

# DATA MINING

## LECTURE 12

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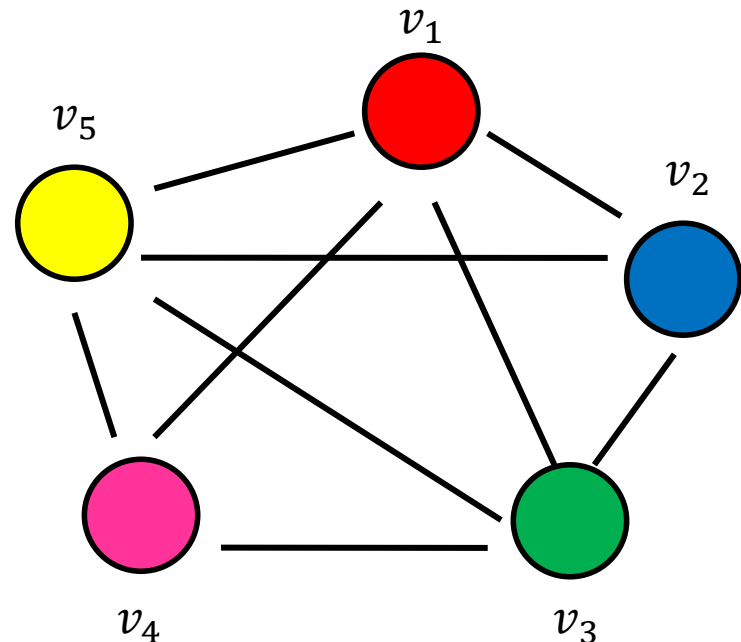
**Graphs, Node importance, Link  
Analysis Ranking, Random walks**

# RANDOM WALKS AND PAGERANK

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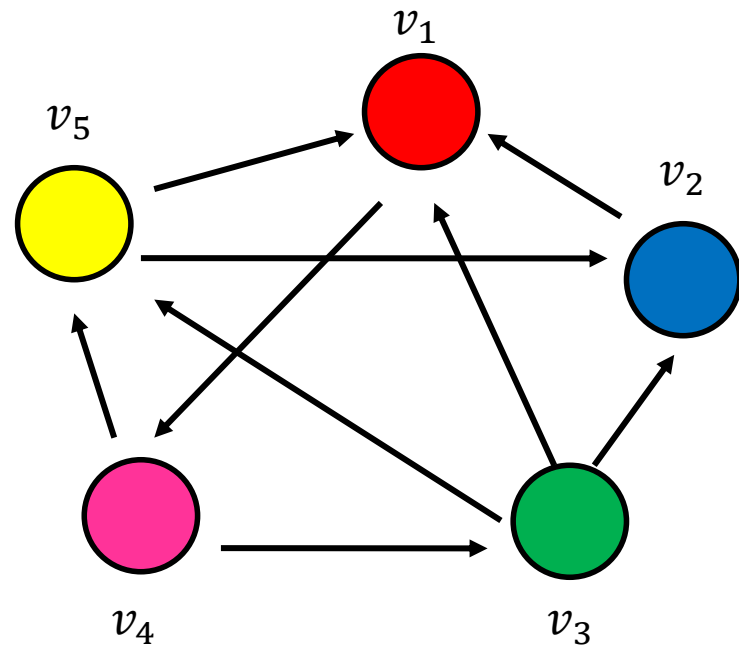
# Graphs

- A graph is a powerful abstraction for modeling entities and their pairwise relationships.
- Examples:
  - Social network
  - Collaboration graphs
  - Twitter Followers
  - Web



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# Mining the graph structure

- A graph is a combinatorial object, with a certain **structure**.
- Mining the structure of the graph reveals information about the entities in the graph
  - E.g., if in the Facebook graph I find that there are 100 people that are all linked to each other, then these people are likely to be a community
    - The **community discovery** problem
  - By measuring the number of friends in the facebook graph I can find the most important nodes
    - The **node importance** problem
- We will now focus on the node importance problem

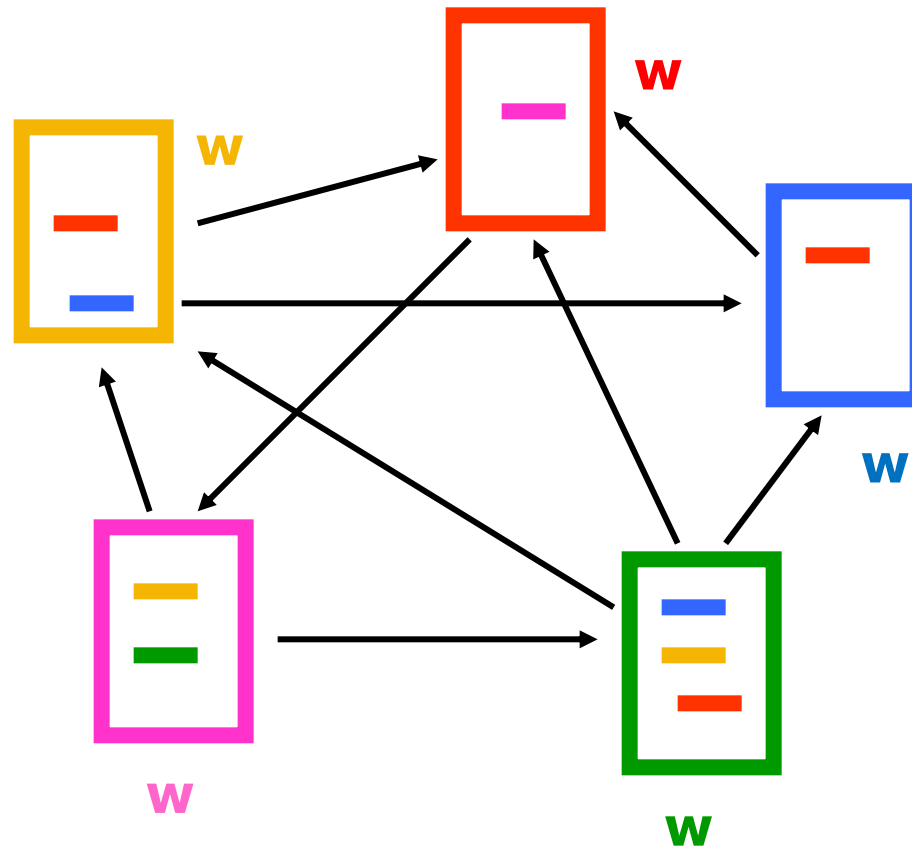
# Link Analysis

- First generation search engines
  - view documents as flat text files
  - could not cope with size, spamming, user needs
- Second generation search engines
  - Ranking becomes critical
  - shift from **relevance** to **authoritativeness**
    - **authoritativeness**: the **static** importance of the page
  - use of Web specific data: Link Analysis of the Web graph
  - a success story for the network analysis + a huge commercial success
  - it all started with two graduate students at Stanford

# Link Analysis: Intuition

- A link from page  $p$  to page  $q$  denotes endorsement
  - page  $p$  considers page  $q$  an authority on a subject
  - use the graph of recommendations
  - assign an **authority value** to every page
- The same idea applies to other graphs as well
  - Twitter graph, where user  $p$  **follows** user  $q$

# Constructing the graph

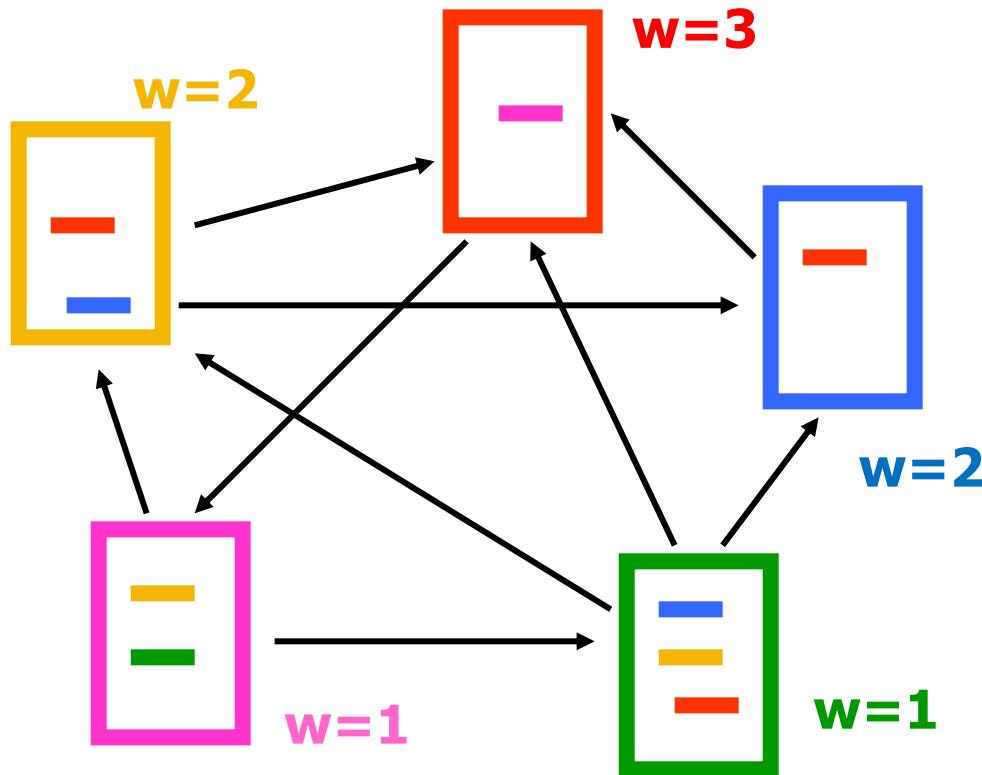


- Goal: output an **authority weight** for each node
  - Also known as **centrality**, or **importance**



