# DATA MINING LECTURE 7

Minimum Description Length Principle Information Theory Co-Clustering

# MINIMUM DESCRIPTION LENGTH

### Occam's razor

- Most data mining tasks can be described as creating a model for the data
  - E.g., the EM algorithm models the data as a mixture of Gaussians, the K-means models the data as a set of centroids.
  - Model vs Hypothesis
- What is the right model?
- Occam's razor: All other things being equal, the simplest model is the best.
  - A good principle for life as well

### Occam's Razor and MDL

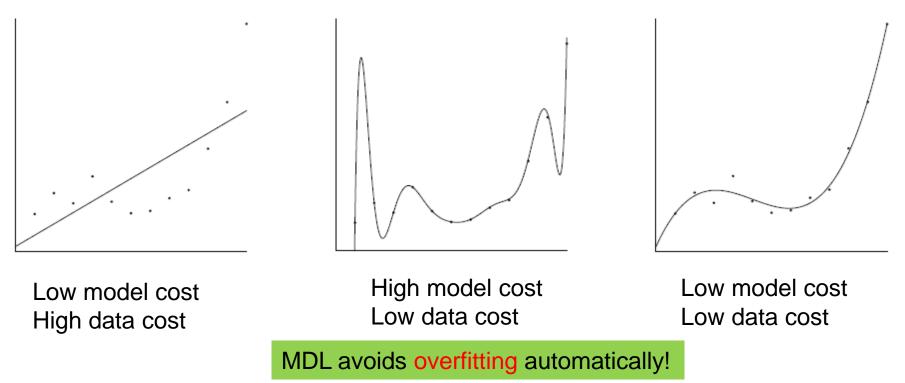
- What is a simple model?
- Minimum Description Length Principle: Every model provides a (lossless) encoding of our data. The model that gives the shortest encoding (best compression) of the data is the best.
  - Related: Kolmogorov complexity. Find the shortest program that produces the data (uncomputable).
  - MDL restricts the family of models considered
  - Encoding cost: cost of party A to transmit to party B the data.

## Minimum Description Length (MDL)

- The description length consists of two terms
  - The cost of describing the model (model cost)
  - The cost of describing the data given the model (data cost).
  - L(D) = L(M) + L(D|M)
- There is a tradeoff between the two costs
  - Very complex models describe the data in a lot of detail but are expensive to describe
  - Very simple models are cheap to describe but require a lot of work to describe the data given the model

### Example

- Regression: find the polynomial for describing the data
  - Complexity of the model vs. Goodness of fit



Source: Grnwald et al. (2005) Advances in Minimum Description Length: Theory and Applications.

### MDL and Data Mining

- Why does the shorter encoding make sense?
  - Shorter encoding implies regularities in the data
  - Regularities in the data imply patterns
  - Patterns are interesting
- Example

• Short description length, just repeat 12 times 00001

• Random sequence, no patterns, no compression

## MDL and Clustering

- If we have a clustering of the data, we can transmit the clusters instead of the data
  - We need to transmit the description of the clusters
  - And the data within each cluster.
- If we have a good clustering the transmission cost is low
  - Why?
  - What happens if all elements of the cluster are identical?
  - What happens if we have very few elements per cluster?
    Homogeneous clusters are cheaper to

Homogeneous clusters are cheaper to encode But we should not have too many

### Issues with MDL

- What is the right model family?
  - This determines the kind of solutions that we can have
    - E.g., polynomials
    - Clusterings
- What is the encoding cost?
  - Determines the function that we optimize
  - Information theory

## **INFORMATION THEORY**

A short introduction

### Encoding

Consider the following sequence

#### AAABBBAAACCCABACAABBAACCABAC

- Suppose you wanted to encode it in binary form, how would you do it?
  - 50% A 25% B 25% C

A is 50% of the sequence We should give it a shorter representation  $\begin{array}{c} A \rightarrow 0 \\ B \rightarrow 10 \\ C \rightarrow 11 \end{array}$ 

This is actually provably the best encoding!

