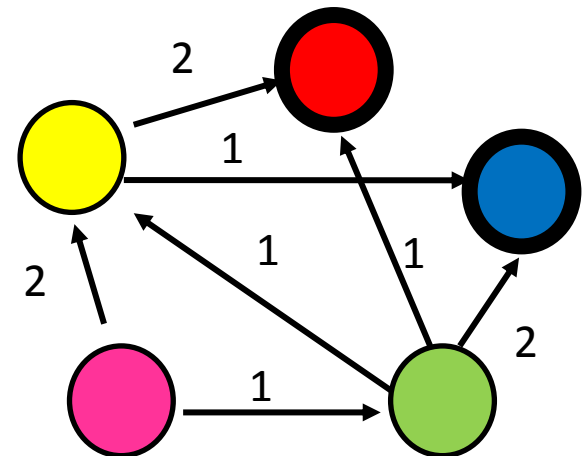
Absorption probability

- Computing the probability of being absorbed:
 - The **absorbing nodes** have probability 1 of being absorbed in themselves and zero of being absorbed in another node.
 - For the **non-absorbing nodes**, take the (weighted) average of the absorption probabilities of your neighbors
 - if one of the neighbors is the absorbing node, it has probability 1
 - Repeat until convergence (= very small change in probs)

$$P(\text{Red}|\text{Pink}) = \frac{2}{3}P(\text{Red}|\text{Yellow}) + \frac{1}{3}P(\text{Red}|\text{Green})$$

$$P(\text{Red}|\text{Green}) = \frac{1}{4}P(\text{Red}|\text{Yellow}) + \frac{1}{4}$$

$$P(\text{Red}|\text{Yellow}) = \frac{2}{3}$$



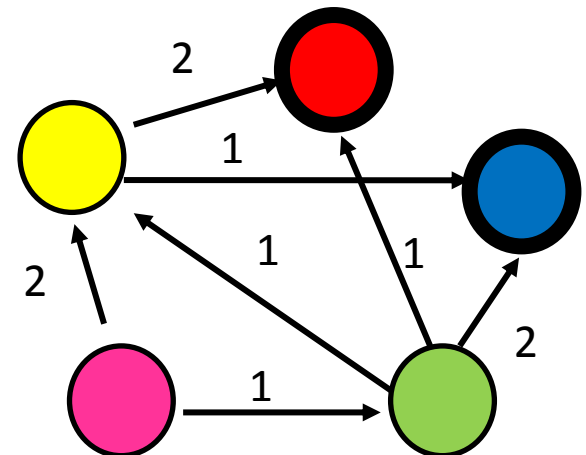
Absorption probability

- Computing the probability of being absorbed:
 - The **absorbing nodes** have probability 1 of being absorbed in themselves and zero of being absorbed in another node.
 - For the **non-absorbing nodes**, take the (weighted) average of the absorption probabilities of your neighbors
 - if one of the neighbors is the absorbing node, it has probability 1
 - Repeat until convergence (= very small change in probs)

$$P(\text{Blue}|\text{Pink}) = \frac{2}{3}P(\text{Blue}|\text{Yellow}) + \frac{1}{3}P(\text{Blue}|\text{Green})$$

$$P(\text{Blue}|\text{Green}) = \frac{1}{4}P(\text{Blue}|\text{Yellow}) + \frac{1}{2}$$

$$P(\text{Blue}|\text{Yellow}) = \frac{1}{3}$$

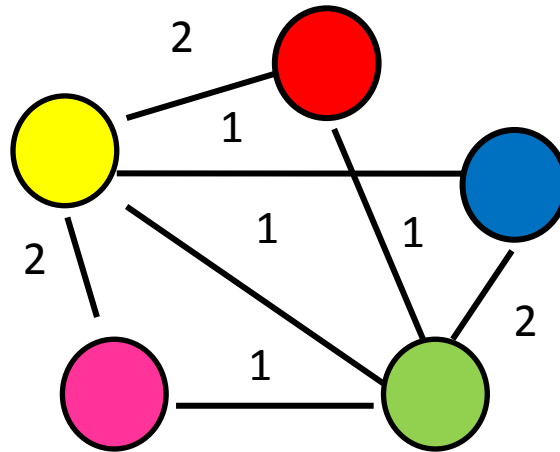


Why do we care?