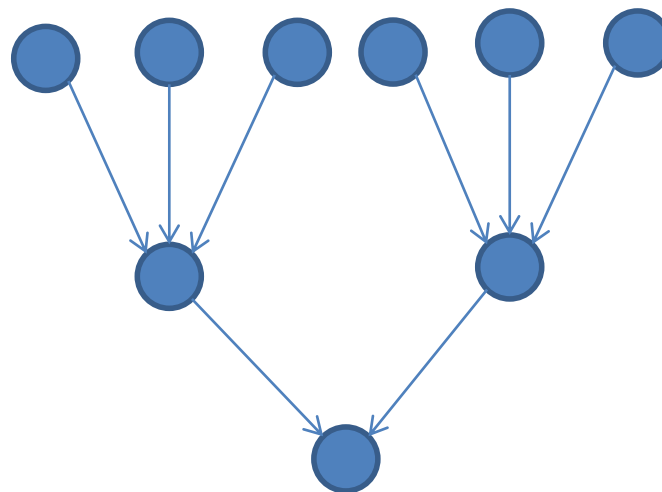
# Rank by Popularity

- Rank pages according to the number of incoming edges (**in-degree**, **degree centrality**)



- 1. Red Page**
- 2. Yellow Page**
- 3. Blue Page**
- 4. Purple Page**
- 5. Green Page**

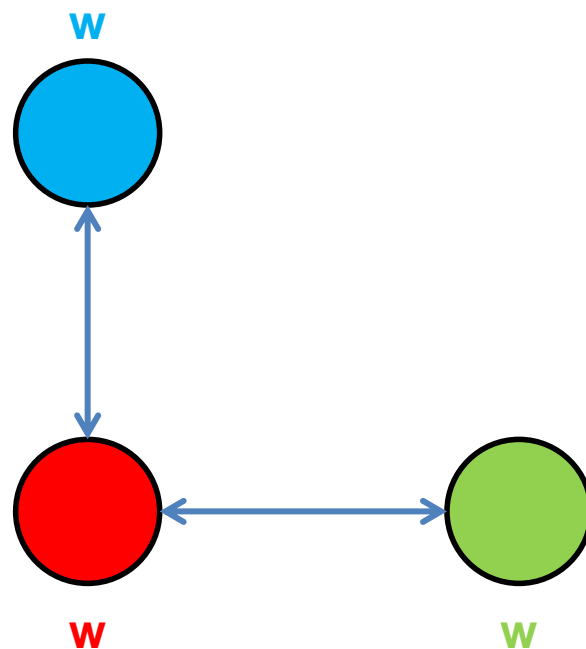
# Popularity



- It is not important only how many link to you, but how important are the people that link to you.
- **Good** authorities are pointed by **good** authorities
  - Recursive definition of importance

# PageRank

- Assume that we have a unity of authority to distribute to all nodes.
- Each node distributes the authority value they have to all their neighbors
- The authority value of each node is the sum of the fractions it collects from its neighbors.
- Solving the system of equations we get the authority values for the nodes
  - $w = \frac{1}{2}$  ,  $w = \frac{1}{4}$  ,  $w = \frac{1}{4}$



$$w + w + w = 1$$

$$w = w + w$$

$$w = \frac{1}{2} w$$

$$w = \frac{1}{2} w$$

# A more complex example

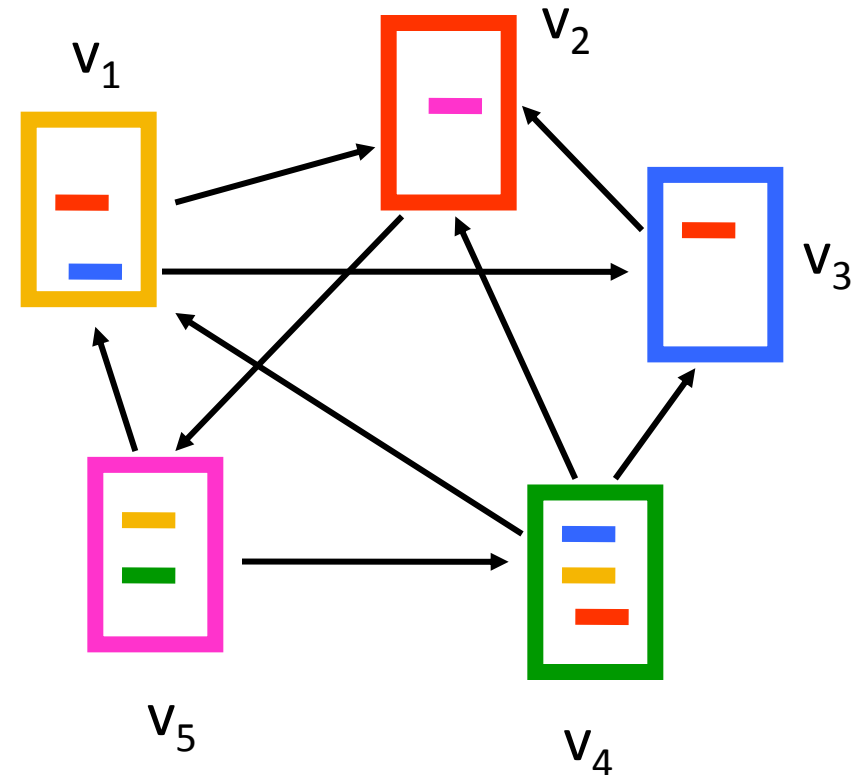
$$w_1 = 1/3 w_4 + 1/2 w_5$$

$$w_2 = 1/2 w_1 + w_3 + 1/3 w_4$$

$$w_3 = 1/2 w_1 + 1/3 w_4$$

$$w_4 = 1/2 w_5$$

$$w_5 = w_2$$



$$PR(p) = \sum_{q \rightarrow p} \frac{PR(q)}{|Out(q)|}$$

# Random walks on graphs

- The equations above describe a **step** of a **random walk** on the graph
  - Random walk: start from some node uniformly at random and then from each node pick a random link to follow.
  - Question: what is the probability of being at a specific node?
    - $p_i$ : probability of being at node  $i$  at this step
    - $p'_i$ : probability of being at node  $i$  in the next step

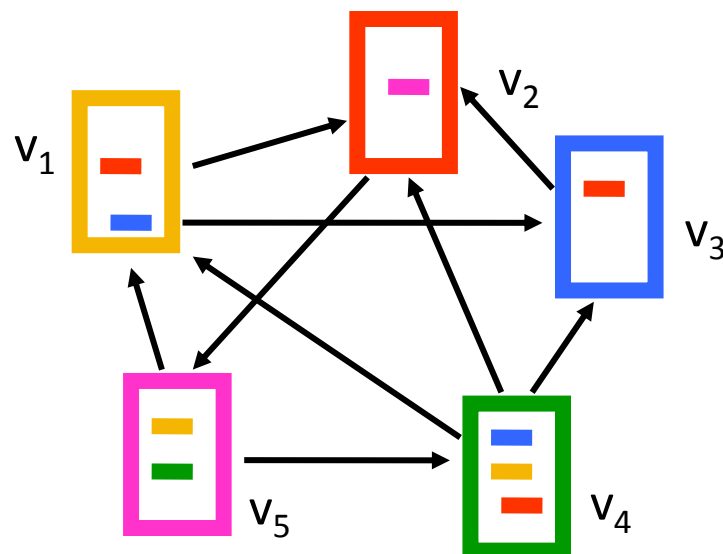
$$p'_1 = 1/3 p_4 + 1/2 p_5$$

$$p'_2 = 1/2 p_1 + p_3 + 1/3 p_4$$

$$p'_3 = 1/2 p_1 + 1/3 p_4$$

$$p'_4 = 1/2 p_5$$

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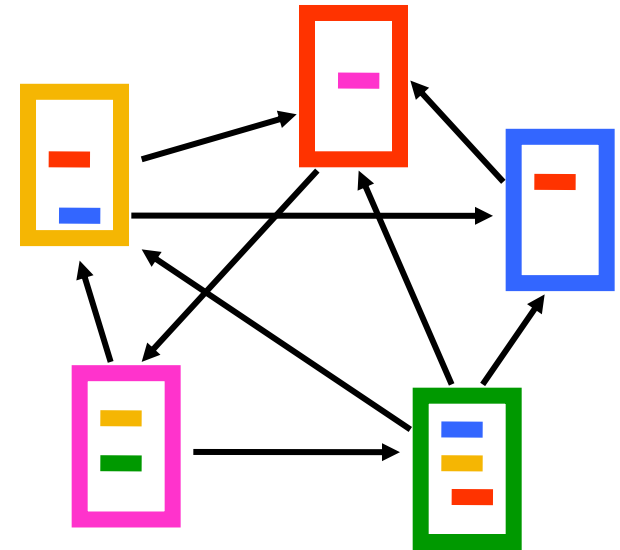


- After many steps the probabilities **converge** to the **stationary distribution** of the random walk.