## Encoding

Prefix Codes: no codeword is a prefix of another

$A \rightarrow 0$	Uniquely directly decodable
$B \rightarrow 10$ $C \rightarrow 11$	For every code we can find a prefix code of equal length

- Codes and Distributions: There is one to one mapping between codes and distributions
  - If P is a distribution over a set of elements (e.g., {A,B,C}) then there exists a (prefix) code C where  $L_C(x) = -\lceil \log P(x) \rceil$ ,  $x \in \{A, B, C\}$
  - For every (prefix) code C of elements {A,B,C}, we can define a distribution  $P(x) = 2^{-C(x)}$
- The code defined has the smallest average codelength!

## Entropy

- Suppose we have a random variable X that takes n distinct values  $X = \{x_1, x_2, \dots, x_n\}$ that have probabilities  $P(X) = \{p_1, \dots, p_n\}$
- This defines a code C with  $L_C(x_i) = -\lceil \log p_i \rceil$ . The average codelength is

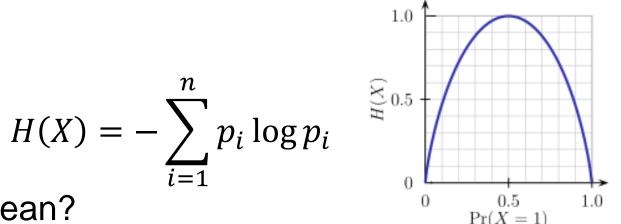
$$-\sum_{i=1}^{n} p_i \lceil \log p_i \rceil$$

• This (more or less) is the entropy H(X) of the random variable X

$$H(X) = -\sum_{i=1}^{n} p_i \log p_i$$

- Shannon's theorem: The entropy is a lower bound on the average codelength of any code that encodes the distribution P(X)
  - When encoding N numbers drawn from P(X), the best encoding length we can hope for is N \* H(X)
  - Reminder: Lossless encoding

## Entropy



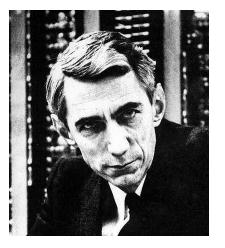
- What does it mean?
- Entropy captures different aspects of a distribution:
  - The compressibility of the data represented by random variable X
    - Follows from Shannon's theorem
  - The uncertainty of the distribution (highest entropy for uniform distribution)
    - How well can I predict a value of the random variable?
  - The information content of the random variable X
    - The number of bits used for representing a value is the information content of this value.

### **Claude Shannon**

Father of Information Theory

Envisioned the idea of communication of information with 0/1 bits

Introduced the word "bit"



The word entropy was suggested by Von Neumann

 Similarity to physics, but also "nobody really knows what entropy really is, so in any conversation you will have an advantage"

### Some information theoretic measures

 Conditional entropy H(Y|X): the uncertainty for Y given that we know X

$$H(Y|X) = -\sum_{x,y} p(x,y) \log \frac{p(x,y)}{p(x)}$$

Mutual Information I(X,Y): The reduction in the uncertainty for X (or Y) given that we know Y (or X)

I(X,Y) = H(X) - H(X|Y) = H(Y) - H(Y|X)

### Some information theoretic measures

 Cross Entropy: The cost of encoding distribution P, using the code of distribution Q

$$-\sum_{x} P(x) \log Q(x)$$

 KL Divergence KL(P||Q): The increase in encoding cost for distribution P when using the code of distribution Q

$$KL(P||Q) = -\sum_{x} P(x) \log Q(x) + \sum_{x} P(x) \log P(x)$$

- Not symmetric
- Problematic if Q not defined for all x of P.

