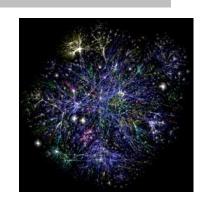
Information Networks

Link Analysis Ranking
Lecture 8





Why Link Analysis?

- § First generation search engines
 - § view documents as flat text files
 - § could not cope with size, spamming, user needs
- § Second generation search engines
 - § Ranking becomes critical
 - § use of Web specific data: Link Analysis
 - § shift from relevance to authoritativeness
 - § a success story for the network analysis



Outline

- § ...in the beginning...
- § previous work
- § some more algorithms
- § some experimental data
- § a theoretical framework



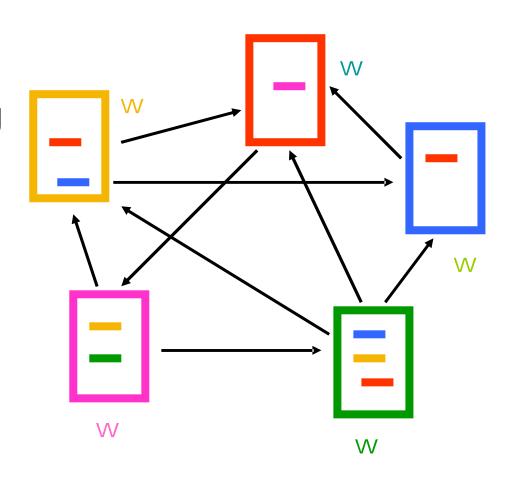
Link Analysis: Intuition

- § A link from page p to page q denotes endorsement
 - § page p considers page q an authority on a subject
 - § mine the web graph of recommendations
 - § assign an authority value to every page



Link Analysis Ranking Algorithms

- § Start with a collection of web pages
- § Extract the underlying hyperlink graph
- § Run the LAR algorithm on the graph
- § Output: an authority weight for each node





Link Analysis: Intuition

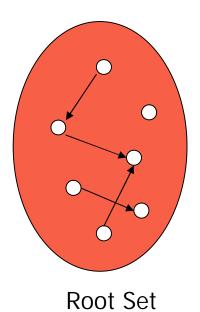
- § A link from page p to page q denotes endorsement
 - § page p considers page q an authority on a subject
 - § mine the web graph of recommendations
 - § assign an authority value to every page



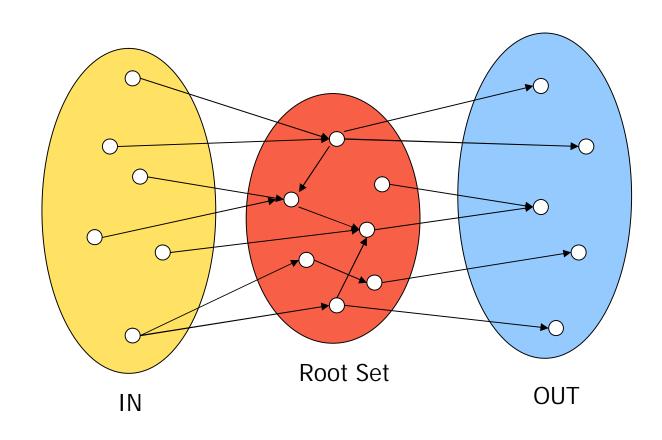
Algorithm input

- § Query independent: rank the whole Web
 - § PageRank (Brin and Page 98) was proposed as query independent
- § Query dependent: rank a small subset of pages related to a specific query
 - § HITS (Kleinberg 98) was proposed as query dependent