# PageRank algorithm [BP98]

- **Good** authorities should be pointed by **good** authorities
  - The value of a page is the value of the people that link to you
- How do we implement that?
  - Each page has a value.
  - Proceed in iterations,
    - in each iteration every page **distributes** the value to the neighbors
  - Continue until there is convergence.

$$PR(p) = \sum_{q \rightarrow p} \frac{PR(q)}{|Out(q)|}$$



1. **Red Page**
2. **Purple Page**
3. **Yellow Page**
4. **Blue Page**
5. **Green Page**

# Markov chains

- A Markov chain describes a discrete time **stochastic process** over a set of states

$$S = \{s_1, s_2, \dots, s_n\}$$

according to a **transition probability matrix**

$$P = \{P_{ij}\}$$

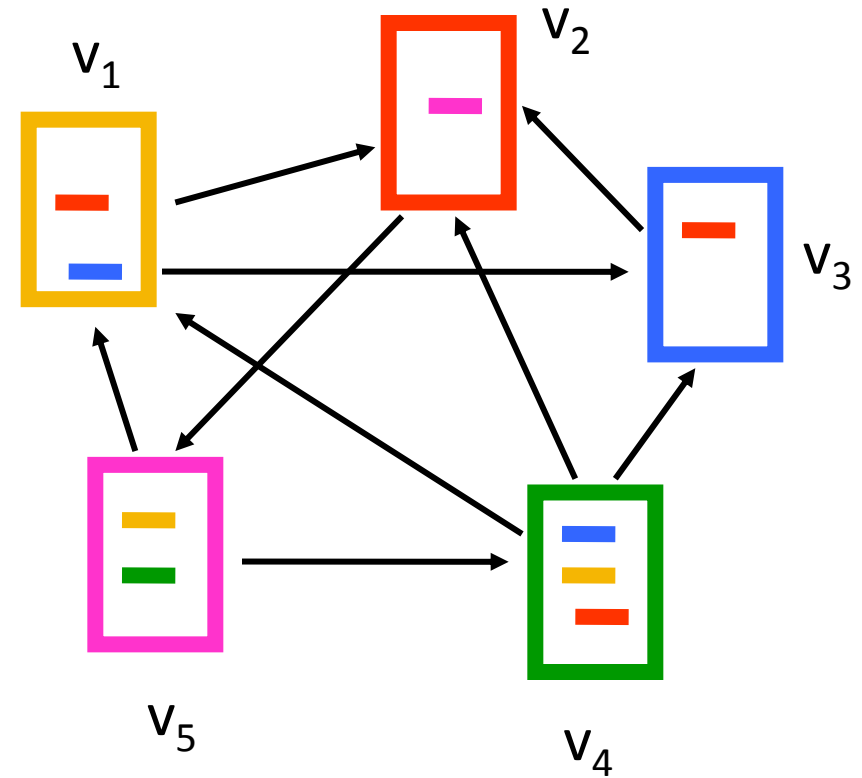
- $P_{ij}$  = probability of moving to state  $j$  when at state  $i$ 
  - $\sum_j P_{ij} = 1$  (**stochastic matrix**)
- **Memorylessness property**: The next state of the chain depends only at the current state and not on the past of the process (first order MC)
  - higher order MCs are also possible

# Random walks

- Random walks on graphs correspond to Markov Chains
  - The set of states  $S$  is the set of nodes of the graph  $G$
  - The **transition probability matrix** is the probability that we follow an edge from one node to another



# An example

$$A = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$
$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 0 & 1/2 \end{bmatrix}$$


# State probability vector

- The vector  $q^t = (q_1^t, q_2^t, \dots, q_n^t)$  that stores the probability of being at state  $i$  at time  $t$ 
  - $q_i^0$  = the probability of starting from state  $i$

$$q^t = q^{t-1} P$$

# An example

$$q^t = q^{t-1} P$$

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

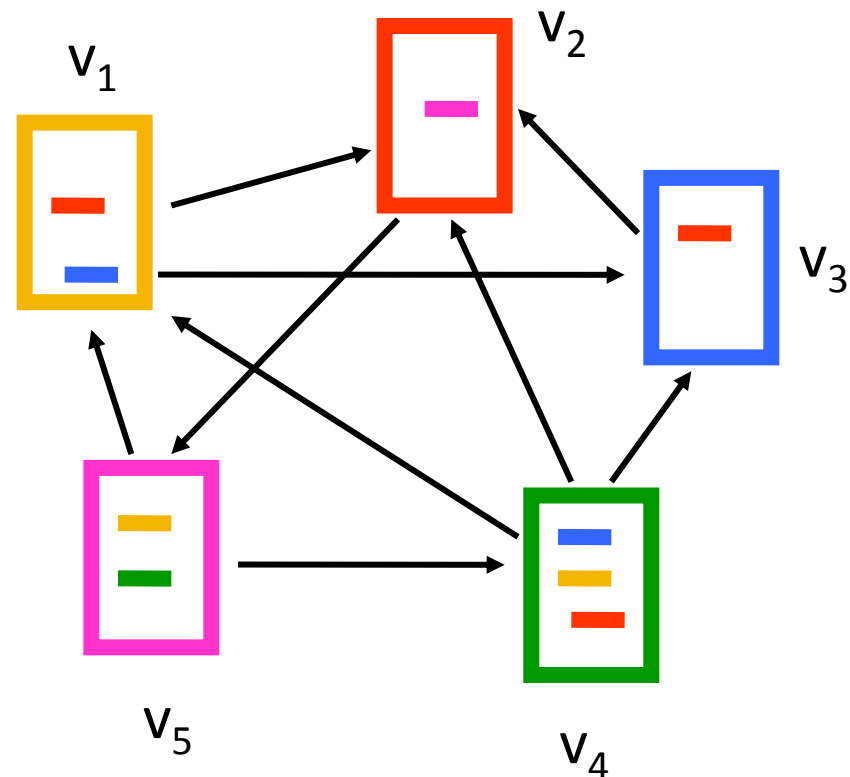
$$q^{t+1}_1 = 1/3 q^t_4 + 1/2 q^t_5$$

$$q^{t+1}_2 = 1/2 q^t_1 + q^t_3 + 1/3 q^t_4$$

$$q^{t+1}_3 = 1/2 q^t_1 + 1/3 q^t_4$$

$$q^{t+1}_4 = 1/2 q^t_5$$

$$q^{t+1}_5 = q^t_2$$



Same equations as before!

# Stationary distribution

- A stationary distribution for a MC with transition matrix  $P$ , is a probability distribution  $\pi$ , such that  $\pi = \pi P$
- A MC has a unique stationary distribution if
  - it is **irreducible**
    - the underlying graph is **strongly connected**
  - it is **aperiodic**
    - for random walks, the underlying graph is **not** bipartite
- The probability  $\pi_i$  is the fraction of times that we visited state  $i$  as  $t \rightarrow \infty$
- The stationary distribution is an eigenvector of matrix  $P$ 
  - the principal left eigenvector of  $P$  – stochastic matrices have maximum eigenvalue 1

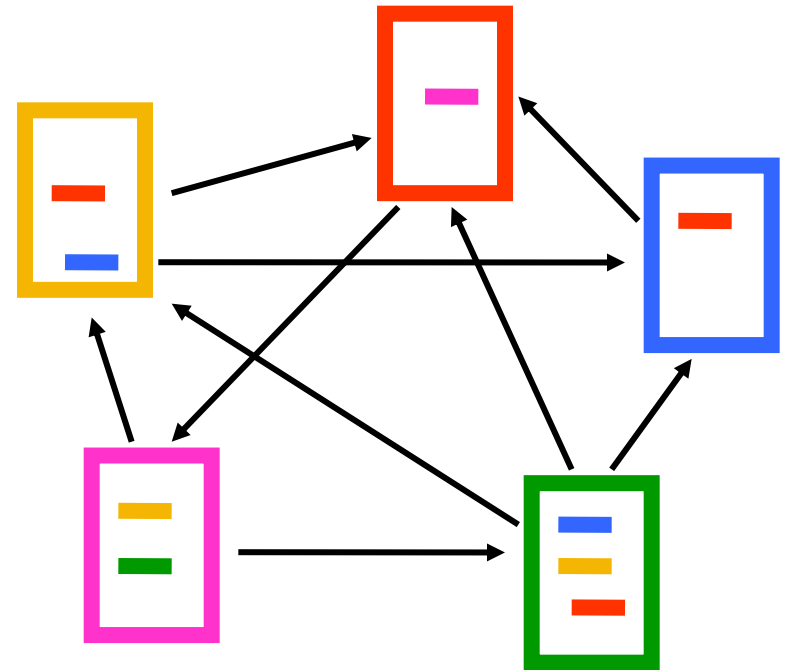
# Computing the stationary distribution

- The Power Method
  - Initialize to some distribution  $q^0$
  - Iteratively compute  $q^t = q^{t-1}P$
  - After enough iterations  $q^t \approx \pi$
  - Power method because it computes  $q^t = q^0 P^t$
- Why does it converge?
  - follows from the fact that any vector can be written as a linear combination of the eigenvectors
    - $q^0 = v_1 + c_2 v_2 + \dots + c_n v_n$
- Rate of convergence
  - determined by  $\lambda_2^t$

# The PageRank random walk

- Vanilla random walk
  - make the adjacency matrix stochastic and run a random walk

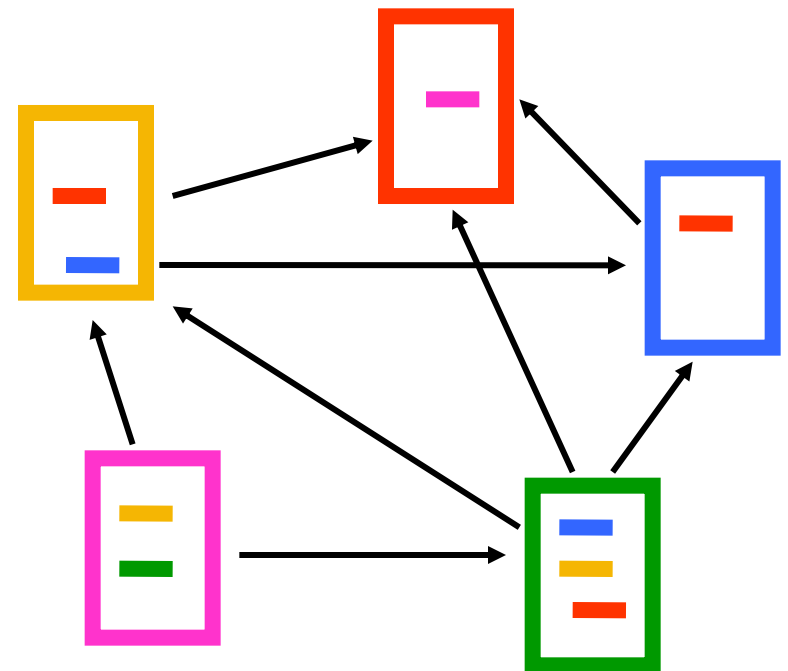
$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$



# The PageRank random walk

- What about **sink** nodes?
  - what happens when the random walk moves to a node without any outgoing links?