### Some information theoretic measures

- Jensen-Shannon Divergence JS(P,Q): distance between two distributions P and Q
  - Deals with the shortcomings of KL-divergence
- If  $M = \frac{1}{2} (P+Q)$  is the mean distribution

$$JS(P,Q) = \frac{1}{2}KL(P||M) + \frac{1}{2}KL(Q||M)$$

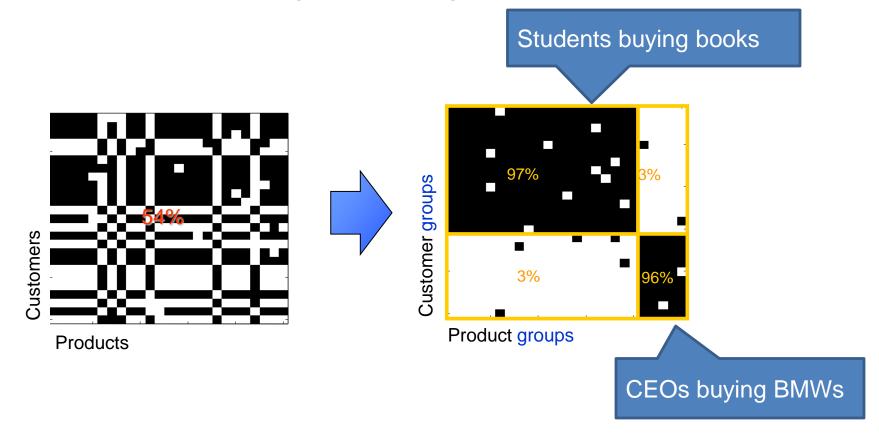
Jensen-Shannon is a metric

## USING MDL FOR CO-CLUSTERING (CROSS-ASSOCIATIONS)

Thanks to Spiros Papadimitriou.