- Why do we care to compute the absorption probability to sink nodes?
- Given a graph (**directed** or **undirected**) we can choose to **make** some nodes **absorbing**.
 - Simply **direct** all edges incident on the chosen nodes towards them.
- The absorbing random walk provides a measure of **proximity** of non-absorbing nodes to the chosen nodes.
 - Useful for **understanding** proximity in graphs
 - Useful for **propagation** in the graph
 - E.g, on a social network some nodes have **high income**, some have **low income**, to which income class is a non-absorbing node closer?

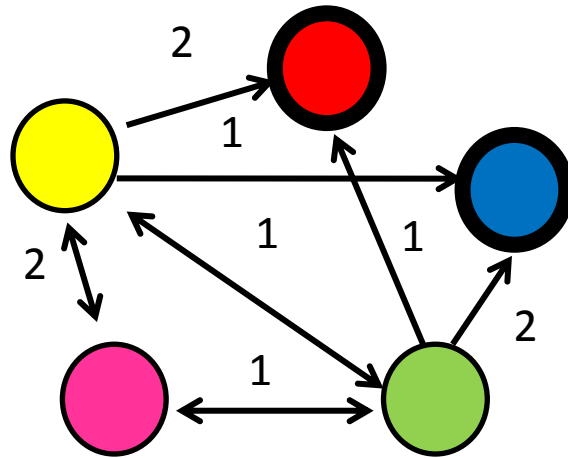
Example

- In this **undirected** graph we want to learn the proximity of nodes to the **red** and **blue** nodes



Example

- Make the nodes absorbing



Absorption probability

- Compute the absorption probabilities for red and blue

$$P(\text{Red}|\text{Pink}) = \frac{2}{3}P(\text{Red}|\text{Yellow}) + \frac{1}{3}P(\text{Red}|\text{Green})$$

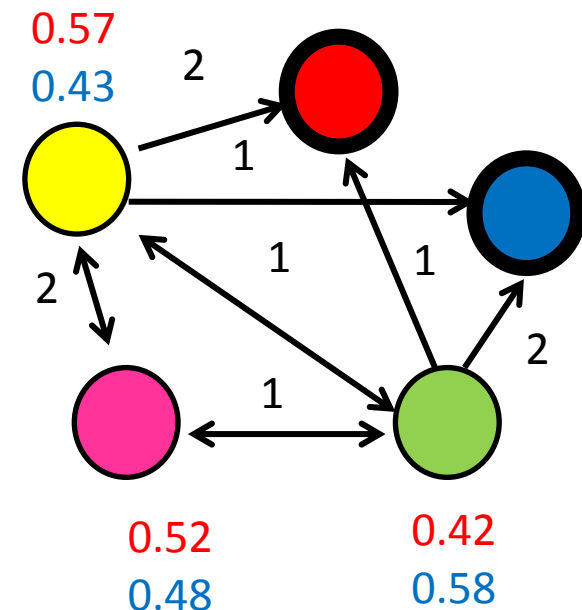
$$P(\text{Red}|\text{Green}) = \frac{1}{5}P(\text{Red}|\text{Yellow}) + \frac{1}{5}P(\text{Red}|\text{Pink}) + \frac{1}{5}$$

$$P(\text{Red}|\text{Yellow}) = \frac{1}{6}P(\text{Red}|\text{Green}) + \frac{1}{3}P(\text{Red}|\text{Pink}) + \frac{1}{3}$$