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

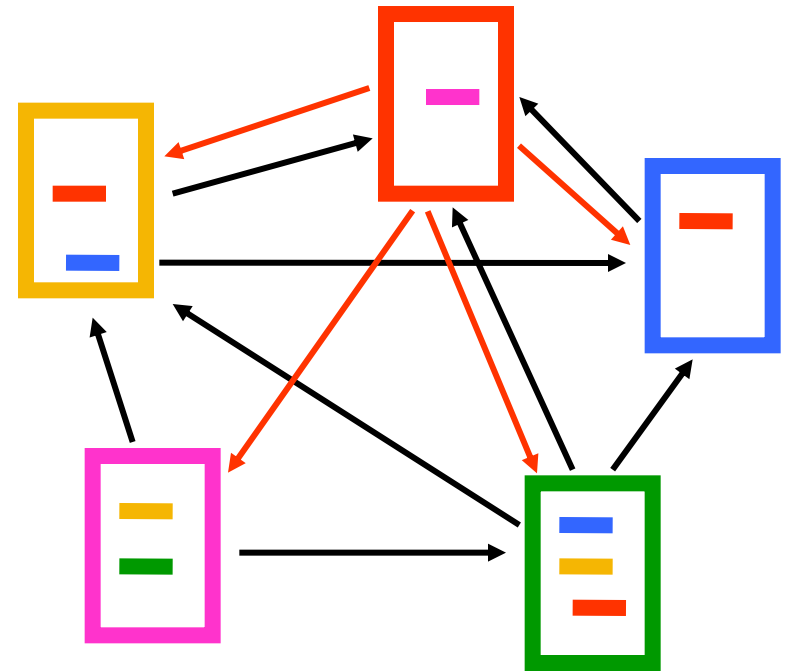


# The PageRank random walk

- Replace these row vectors with a vector  $\mathbf{v}$ 
  - typically, the uniform vector

$$P' = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

$$P' = P + d\mathbf{v}^T \quad d = \begin{cases} 1 & \text{if } i \text{ is sink} \\ 0 & \text{otherwise} \end{cases}$$





# The PageRank random walk

- How do we guarantee irreducibility?
- How do we guarantee not getting stuck in loops?
  - add a random jump to vector  $\mathbf{v}$  with prob  $\alpha$ 
    - typically, to a uniform vector

$$P'' = \alpha \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 0 & 1/2 \end{bmatrix} + (1-\alpha) \begin{bmatrix} 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \end{bmatrix}$$

$P'' = \alpha P' + (1-\alpha)uv^T$ , where  $u$  is the vector of all 1s

Random walk with restarts

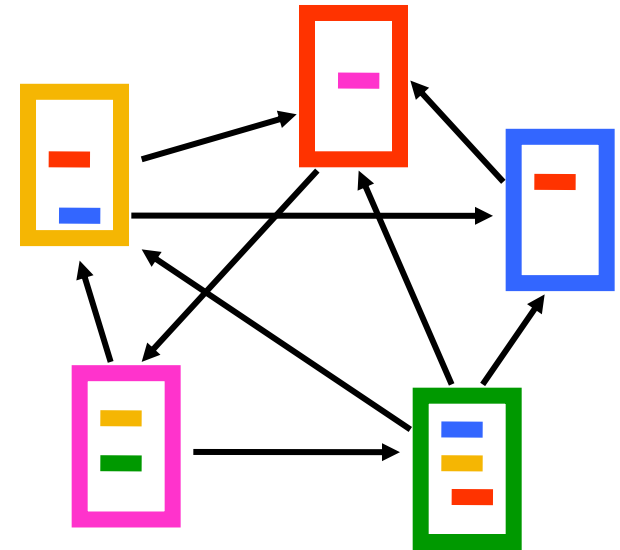
# PageRank algorithm [BP98]

- The Random Surfer model
  - pick a page at random
  - with probability  $1 - \alpha$  jump to a random page
  - with probability  $\alpha$  follow a random outgoing link

- Rank according to the stationary distribution

$$PR(p) = \alpha \sum_{q \rightarrow p} \frac{PR(q)}{|Out(q)|} + (1 - \alpha) \frac{1}{n}$$

$\alpha = 0.85$  in most cases



1. Red Page
2. Purple Page
3. Yellow Page
4. Blue Page
5. Green Page

# The stationary distribution

- What is the meaning of the stationary distribution  $\pi$  of a random walk?
- $\pi(i)$ : the probability of being at node  $i$  after very large (infinite) number of steps
- $\pi = p_0 P^\infty$ , where  $P$  is the transition matrix,  $p_0$  the original vector
  - $P(i, j)$ : probability of going from  $i$  to  $j$  in one step
  - $P^2(i, j)$ : probability of going from  $i$  to  $j$  in two steps (probability of all paths of length 2)
  - $P^\infty(i, j) = \pi(j)$ : probability of going from  $i$  to  $j$  in infinite steps – starting point does not matter.

# Stationary distribution with random jump

- If  $v$  is the jump vector

$$p_0 = v$$

$$p_1 = \alpha p_0 P + (1 - \alpha)v = \alpha v P + (1 - \alpha)v$$

$$p_2 = \alpha p_1 P + (1 - \alpha)v = \alpha^2 v P^2 + (1 - \alpha)v \alpha P + (1 - \alpha)v$$

$\vdots$

$$p^\infty = (1 - \alpha)v + (1 - \alpha)v \alpha P + (1 - \alpha)v \alpha^2 P^2 + \dots$$

$$= (1 - \alpha)(I - \alpha P)^{-1}$$

- With the random jump the shorter paths are more important, since the weight decreases exponentially
  - makes sense when thought of as a restart
- If  $v$  is not uniform, we can bias the random walk towards the pages that are close to  $v$ 
  - **Personalized** and **Topic Specific** Pagerank.

# Effects of random jump

- Guarantees irreducibility
- Motivated by the concept of random surfer
- Offers additional flexibility
  - personalization
  - anti-spam
- Controls the rate of convergence
  - the second eigenvalue of matrix  $P''$  is  $\alpha$

# Random walks on undirected graphs

- For undirected graphs, the stationary distribution is proportional to the degrees of the nodes
  - Thus in this case a random walk is the same as degree popularity
- This is not longer true if we do random jumps
  - Now the short paths play a greater role, and the previous distribution does not hold.

# A PageRank algorithm

- Performing vanilla power method is now too expensive – the matrix is not sparse

$$q^0 = v$$

$$t = 1$$

repeat

$$q^t = (P'')^T q^{t-1}$$

$$\delta = \|q^t - q^{t-1}\|$$

$$t = t + 1$$

until  $\delta < \epsilon$

Efficient computation of  $y = (P'')^T x$

$$y = \alpha P^T x$$

$$\beta = \|x\|_1 - \|y\|_1$$

$$y = y + \beta v$$

$P$  = normalized adjacency matrix

$P' = P + dv^T$ , where  $d_i$  is 1 if  $i$  is sink and 0 o.w.

$P'' = \alpha P' + (1-\alpha)uv^T$ , where  $u$  is the vector of all 1s

# Pagerank history

- Huge advantage for Google in the early days
  - It gave a way to get an idea for the value of a page, which was useful in many different ways
    - Put an order to the web.
  - After a while it became clear that the anchor text was probably more important for ranking
  - Also, link spam became a new (dark) art
- Flood of research
  - Numerical analysis got rejuvenated
  - Huge number of variations
  - Efficiency became a great issue.
  - Huge number of applications in different fields
    - Random walk is often referred to as PageRank.