### **Co-clustering**

 Simultaneous grouping of rows and columns of a matrix into homogeneous groups

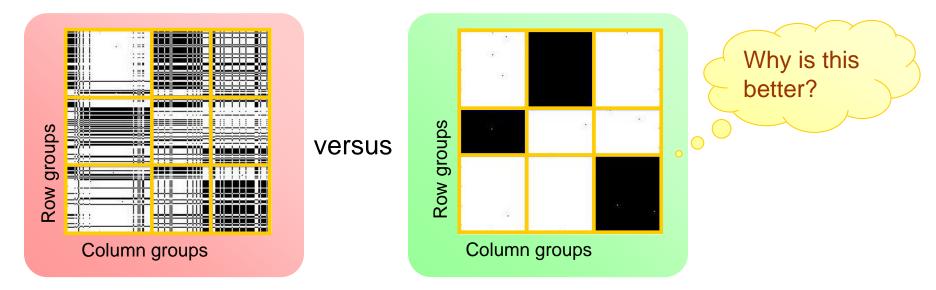


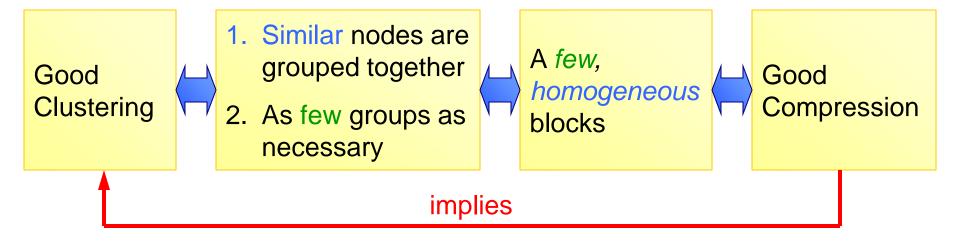
### **Co-clustering**

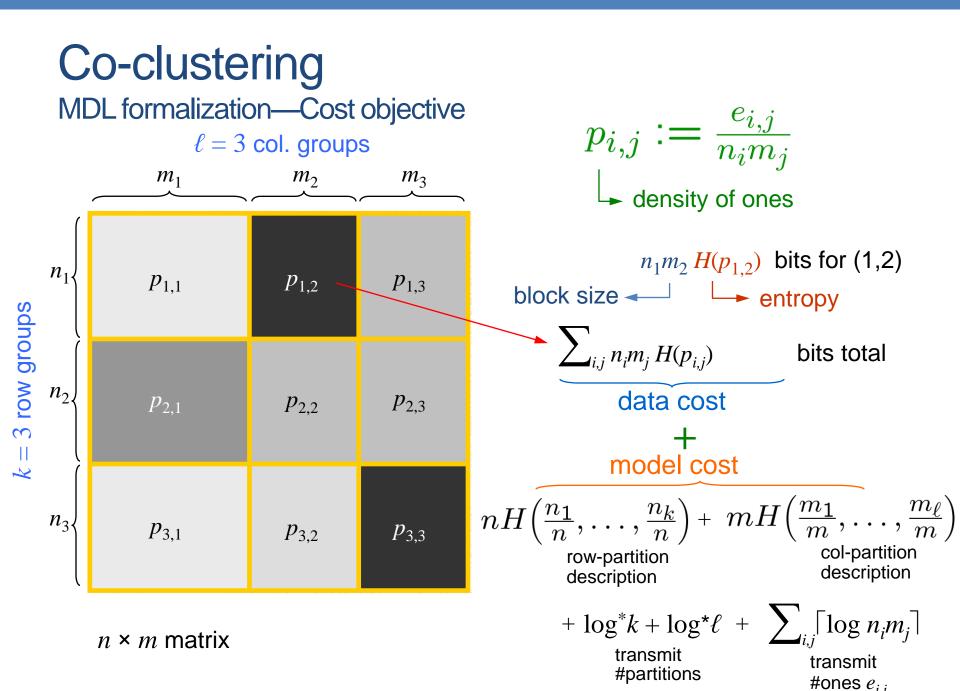
### Step 1: How to define a "good" partitioning? Intuition and formalization

• Step 2: How to find it?