$$P(\text{Blue}|\text{Pink}) = 1 - P(\text{Red}|\text{Pink})$$

$$P(\text{Blue}|\text{Green}) = 1 - P(\text{Red}|\text{Green})$$

$$P(\text{Blue}|\text{Yellow}) = 1 - P(\text{Red}|\text{Yellow})$$

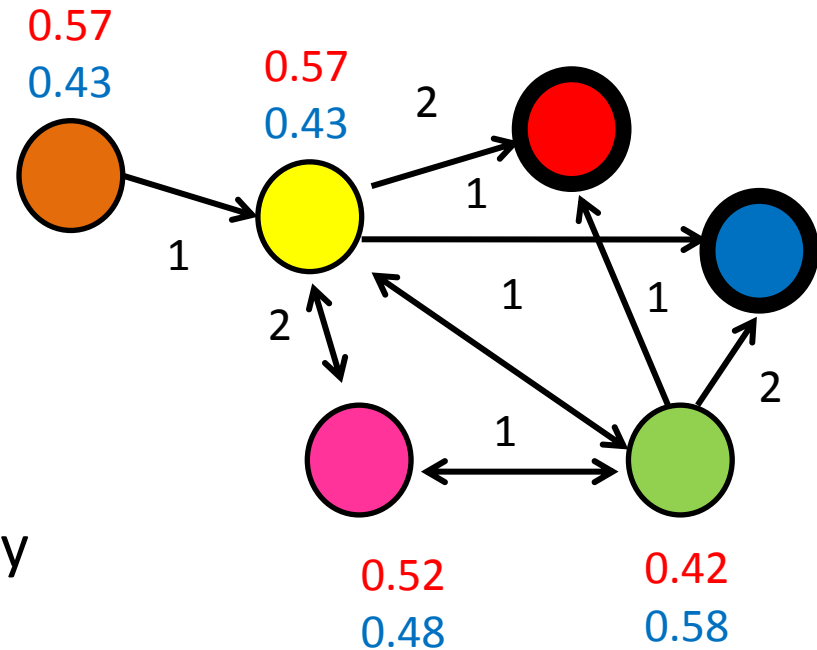


Penalizing long paths

- The orange node has the same probability of reaching red and blue as the yellow one

$$P(\text{Red}|\text{Orange}) = P(\text{Red}|\text{Yellow})$$

$$P(\text{Blue}|\text{Orange}) = P(\text{Blue}|\text{Yellow})$$



- Intuitively though it is further away

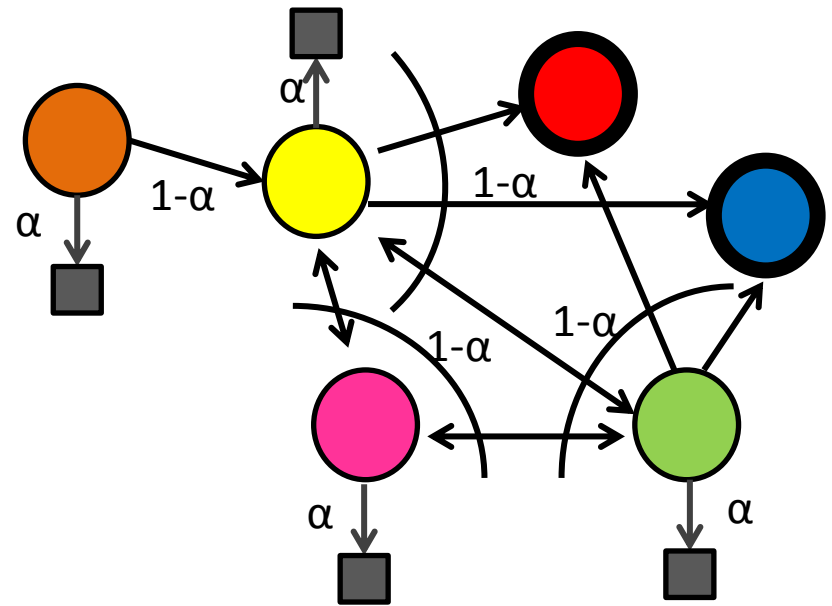
Penalizing long paths

- Add an **universal absorbing node** to which each node gets absorbed with probability α .