# THE HITS ALGORITHM

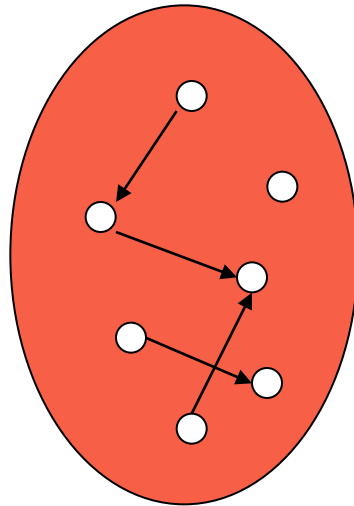
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# The HITS algorithm

- Another algorithm proposed around the same time as Pagerank for using the hyperlinks to rank pages
  - Kleinberg: then an intern at IBM Almaden
  - IBM never made anything out of it

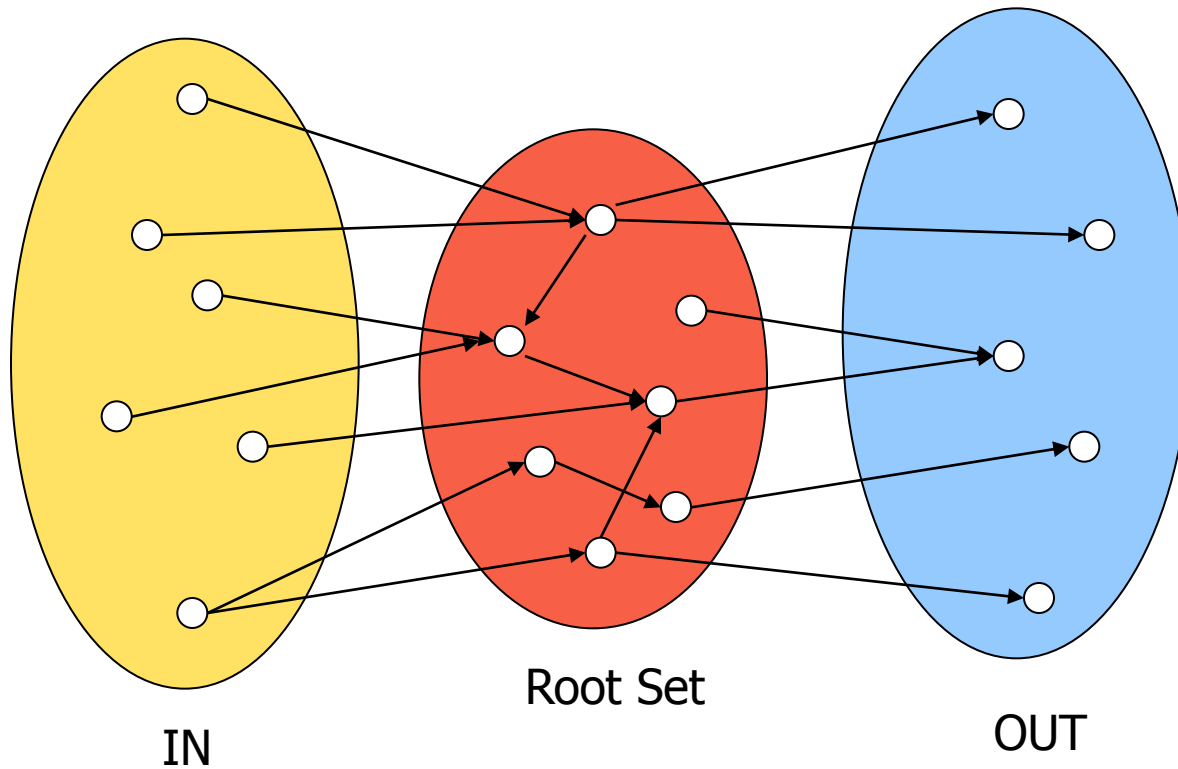
# Query dependent input

Root set obtained from a text-only search engine

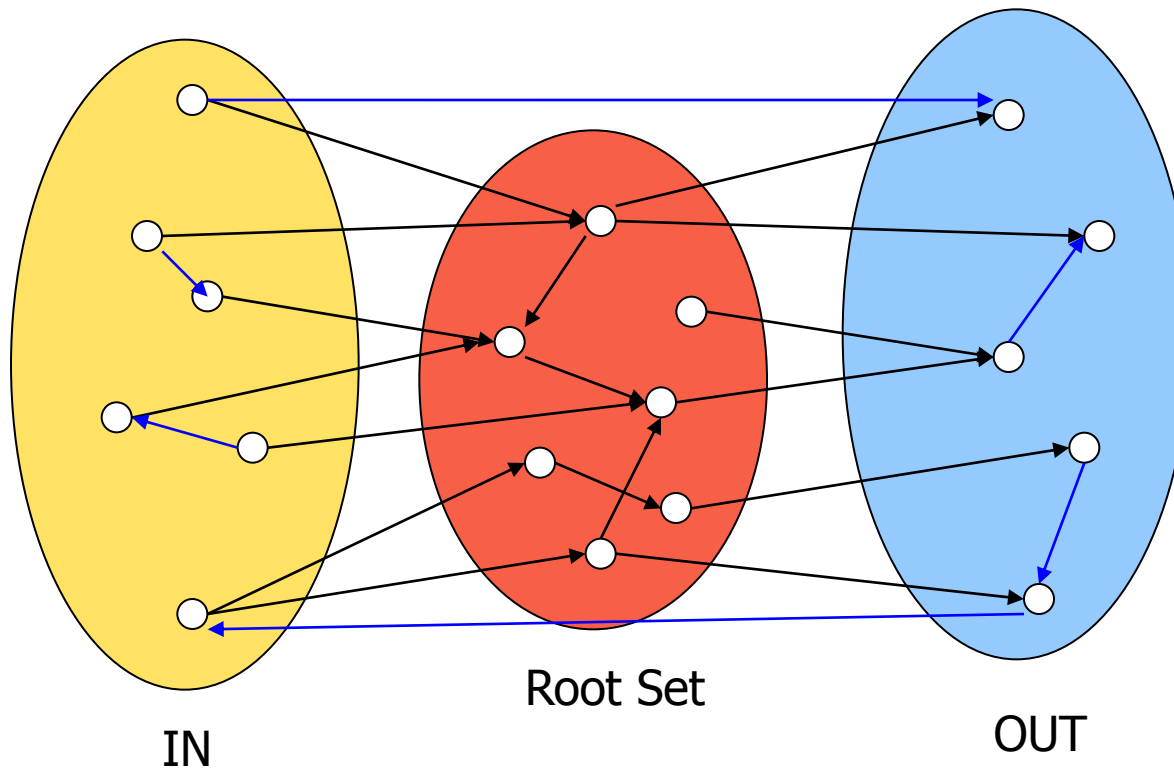


Root Set

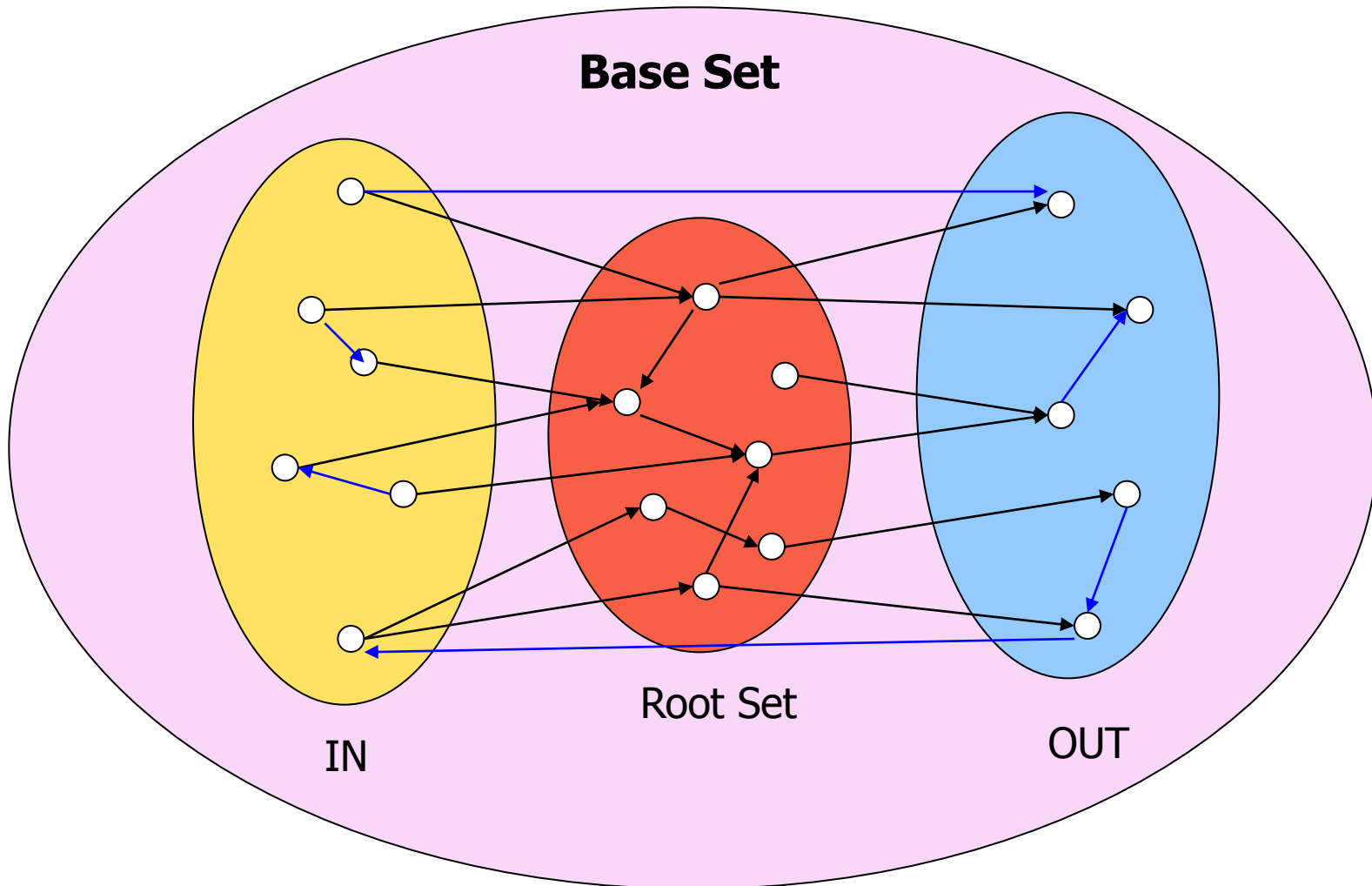
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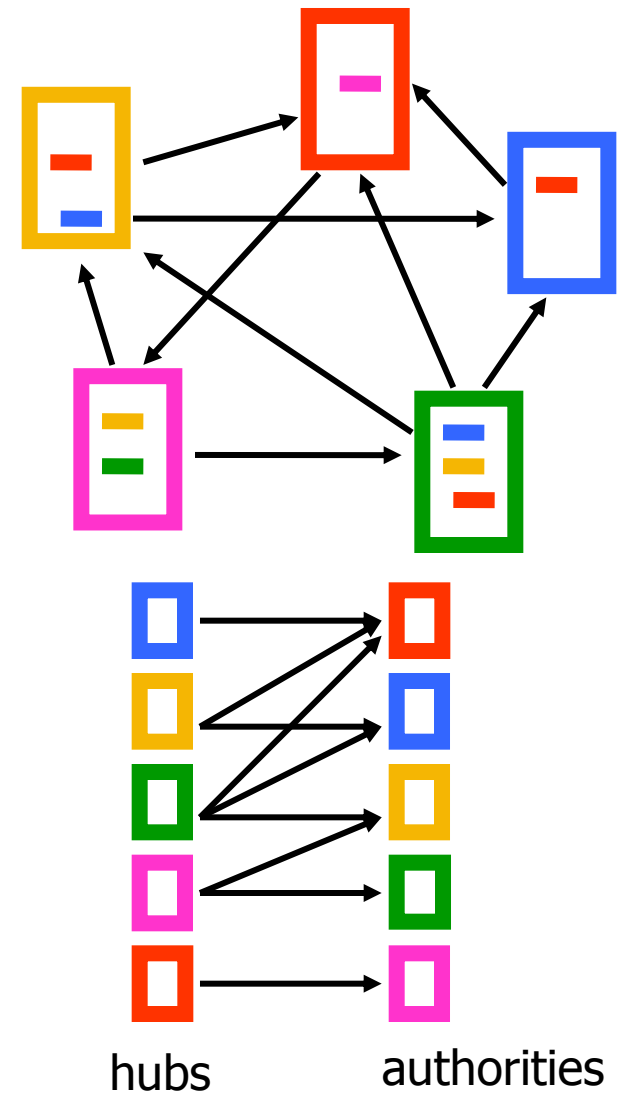