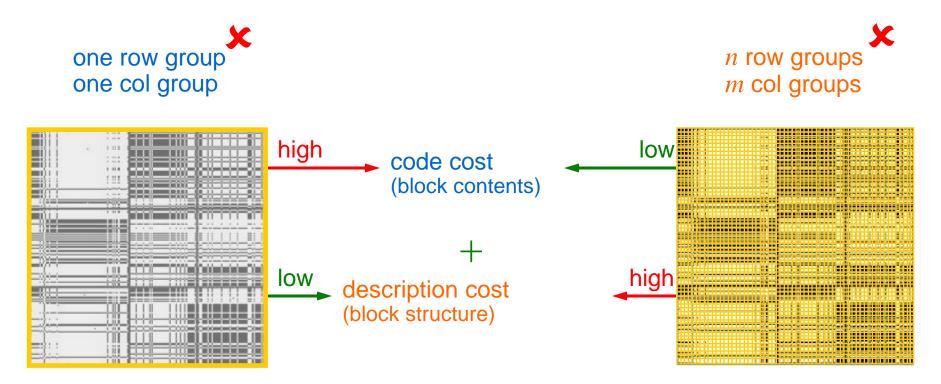
#### Co-clustering Intuition





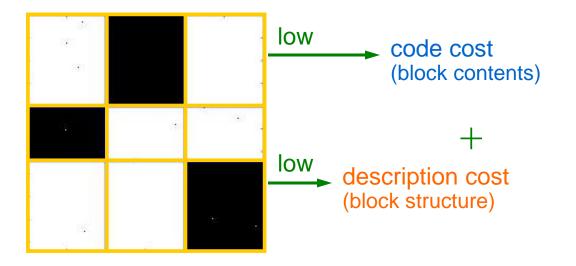


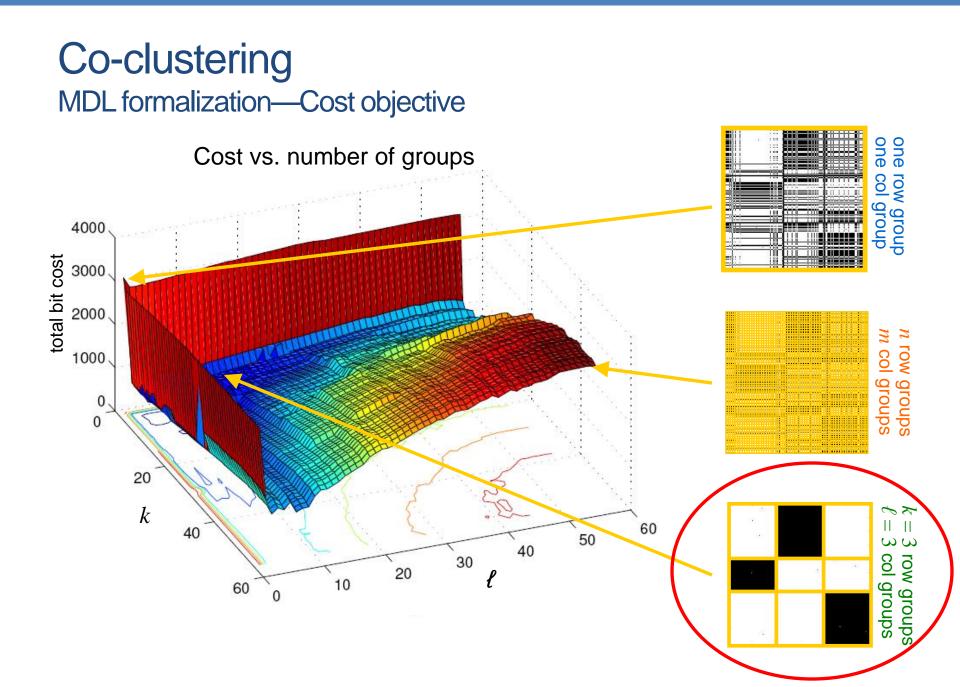
#### Co-clustering MDL formalization—Cost objective





k = 3 row groups  $\ell = 3$  col groups





### **Co-clustering**

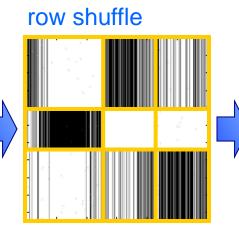
• Step 1: How to define a "good" partitioning? Intuition and formalization

• Step 2: How to find it?

Overview: assignments w/ fixed number of groups (shuffles)

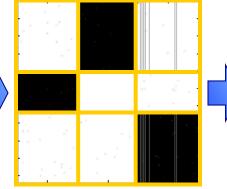
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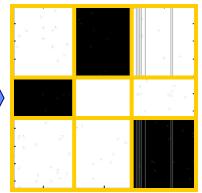


reassign all rows, holding column assignments fixed





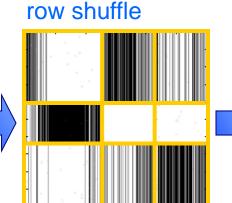
row shuffle

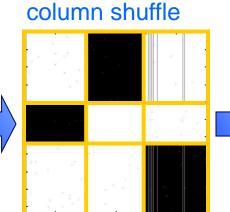


reassign all col**Noncost improvement:** holding row Discard assignments fixed

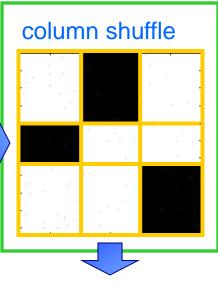
Overview: assignments w/ fixed number of groups (shuffles)