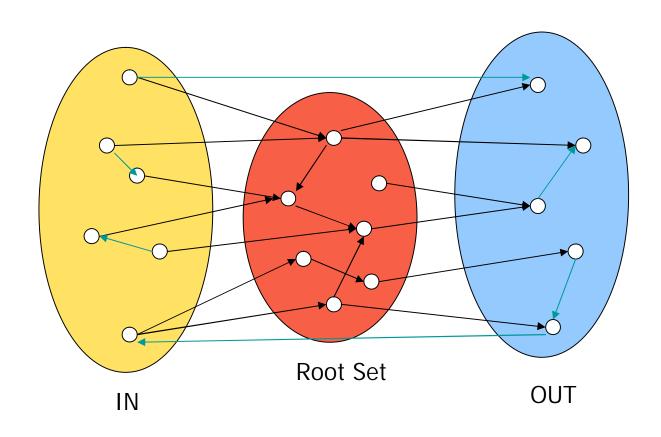




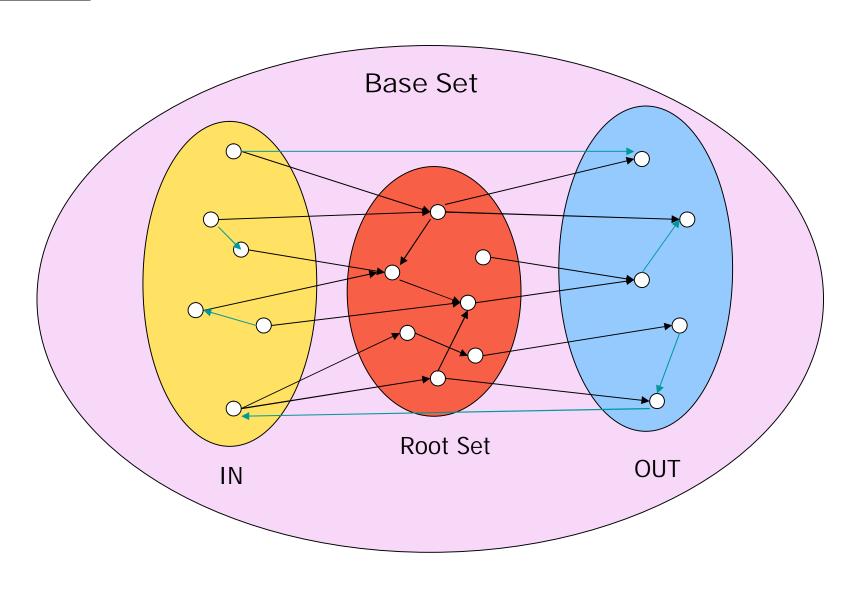














Link Filtering

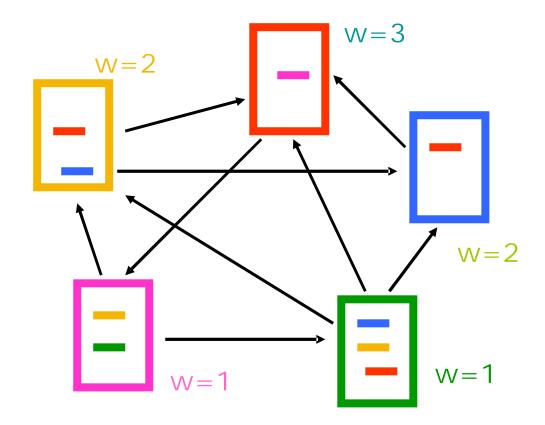
- § Navigational links: serve the purpose of moving within a site (or to related sites)
 - www.espn.com → www.espn.com/nba
 - www.yahoo.com → www.yahoo.it
 - www.espn.com → www.msn.com
- § Filter out navigational links
 - § same domain name
 - www.yahoo.com VS yahoo.com
 - § same IP address
 - § other way?



InDegree algorithm

§ Rank pages according to in-degree

$$\S$$
 $W_i = |B(i)|$



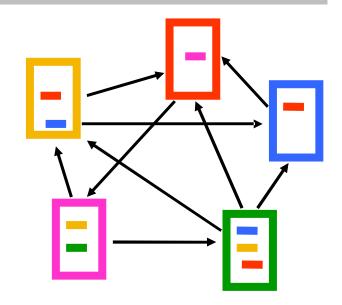
- 1. Red Page
- 2. Yellow Page
- 3. Blue Page
- 4. Purple Page
- 5. Green Page



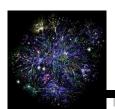
PageRank algorithm [BP98]

- § Good authorities should be pointed by good authorities
- § Random walk on the web graph
 - § pick a page at random
 - § with probability 1- α jump to a random page
 - § with probability α follow a random outgoing link
- § Rank according to the stationary distribution

§
$$PR(p) = \alpha \sum_{q \to p} \frac{PR(q)}{|F(q)|} + (1 - \alpha) \frac{1}{n}$$