# Query dependent input



# Hubs and Authorities [K98]

- Authority is not necessarily transferred directly between authorities
- Pages have double identity
  - **hub** identity
  - **authority** identity
- **Good** hubs point to **good** authorities
- **Good** authorities are pointed by **good** hubs



# Hubs and Authorities

- Two kind of weights:
  - Hub weight
  - Authority weight
- The hub weight is the sum of the authority weights of the authorities pointed to by the hub
- The authority weight is the sum of the hub weights that point to this authority.



# HITS Algorithm

- Initialize all weights to 1.
- Repeat until convergence
  - *O* operation : hubs collect the weight of the authorities

$$h_i = \sum_{j:i \rightarrow j} a_j$$

- *I* operation: authorities collect the weight of the hubs

$$a_i = \sum_{j:j \rightarrow i} h_j$$

- Normalize weights under some norm

# HITS and eigenvectors

- The HITS algorithm is a power-method eigenvector computation
  - in vector terms  $\mathbf{a}^t = A^T \mathbf{h}^{t-1}$  and  $\mathbf{h}^t = A \mathbf{a}^{t-1}$
  - so  $\mathbf{a} = A^T A \mathbf{a}^{t-1}$  and  $\mathbf{h}^t = A A^T \mathbf{h}^{t-1}$
  - The authority weight vector  $\mathbf{a}$  is the eigenvector of  $A^T A$  and the hub weight vector  $\mathbf{h}$  is the eigenvector of  $A A^T$
  - Why do we need normalization?
- The vectors  $\mathbf{a}$  and  $\mathbf{h}$  are **singular vectors** of the matrix  $A$

# Singular Value Decomposition

$$A = U \Sigma V^T = \begin{bmatrix} \vec{u}_1 & \vec{u}_2 & \cdots & \vec{u}_r \end{bmatrix} \begin{bmatrix} \sigma_1 & & & \\ & \sigma_2 & & \\ & & \ddots & \\ & & & \sigma_r \end{bmatrix} \begin{bmatrix} \vec{v}_1 \\ \vec{v}_2 \\ \vdots \\ \vec{v}_r \end{bmatrix}$$

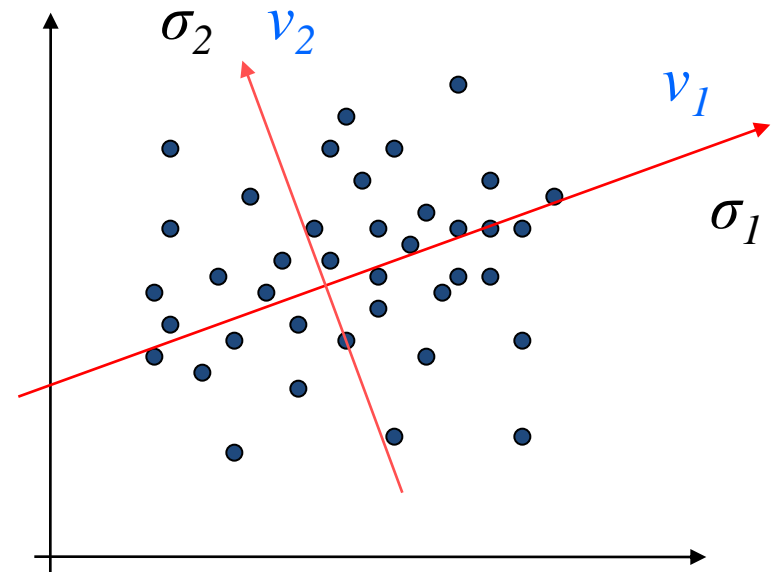
$[n \times r] \quad [r \times r] \quad [r \times n]$

- $r$ : rank of matrix  $A$
- $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_r$ : singular values (square roots of eig-vals  $AA^T, A^T A$ )
- $\vec{u}_1, \vec{u}_2, \dots, \vec{u}_r$ : left singular vectors (eig-vectors of  $AA^T$ )
- $\vec{v}_1, \vec{v}_2, \dots, \vec{v}_r$ : right singular vectors (eig-vectors of  $A^T A$ )

$$A = \sigma_1 \vec{u}_1 \vec{v}_1^T + \sigma_2 \vec{u}_2 \vec{v}_2^T + \cdots + \sigma_r \vec{u}_r \vec{v}_r^T$$

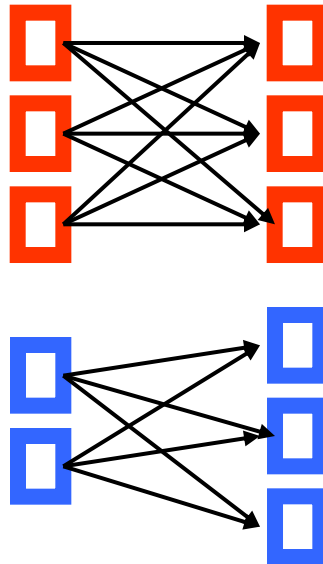
# Singular Value Decomposition

- **Linear trend  $\mathbf{v}$**  in matrix  $A$ :
    - the tendency of the row vectors of  $A$  to align with vector  $\mathbf{v}$
    - strength of the linear trend:  $A\mathbf{v}$
  - SVD discovers the linear trends in the data
  - $\mathbf{u}_i, \mathbf{v}_i$ : the  $i$ -th strongest linear trends
  - $\sigma_i$ : the strength of the  $i$ -th strongest linear trend
- 
- HITS discovers the **strongest linear trend** in the authority space



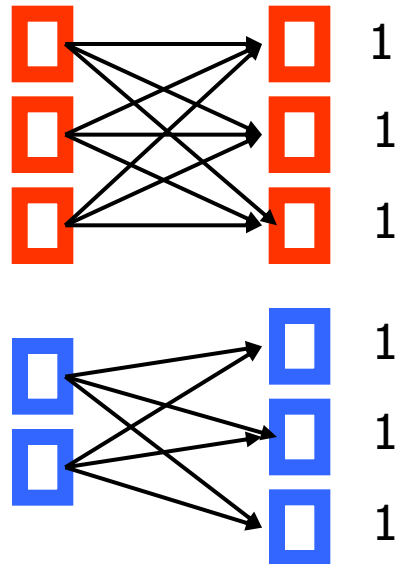
# HITS and the TKC effect

- The HITS algorithm favors the most **dense community** of hubs and authorities
  - Tightly Knit Community (TKC) effect