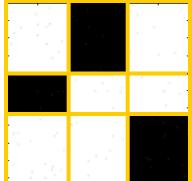




#### Final shuffle result



#### cownstrustlentfle



#### No cost improvement: Discard

#### Search for solution Shuffles

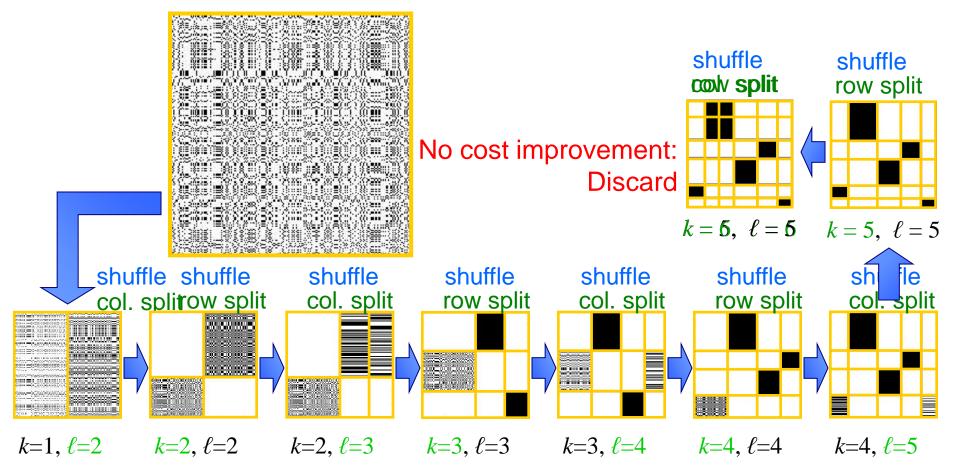
<i>P</i> <sub>1,1</sub>	<i>p</i> <sub>1,2</sub>	<i>p</i> <sub>1,3</sub>	Similarity ("KL-divergences") of row fragments to blocks of a row group Assign to second row-group	eration						
P <sub>2,1</sub> P <sub>3,1</sub>	р <sub>2,2</sub> р <sub>3,2</sub>	Р <sub>2,3</sub> Р <sub>3,3</sub>		ach part that, for all						
$-\sum_{j=1}^{\infty} \left(\nu_j \log p_{i^*,j} + (n - \nu_j) \log(1 - p_{i^*,j})\right)$ $\leq -\sum_{j=1}^{\ell} \left(\nu_j \log p_{i,j} + (n - \nu_j) \log(1 - p_{i,j})\right)$										

Overview: number of groups k and  $\ell$  (splits & shuffles)

 $k = 5, \ \ell = 5$ 

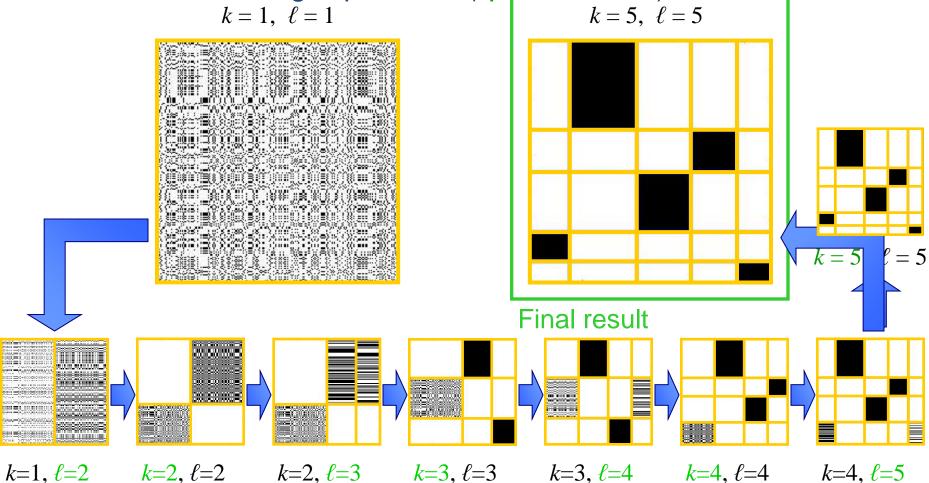
Overview: number of groups k and  $\ell$  (splits & shuffles)