- 1. Red Page
- 2. Purple Page
- 3. Yellow Page
- 4. Blue Page
- 5. Green Page



Markov chains

§ A Markov chain describes a discrete time stochastic process over a set of states

$$S = \{s_1, s_2, \dots s_n\}$$

according to a transition probability matrix

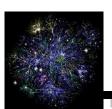
$$P = \{P_{ij}\}$$

- P_{ii} = probability of moving to state j when at state i
 - $\sum_{i} P_{ij} = 1$ (stochastic matrix)
- § Memorylessness property: The next state of the chain depends only at the current state and not on the past of the process (first order MC)
 - § higher order MCs are also possible

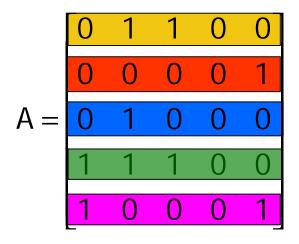


Random walks

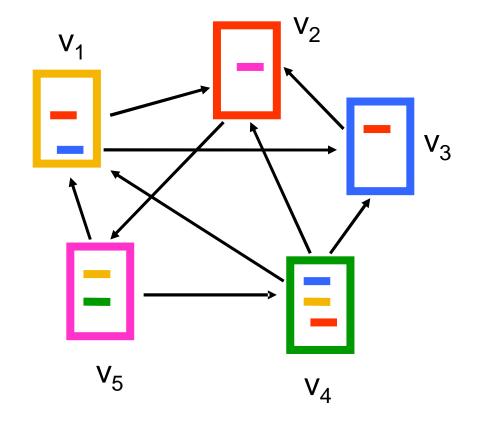
- § Random walks on graphs correspond to Markov Chains
 - § The set of states S is the set of nodes of the graph G
 - § The transition probability matrix is the probability that we follow an edge from one node to another



An example



$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 0 & 1/2 \end{bmatrix}$$

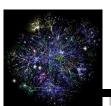




State probability vector

§ The vector $\mathbf{q}^t = (\mathbf{q}^t_1, \mathbf{q}^t_2, \dots, \mathbf{q}^t_n)$ that stores the probability of being at state i at time t § \mathbf{q}^0_i = the probability of starting from state i

$$q^t = q^{t-1} P$$



An example

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

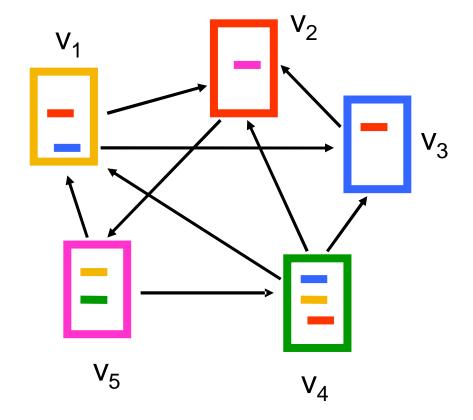
$$q^{t+1}_{1} = 1/3 \ q^{t}_{4} + 1/2 \ q^{t}_{5}$$

$$q^{t+1}_{2} = 1/2 \ q^{t}_{1} + q^{t}_{3} + 1/3 \ q^{t}_{4}$$

$$q^{t+1}_{3} = 1/2 \ q^{t}_{1} + 1/3 \ q^{t}_{4}$$

$$q^{t+1}_{4} = 1/2 \ q^{t}_{5}$$

$$q^{t+1}_{5} = q^{t}_{2}$$





Stationary distribution

- § A stationary distribution for a MC with transition matrix P, is a probability distribution π , such that $\pi = \pi P$
- § A MC has a unique stationary distribution if
 - § it is irreducible
 - the underlying graph is strongly connected
 - § it is aperiodic
 - for random walks, the underlying graph is not bipartite
- § The probability π_i is the fraction of times that we visited state i as $t \to \infty$
- § The stationary distribution is an eigenvector of matrix P
 - § the principal left eigenvector of P stochastic matrices have maximum eigenvalue 1