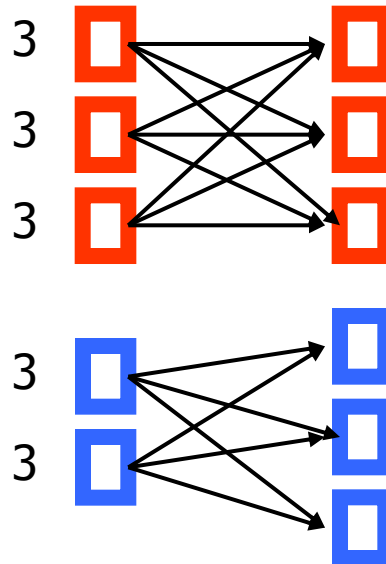
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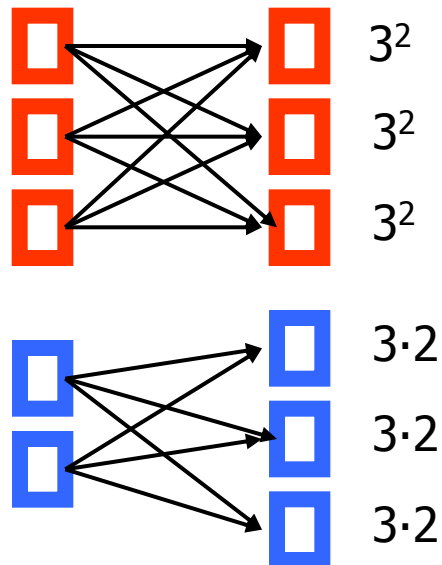
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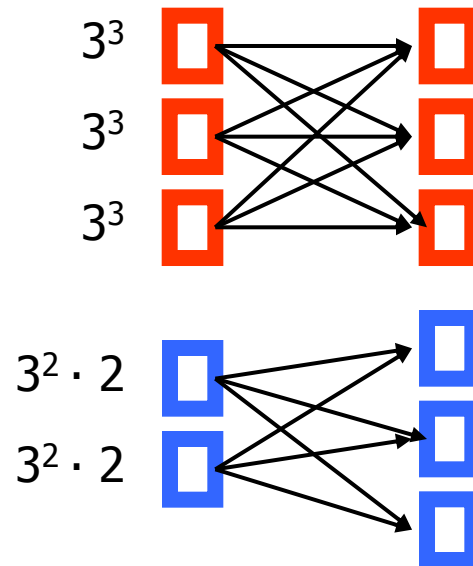
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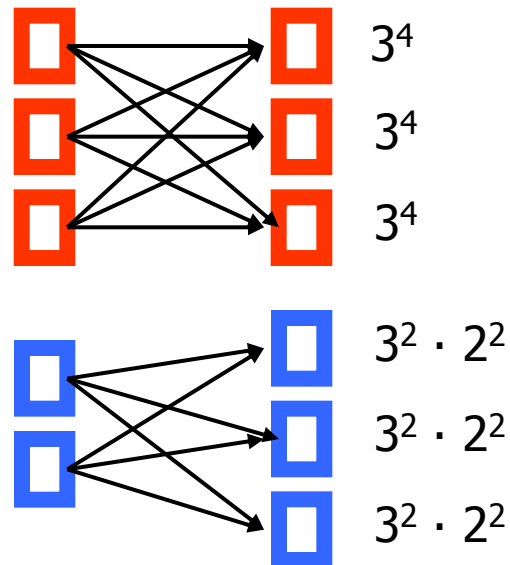
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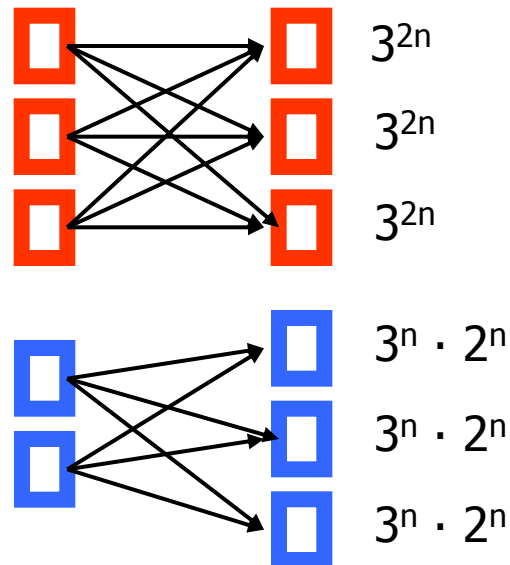
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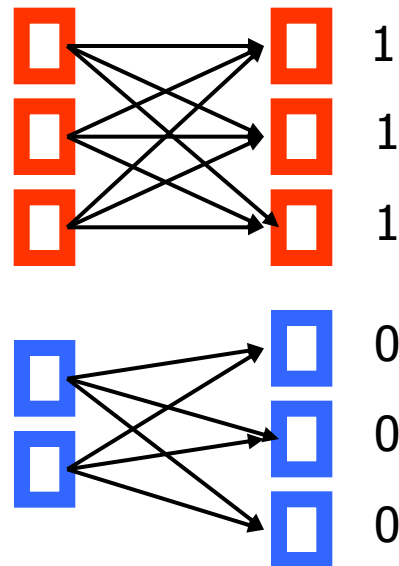
weight of node  $p$  is proportional to the number of  $(BF)^n$  paths that leave node  $p$



after  $n$  iterations

# HITS and the TKC effect

- The HITS algorithm favors the most **dense community** of hubs and authorities
  - Tightly Knit Community (TKC) effect



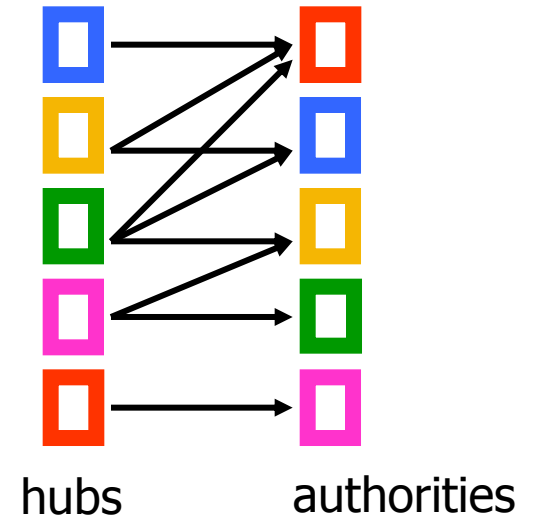
after normalization  
with the max  
element as  $n \rightarrow \infty$

# OTHER ALGORITHMS

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# The SALSA algorithm [LM00]

- Perform a random walk alternating between hubs and authorities
- What does this random walk converge to?
- The graph is essentially undirected, so it will be proportional to the degree.



# Social network analysis

- Evaluate the **centrality** of individuals in social networks

- **degree centrality**

- the (weighted) degree of a node

- **distance centrality**

- the average (weighted) distance of a node to the rest in the graph

$$D_c(v) = \frac{1}{\sum_{u \neq v} d(v, u)}$$

- **betweenness centrality**

- the average number of (weighted) shortest paths that use node  $v$

$$B_c(v) = \sum_{s \neq v \neq t} \frac{\sigma_{st}(v)}{\sigma_{st}}$$

# Counting paths – Katz 53

- The importance of a node is measured by the weighted sum of paths that lead to this node
- $A^m[i,j]$  = number of paths of length  $m$  from  $i$  to  $j$
- Compute

$$P = bA + b^2A^2 + \dots + b^m A^m + \dots = (I - bA)^{-1} - I$$

- converges when  $b < \lambda_1(A)$
- Rank nodes according to the column sums of the matrix  $P$



# Bibliometrics

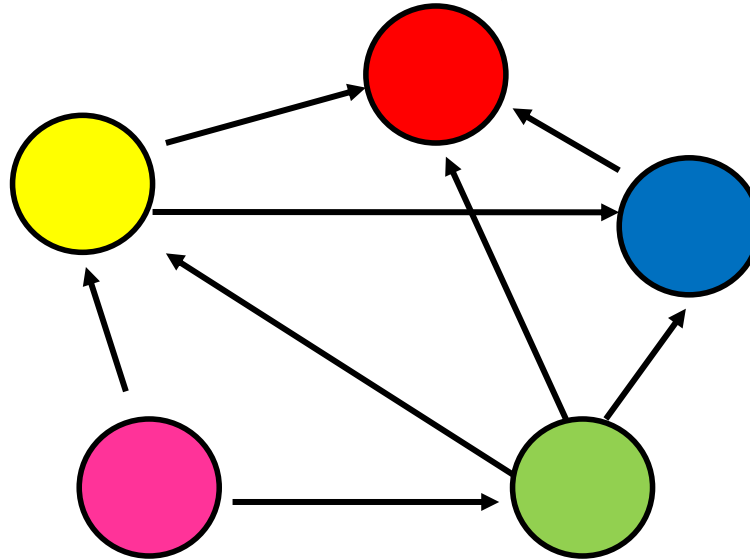
- Impact factor (E. Garfield 72)
  - counts the number of citations received for papers of the journal in the previous two years
- Pinsky-Narin 76
  - perform a random walk on the set of journals
  - $P_{ij}$  = the fraction of citations from journal  $i$  that are directed to journal  $j$

# ABSORBING RANDOM WALKS

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# Random walk with absorbing nodes

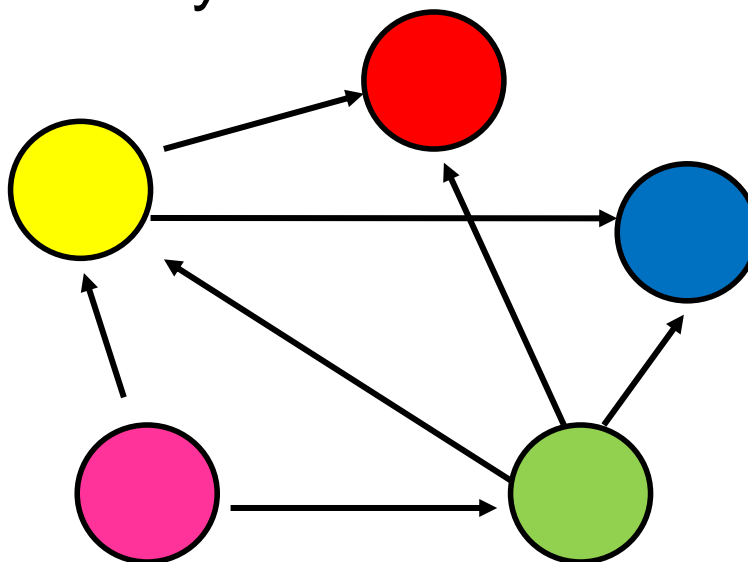
- What happens if we do a random walk on this graph? What is the stationary distribution?



