

9.8 Node Arrays (`node_array`)

1. Definition

An instance A of the parameterized data type `node_array<E>` is a partial mapping from the node set of a graph G to the set of variables of type E , called the element type of the array. The domain I of A is called the index set of A and $A(v)$ is called the element at position v . A is said to be valid for all nodes in I . The array access operator $A[v]$ checks its precondition (A must be valid for v). The check can be turned off by compiling with the flag `-DLEDA_CHECKING_OFF`.

```
#include <LEDA/node_array.h >
```

2. Creation

`node_array<E> A;` creates an instance A of type `node_array<E>` with empty index set.

`node_array<E> A(const graph_t& G);`
creates an instance A of type `node_array<E>` and initializes the index set of A to the current node set of graph G .

`node_array<E> A(const graph_t& G, E x);`
creates an instance A of type `node_array<E>`, sets the index set of A to the current node set of graph G and initializes $A(v)$ with x for all nodes v of G .

`node_array<E> A(const graph_t& G, int n, E x);`
creates an instance A of type `node_array<E>` valid for up to n nodes of graph G and initializes $A(v)$ with x for all nodes v of G .
Precondition: $n \geq |V|$.
 A is also valid for the next $n - |V|$ nodes added to G .

3. Operations

`const graph_t& A.get_graph()` returns a reference to the graph of A .

`E& A[node v]` returns the variable $A(v)$.
Precondition: A must be valid for v .

`void A.init(const graph_t& G)` sets the index set I of A to the node set of G , i.e., makes A valid for all nodes of G .

`void A.init(const graph_t& G, E x)`
makes A valid for all nodes of G and sets $A(v) = x$ for all nodes v of G .

void *A*.init(*const graph_t& G*, *int n*, *E x*)
makes *A* valid for at most *n* nodes of *G* and sets $A(v) = x$ for all nodes *v* of *G*.
Precondition: $n \geq |V|$.
A is also valid for the next $n - |V|$ nodes added to *G*.

bool *A*.use_node_data(*const graph_t& G*)
use free data slots in the nodes of *G* (if available) for storing the entries of *A*. If no free data slot is available in *G*, an ordinary *node_array<E>* is created. The number of additional data slots in the nodes and edges of a graph can be specified in the *graph::graph(int n_slots, int e_slots)* constructor. The result is *true* if a free slot is available and *false* otherwise.

bool *A*.use_node_data(*const graph_t& G*, *E x*)
use free data slots in the nodes of *G* (if available) for storing the entries of *A* and initializes $A(v) = x$ for all nodes *v* of *G*. If no free data slot is available in *G*, an ordinary *node_array<E>* is created. The number of additional data slots in the nodes and edges of a graph can be specified in the *graph::graph(int n_slots, int e_slots)* constructor. The result is *true* if a free slot is available and *false* otherwise.

4. Implementation

Node arrays for a graph *G* are implemented by C++vectors and an internal numbering of the nodes and edges of *G*. The access operation takes constant time, *init* takes time $O(n)$, where *n* is the number of nodes in *G*. The space requirement is $O(n)$.

Remark: A node array is only valid for a bounded number of nodes of *G*. This number is either the number of nodes of *G* at the moment of creation of the array or it is explicitly set by the user. Dynamic node arrays can be realized by node maps (cf. section 9.11).

9.9 Edge Arrays (`edge_array`)

1. Definition

An instance A of the parameterized data type `edge_array<E>` is a partial mapping from the edge set of a graph G to the set of variables of type E , called the element type of the array. The domain I of A is called the index set of A and $A(e)$ is called the element at position e . A is said to be valid for all edges in I . The array access operator $A[e]$ checks its precondition (A must be valid for e). The check can be turned off by compiling with the flag `-DLEDA_CHECKING_OFF`.

```
#include < LEDA/edge_array.h >
```

2. Creation

`edge_array<E> A;` creates an instance A of type `edge_array<E>` with empty index set.

`edge_array<E> A(const graph_t& G);`
creates an instance A of type `edge_array<E>` and initializes the index set of A to be the current edge set of graph G .

`edge_array<E> A(const graph_t& G, E x);`
creates an instance A of type `edge_array<E>`, sets the index set of A to the current edge set of graph G and initializes $A(v)$ with x for all edges v of G .

`edge_array<E> A(const graph_t& G, int n, E x);`
creates an instance A of type `edge_array<E>` valid for up to n edges of graph G and initializes $A(e)$ with x for all edges e of G .
Precondition: $n \geq |E|$.
 A is also valid for the next $n - |E|$ edges added to G .

3. Operations

`const graph_t& A.get_graph()` returns a reference to the graph of A .

`E& A[edge e]` returns the variable $A(e)$.
Precondition: A must be valid for e .

`void A.init(const graph_t& G)` sets the index set I of A to the edge set of G , i.e., makes A valid for all edges of G .

`void A.init(const graph_t& G, E x)`
makes A valid for all edges of G and sets $A(e) = x$ for all edges e of G .

void `A.init(const graph_t& G, int n, E x)`
 makes A valid for at most n edges of G and sets $A(e) = x$ for all edges e of G .
Precondition: $n \geq |E|$.
 A is also valid for the next $n - |E|$ edges added to G .

bool `A.use_edge_data(const graph_t& G, E x)`
 use free data slots in the edges of G (if available) for storing the entries of A . The number of additional data slots in the nodes and edges of a graph can be specified in the `graph::graph(int n_slots, int e_slots)` constructor. The result is *true* if a free slot is available and *false* otherwise.

4. Implementation

Edge arrays for a graph G are implemented by C++vectors and an internal numbering of the nodes and edges of G . The access operation takes constant time, *init* takes time $O(n)$, where n is the number of edges in G . The space requirement is $O(n)$.

Remark: An edge array is only valid for a bounded number of edges of G . This number is either the number of edges of G at the moment of creation of the array or it is explicitly set by the user. Dynamic edge arrays can be realized by edge maps (cf. section 9.12).