Computing the stationary distribution

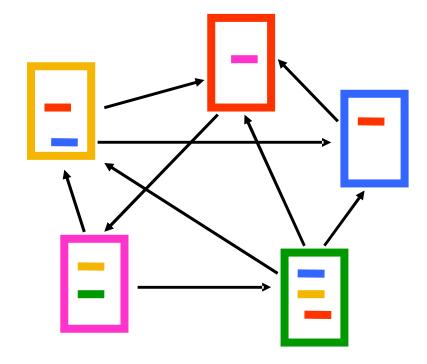
- § The Power Method
 - § Initialize to some distribution q⁰
 - § Iteratively compute q^t = q^{t-1}P
 - § After enough iterations $q^t \approx \pi$
 - § Power method because it computes $q^t = q^0P^t$
- § Why does it converge?
 - § follows from the fact that any vector can be written as a linear combination of the eigenvectors
 - $q^0 = V_1 + C_2 V_2 + ... C_n V_n$
- § Rate of convergence
 - § determined by λ_2^t

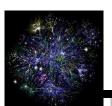


§ Vanilla random walk

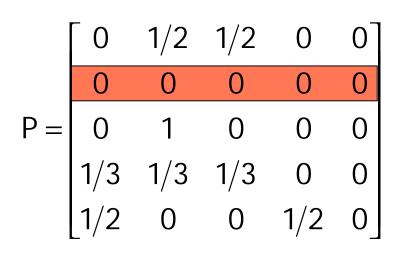
§ make the adjacency matrix stochastic and run a random walk

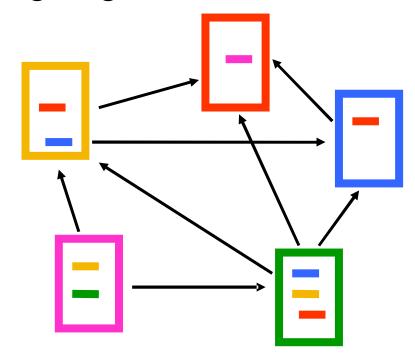
$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$





- § What about sink nodes?
 - § what happens when the random walk moves to a node without any outgoing inks?



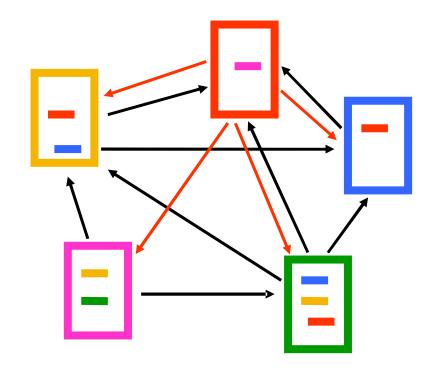




- § Replace these row vectors with a vector v
 - § typically, the uniform vector

$$P' = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

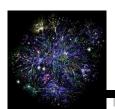
$$P' = P + dv^{T} \qquad d = \begin{cases} 1 & \text{if i is sink} \\ 0 & \text{otherwise} \end{cases}$$



- § How do we guarantee irreducibility?
 - § add a random jump to vector v with prob α
 - typically, to a uniform vector

$$\mathsf{P''} = \alpha \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1 & 0 & 0 & 0 \\ 1/3 & 1/3 & 1/3 & 0 & 0 \\ 1/2 & 0 & 0 & 0 & 1/2 \end{bmatrix} + (1-\alpha) \begin{bmatrix} 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \end{bmatrix}$$

 $P'' = \alpha P' + (1-\alpha)uv^T$, where u is the vector of all 1s



Effects of random jump

- § Guarantees irreducibility
- § Motivated by the concept of random surfer
- § Offers additional flexibility
 - § personalization
 - § anti-spam
- § Controls the rate of convergence
 - § the second eigenvalue of matrix P" is α