 $k = 1, \ \ell = 1$ 



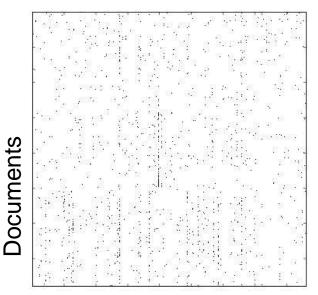
Split:Shuffle:Increase k or  $\ell$ Rearrange rows or cols

Overview: number of groups k and  $\ell$  (splits & shuffles)



Split:Shuffle:Increase k or lRearrange rows or cols

# CLASSIC



Words

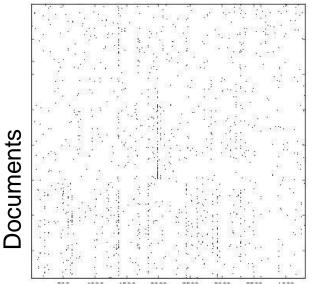
CLASSIC corpus

- 3,893 documents
- 4,303 words
- 176,347 "dots" (edges)

Combination of 3 sources:

- MEDLINE (medical)
- CISI (info. retrieval)
- CRANFIELD (aerodynamics)

#### Graph co-clustering CLASSIC



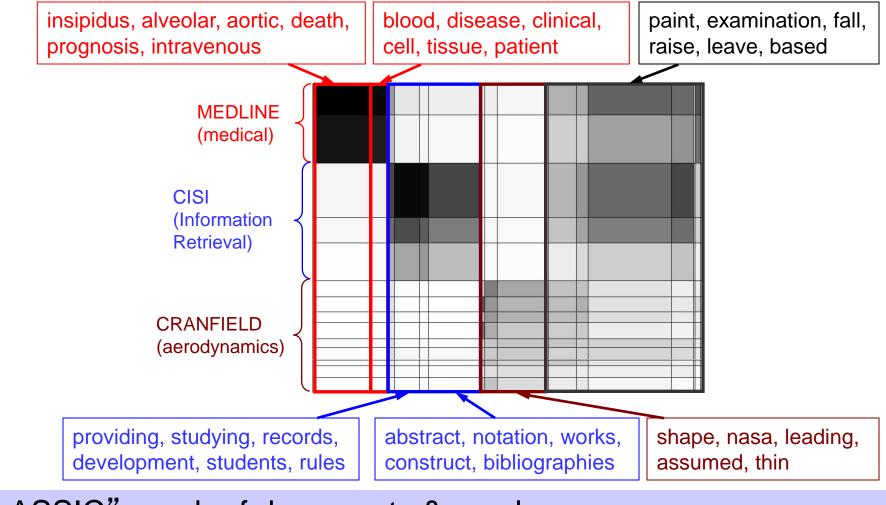
Words

-	8			-	-	_		
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# "CLASSIC" graph of documents & words: $k = 15, \ell = 19$

### Co-clustering

#### CLASSIC



"CLASSIC" graph of documents & words:  $k = 15, \ell = 19$ 

# CLASSIC

Document	Doc	ument clas	Precision			
cluster #	CRANFIELD	CISI	MEDLINE			
1	0	1	390	0.997		
2	0	0	610	1.000	}0.999	
3	2	676	9	0.984		
4	1	317	6	0.978	≻0.975	
5	3	452	16	0.960		
6	207	0	0	1.000		00
7	188	0	0	1.000		0.94-1.00
8	131	0	0	1.000		.94
9	209	0	0	1.000		0
10	107	2	0	0.982	≻0.987	
11	152	3	2	0.968	0.907	
12	74	0	0	1.000		
13	139	9	0	0.939		
14	163	0	0	1.000		
15	24	0	0	1.000		
Recall	0.996	0.990	0.968			

0.97-0.99