With probability α the random walk **dies**

With probability $(1-\alpha)$ the random walk continues as before

The longer the path from a node to an absorbing node the more likely the random walk dies along the way, the lower the absorption probability



$$P(\text{Red}|\text{Green}) = (1 - \alpha) \left(\frac{1}{5} P(\text{Red}|\text{Yellow}) + \frac{1}{5} P(\text{Red}|\text{Pink}) + \frac{1}{5} \right)$$

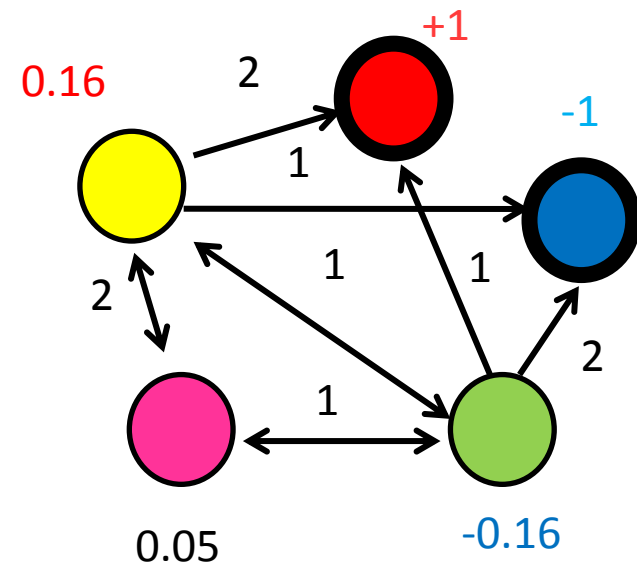
Propagating values

- Assume that **Red** has a positive value and **Blue** a negative value
 - Positive/Negative **class**, Positive/Negative **opinion**
- We can compute a value for all the other nodes in the same way
 - This is the **expected** value for the node

$$V(\text{Pink}) = \frac{2}{3}V(\text{Yellow}) + \frac{1}{3}V(\text{Green})$$

$$V(\text{Green}) = \frac{1}{5}V(\text{Yellow}) + \frac{1}{5}V(\text{Pink}) + \frac{1}{5} - \frac{2}{5}$$

$$V(\text{Yellow}) = \frac{1}{6}V(\text{Green}) + \frac{1}{3}V(\text{Pink}) + \frac{1}{3} - \frac{1}{6}$$



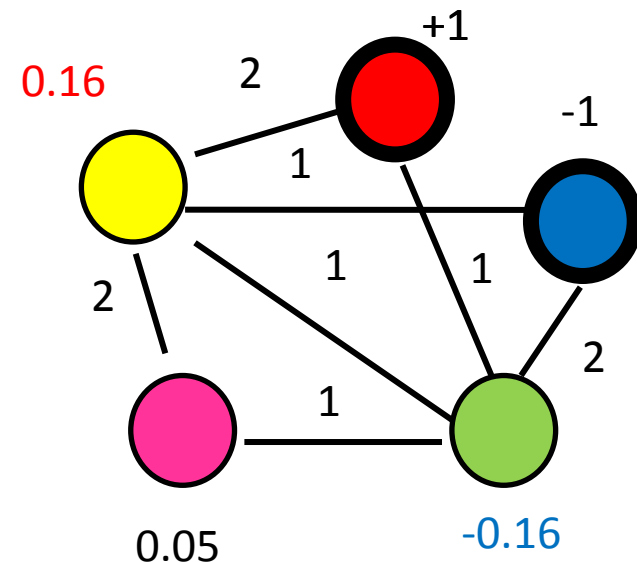
Electrical networks and random walks

- Our graph corresponds to an **electrical network**
- There is a positive **voltage** of **+1** at the Red node, and a negative voltage **-1** at the Blue node
- There are **resistances** on the edges **inversely proportional** to the weights (or **conductance proportional** to the weights)
- The computed values are the **voltages** at the nodes

$$V(\text{Pink}) = \frac{2}{3}V(\text{Yellow}) + \frac{1}{3}V(\text{Green})$$

$$V(\text{Green}) = \frac{1}{5}V(\text{Yellow}) + \frac{1}{5}V(\text{Pink}) + \frac{1}{5} - \frac{2}{5}$$

$$V(\text{Yellow}) = \frac{1}{6}V(\text{Green}) + \frac{1}{3}V(\text{Pink}) + \frac{1}{3} - \frac{1}{6}$$