A PageRank algorithm

§ Performing vanilla power method is now too expensive – the matrix is not sparse

$$q^{0} = v$$

$$t = 1$$

$$repeat$$

$$q^{t} = (P^{t})^{T} q^{t-1}$$

$$\delta = \|q^{t} - q^{t-1}\|$$

$$t = t + 1$$

$$until \delta < \epsilon$$

Efficient computation of $y = (P'')^T x$

$$y = aP^{T}x$$

$$\beta = ||x||_{1} - ||y||_{1}$$

$$y = y + \beta v$$



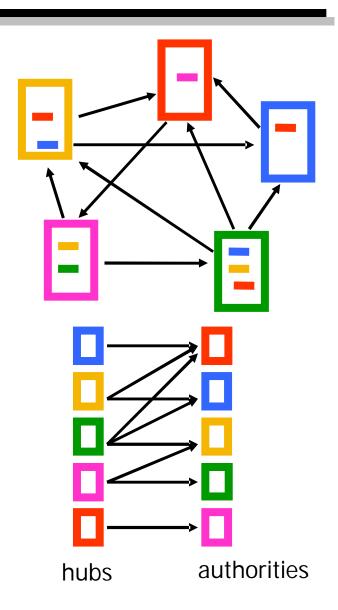
Research on PageRank

- § Specialized PageRank
 - § personalization [BP98]
 - instead of picking a node uniformly at random favor specific nodes that are related to the user
 - § topic sensitive PageRank [H02]
 - compute many PageRank vectors, one for each topic
 - estimate relevance of query with each topic
 - produce final PageRank as a weighted combination
- § Updating PageRank [Chien et al 2002]
- § Fast computation of PageRank
 - § numerical analysis tricks
 - § node aggregation techniques
 - § dealing with the "Web frontier"



Hubs and Authorities [K98]

- § Authority is not necessarily transferred directly between authorities
- § Pages have double identity
 - § hub identity
 - § authority identity
- § Good hubs point to good authorities
- § Good authorities are pointed by good hubs



HITS Algorithm

- § Initialize all weights to 1.
- § Repeat until convergence
 - § O operation: hubs collect the weight of the authorities

$$h_i = \sum_{j:i \to j} a_j$$

§ I operation: authorities collect the weight of the hubs

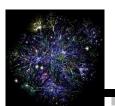
$$a_i = \sum_{j: j \to i} h_j$$

§ Normalize weights under some norm



HITS and eigenvectors

- § The HITS algorithm is a power-method eigenvector computation
 - § in vector terms $\mathbf{a}^t = \mathbf{A}^T \mathbf{h}^{t-1}$ and $\mathbf{h}^t = \mathbf{A} \mathbf{a}^{t-1}$
 - § so $a = A^{T}Aa^{t-1}$ and $h^{t} = AA^{T}h^{t-1}$
 - § The authority weight vector a is the eigenvector of A^TA and the hub weight vector h is the eigenvector of AA^T
 - § Why do we need normalization?
- § The vectors a and h are singular vectors of the matrix A



Singular Value Decomposition

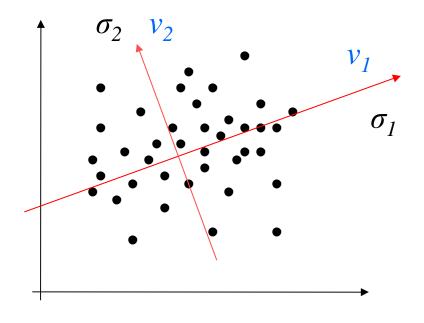
- § r: rank of matrix A
- § $\sigma_1 \ge \sigma_2 \ge ... \ge \sigma_r$: singular values (square roots of eig-vals AA^T , A^TA)
- § $\ddot{u}_1, \ddot{u}_2, \ddot{l}_r$: left singular vectors (eig-vectors of AA^T)
- § V_1 , V_2 , V_r : right singular vectors (eig-vectors of A^TA)

$$A = \sigma_1 \ddot{\mathbf{u}}_1 \ddot{\mathbf{v}}_1^{\mathsf{T}} + \sigma_2 \ddot{\mathbf{u}}_2 \ddot{\mathbf{v}}_2^{\mathsf{T}} + \mathbf{1} + \sigma_r \ddot{\mathbf{u}}_r \ddot{\mathbf{v}}_r^{\mathsf{T}}$$



Singular Value Decomposition

- § Linear trend v in matrix A:
 - § the tendency of the row vectors of A to align with vector v
 - strength of the linear trend:
 Av
- § SVD discovers the linear trends in the data
- § **u**_i, **v**_i: the i-th strongest linear trends
- § o_i: the strength of the i-th strongest linear trend