- All the probability mass on the red sink node:
  - The red node is an **absorbing node**

# Random walk with absorbing nodes

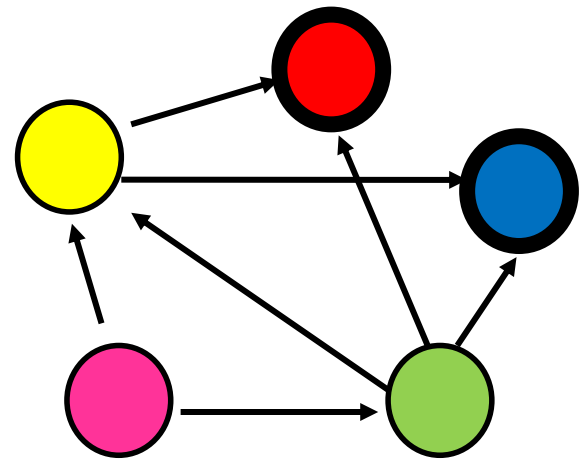
- What happens if we do a random walk on this graph? What is the stationary distribution?



- There are two absorbing nodes: the red and the blue.
- The probability mass will be divided between the two

# Absorption probability

- If there are more than one absorbing nodes in the graph a random walk that starts from a non-absorbing node will be absorbed in one of them with some probability
  - The probability of absorption gives an estimate of how close the node is to red or blue



- Why care?
  - Red and Blue may be different categories

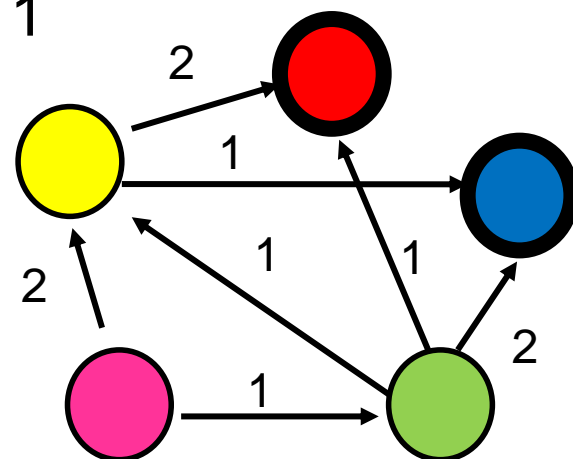
# Absorption probability

- Computing the probability of being absorbed is very easy
  - Take the (weighted) average of the absorption probabilities of your neighbors
    - if one of the neighbors is the absorbing node, it has probability 1
  - Repeat until convergence
  - Initially only the absorbing have prob 1

$$P(\text{Red}|\text{Pink}) = \frac{2}{3}P(\text{Red}|\text{Yellow}) + \frac{1}{3}P(\text{Red}|\text{Green})$$

$$P(\text{Red}|\text{Green}) = \frac{1}{4}P(\text{Red}|\text{Yellow}) + \frac{1}{4}$$

$$P(\text{Red}|\text{Yellow}) = \frac{2}{3}$$



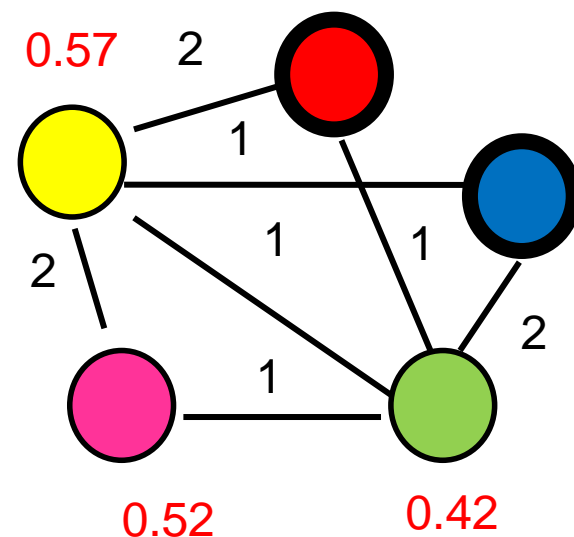
# Absorption probability

- The same idea can be applied to the case of undirected graphs
  - The absorbing nodes are still absorbing, so the edges to them are (implicitly) directed.

$$P(\text{Red}|\text{Pink}) = \frac{2}{3}P(\text{Red}|\text{Yellow}) + \frac{1}{3}P(\text{Red}|\text{Green})$$

$$P(\text{Red}|\text{Green}) = \frac{1}{5}P(\text{Red}|\text{Yellow}) + \frac{1}{5}P(\text{Red}|\text{Pink}) + \frac{1}{5}$$

$$P(\text{Red}|\text{Yellow}) = \frac{1}{6}P(\text{Red}|\text{Green}) + \frac{1}{3}P(\text{Red}|\text{Pink}) + \frac{1}{3}$$



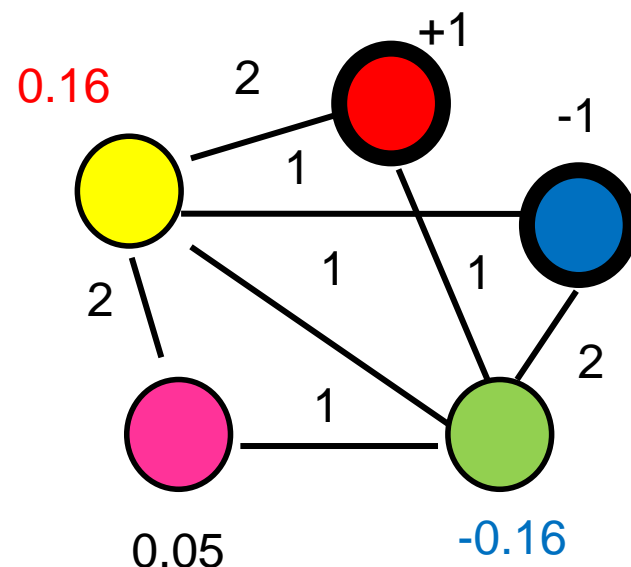
# Propagating values

- Assume that Red corresponds to a positive class and Blue to a negative class
  - We can compute a value for all the other nodes in the same way
    - This is the expected value for the node

$$V(\text{Pink}) = \frac{2}{3}V(\text{Yellow}) + \frac{1}{3}V(\text{Green})$$

$$V(\text{Green}) = \frac{1}{5}V(\text{Yellow}) + \frac{1}{5}V(\text{Pink}) + \frac{1}{5} - \frac{2}{5}$$

$$V(\text{Yellow}) = \frac{1}{6}V(\text{Green}) + \frac{1}{3}V(\text{Pink}) + \frac{1}{3} - \frac{1}{6}$$





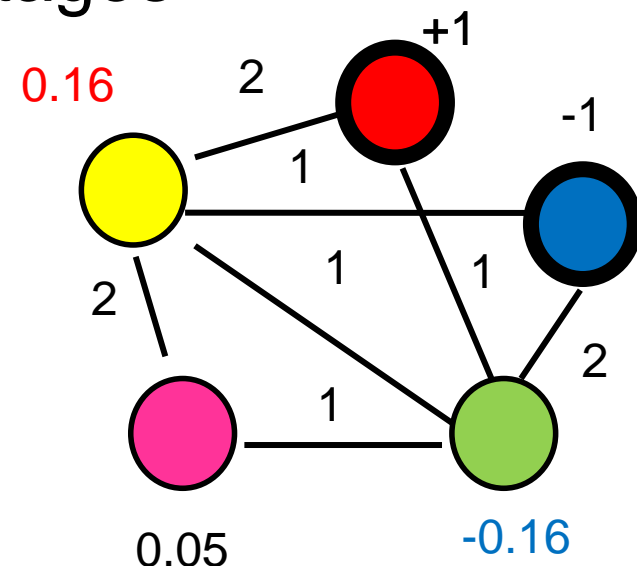
# Electrical networks and random walks

- If Red corresponds to a positive voltage and Blue to a negative voltage
- There are resistances on the edges inversely proportional to the weights
- The computed values are the voltages

$$V(\text{Pink}) = \frac{2}{3}V(\text{Yellow}) + \frac{1}{3}V(\text{Green})$$

$$V(\text{Green}) = \frac{1}{5}V(\text{Yellow}) + \frac{1}{5}V(\text{Pink}) + \frac{1}{5} - \frac{2}{5}$$

$$V(\text{Yellow}) = \frac{1}{6}V(\text{Green}) + \frac{1}{3}V(\text{Pink}) + \frac{1}{3} - \frac{1}{6}$$



# Transductive learning

- If we have a graph of relationships and some labels on these edges we can propagate them to the remaining nodes
  - E.g., a social network where some people are tagged as spammers
- This is a form of semi-supervised learning
  - We make use of the unlabeled data, and the relationships
- It is also called transductive learning because it does not produce a model, and labels only what is at hand.