Opinion formation

- The value propagation can be used as a model of opinion formation.
- Model:
 - Opinions are **values** in $[-1,1]$
 - Every user u has an **internal opinion** s_u , and **expressed opinion** z_u .
 - The expressed opinion **minimizes** the **personal cost** of user u :

$$c(z_u) = (s_u - z_u)^2 + \sum_{v:v \text{ is a friend of } u} w_u (z_u - z_v)^2$$

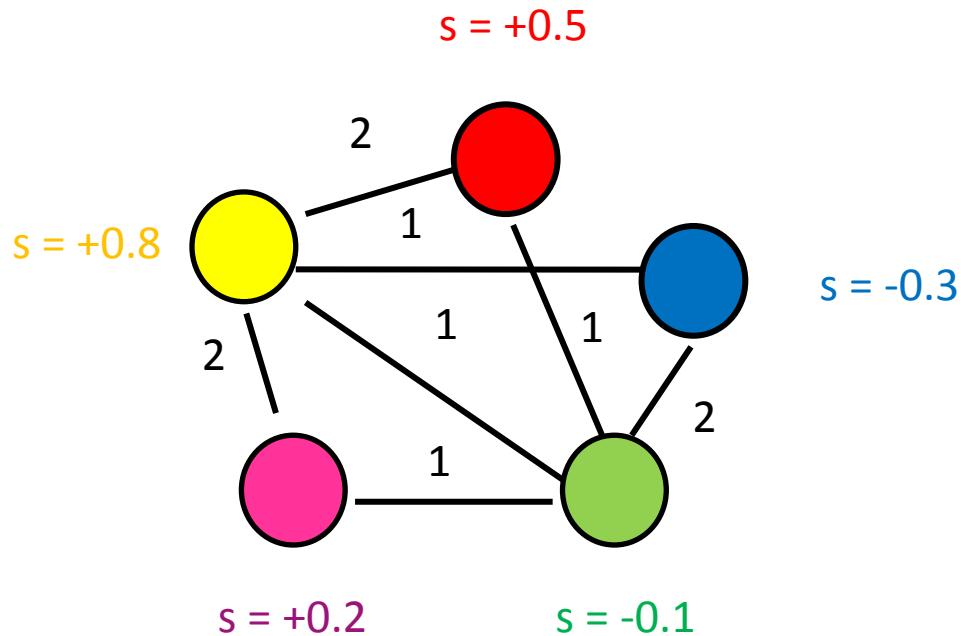
- Minimize deviation from your beliefs and conflicts with the society
- If every user tries **independently (selfishly)** to minimize their personal cost then the best thing to do is to set z_u to the **average** of all opinions:

$$z_u = \frac{s_u + \sum_{v:v \text{ is a friend of } u} w_u z_v}{1 + \sum_{v:v \text{ is a friend of } u} w_u}$$

- This is the same as the value propagation we described before!