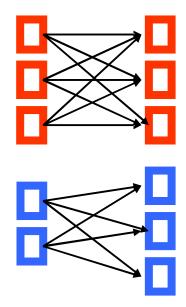


§ HITS discovers the strongest linear trend in the authority space



HITS and the TKC effect

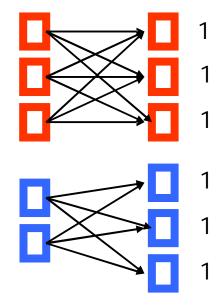
- § The HITS algorithm favors the most dense community of hubs and authorities
 - § Tightly Knit Community (TKC) effect





HITS and the TKC effect

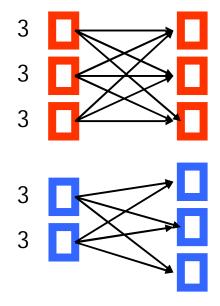
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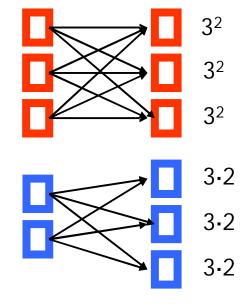
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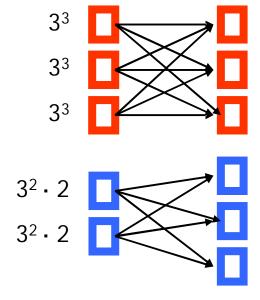


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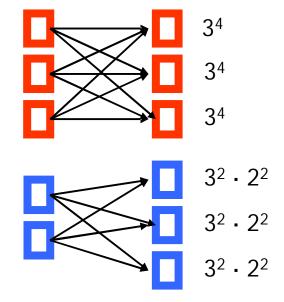


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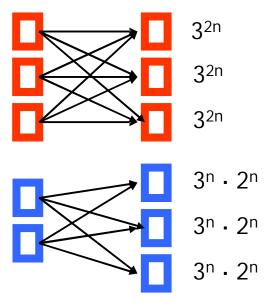
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- § The HITS algorithm favors the most dense community of hubs and authorities
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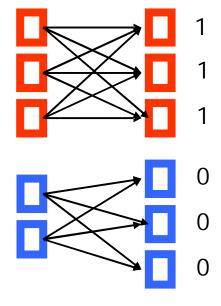
weight of node p is proportional to the number of (BF)ⁿ paths that leave node p



after n iterations



- § The HITS algorithm favors the most dense community of hubs and authorities
 - § Tightly Knit Community (TKC) effect



after normalization with the max element as $n \rightarrow \infty$



Outline

- § ...in the beginning...
- § previous work
- § some more algorithms
- § some experimental data
- § a theoretical framework



Previous work

- § The problem of identifying the most important nodes in a network has been studied before in social networks and bibliometrics
- § The idea is similar
 - § A link from node p to node q denotes endorsement
 - § mine the network at hand
 - § assign an centrality/importance/standing value to every node



Social network analysis

- § Evaluate the centrality of individuals in social networks
 - § degree centrality
 - the (weighted) degree of a node
 - § distance centrality
 - the average (weighted) distance of a node to the rest in the graph $D_c(v) = \frac{1}{\sum_{u,v} d(v,u)}$

 the average number of (weighted) shortest paths that use node v

$$B_{c}(v) = \sum_{s \neq v \neq t} \frac{\sigma_{st}(v)}{\sigma_{st}}$$



Random walks on undirected graphs

§ In the stationary distribution of a random walk on an undirected graph, the probability of being at node i is proportional to the (weighted) degree of the vertex

§ Random walks on undirected graphs are not "interesting"



Counting paths – Katz 53

- § The importance of a node is measured by the weighted sum of paths that lead to this node
- § Am[i,j] = number of paths of length m from i to j
- § Compute

$$P = bA + b^2A^2 +] + b^mA^m +] = (I - bA)^{-1} - I$$

- § converges when $b < \lambda_1(A)$
- § Rank nodes according to the column sums of the matrix P



Bibliometrics

- § Impact factor (E. Garfield 72)
 - § counts the number of citations received for papers of the journal in the previous two years
- § Pinsky-Narin 76
 - § perform a random walk on the set of journals
 - P_{ij} = the fraction of citations from journal i that are directed to journal j



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