Example

- Social network with **internal opinions**



Example

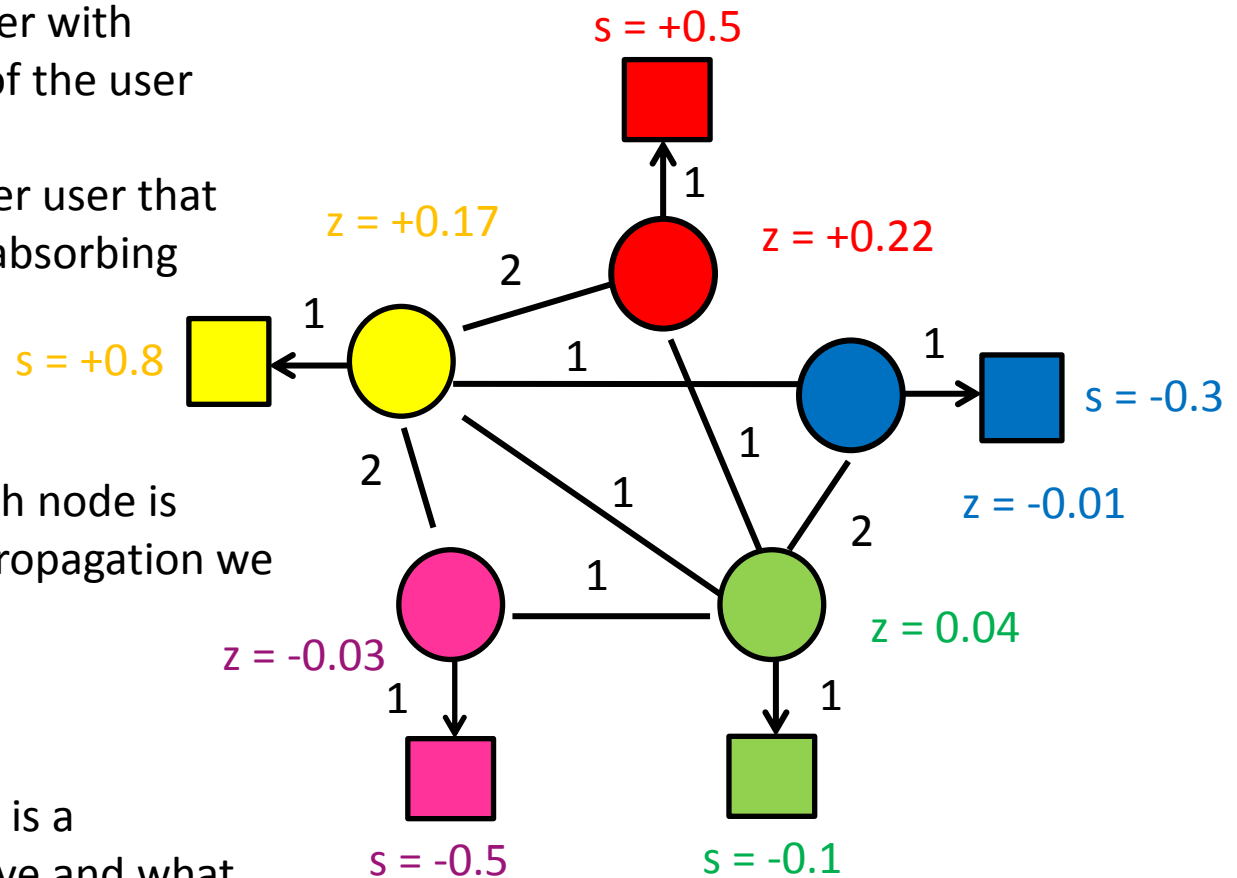
One absorbing node per user with value the **internal opinion** of the user

One non-absorbing node per user that links to the corresponding absorbing node

The **external opinion** for each node is computed using the value propagation we described before

- Repeated averaging

Intuitive model: my opinion is a combination of what I believe and what my social network believes.



Transductive learning

- If we have a graph of relationships and some **labels** on some nodes we can **propagate** them to the remaining nodes
 - Make the labeled nodes to be absorbing and compute the probability for the rest of the graph
 - E.g., a social network where some people are tagged as spammers
 - E.g., the movie-actor graph where some movies are tagged as action or comedy.
- This is a form of **semi-supervised learning**
 - We make use of the unlabeled data, and the relationships
- It is also called **transductive learning** because it does not produce a model, but just labels the unlabeled data that is at hand.
 - Contrast to **inductive learning** that learns a model and can label any new example

Implementation details

- Implementation is in many ways similar to the PageRank implementation
 - For an edge (u, v) instead of updating the value of v we update the value of u .
 - The value of a node is the average of its neighbors
 - We need to check for the case that a node u is absorbing, in which case the value of the node is not updated.
 - Repeat the updates until the change in values is very small.