## Λ14 Διαδικτυακά Κοινωνικά Δίκτυα και Μέσα

**Link Prediction** 

#### Motivation

- Recommending new friends in online social networks.
- Predicting the participation of actors in events
- Suggesting interactions between the members of a company/organization that are external to the hierarchical structure of the organization itself.
- Predicting *connections* between members of terrorist organizations who have not been directly observed to work together.
- Suggesting collaborations between researchers based on coauthorship.
- Overcoming the data-sparsity problem in recommender systems using collaborative filtering

### Motivation

#### In social networks:

- Increases user engagement
- Controls the growth of the network

## Outline

 Estimating a score for each edge (seminal work of Liben-Nowell&Kleinberg)

Classification approach

The who to follow service at Twitter

#### **Problem Definition**

Link prediction problem: Given the links in a social network at time *t*, *predict* the edges that will be added to the network during the time interval from time *t* to a given future time *t'* 

- Based solely on the *topology* of the network (social proximity) (the more general problem also considers attributes of the nodes and links)
- Different from the problem of *inferring missing* (hidden) links (there is a temporal aspect)

To save experimental effort in the laboratory or in the field

## Problem Formulation (details)

Consider a social network G = (V, E) where each edge  $e = \langle u, v \rangle \in E$  represents an interaction between u and v that took place at a particular time t(e)

(multiple interactions between two nodes as parallel edges with different timestamps)

For two times, t < t', let G[t, t'] denote subgraph of G consisting of all edges with a timestamp between t and t'

■ For four times,  $t_0 < t'_0 < t_1 < t'_1$ , given  $G[t_0, t'_0]$ , we wish to output a list of edges not in  $G[t_0, t'_0]$  that are predicted to appear in  $G[t_1, t'_1]$ 

- $\checkmark$  [ $t_o$ ,  $t'_o$ ] training interval
- $\checkmark$  [ $t_1$ ,  $t'_1$ ] test interval

## Problem Formulation (details)

#### Prediction for a subset of nodes

Two parameters:  $\kappa_{\text{training}}$  and  $\kappa_{\text{test}}$ 

Core: all nodes that are incident to at least  $\kappa_{\text{training}}$  edges in  $G[t_0, t'_0]$ , and at least  $\kappa_{\text{test}}$  edges in  $G[t_1, t'_1]$ 

Predict new edges between the nodes in Core

### Example Dataset: co-authorship

	training period			Core		
	authors	papers	collaborations <sup>1</sup>	authors	$ E_{old} $	$ E_{new} $
astro-ph	5343	5816	41852	1561	6178	5751
cond-mat	5469	6700	19881	1253	1899	1150
gr-qc	2122	3287	5724	486	519	400
hep-ph	5414	10254	47806	1790	6654	3294
hep-th	5241	9498	15842	1438	2311	1576

 $t_0$  = 1994,  $t'_0$  = 1996: training interval -> [1994, 1996]

 $t_1$  = 1997,  $t'_1$  = 1999: **test interval** -> **[1997, 1999]** 

- $-G_{collab} = \langle V, E_{old} \rangle = G[1994, 1996]$
- $\boldsymbol{E}_{\text{new}}\!\!:\!$  authors in V that co-author a paper during the test interval but not during the training interval

 $\kappa_{\text{training}} = 3$ ,  $\kappa_{\text{test}} = 3$ : Core consists of all authors who have written at least 3 papers during the training period and at least 3 papers during the test period

## Methods for Link Prediction (outline)

- Assign a connection weight score(x, y) to each pair of nodes <x, y> based on the input graph
- Produce a ranked list of decreasing order of score

- We can consider all links incident to a specific node x, and recommend to x the top ones
- If we focus to a specific x, the score can be seen as a centrality measure for x

## Methods for Link Prediction (outline)

How to assign the score(x, y) between two nodes x and y?

✓ Some form of similarity or node proximity

The larger the overlap of the neighbors of two nodes, the more likely the nodes to be linked in the future

Let  $\Gamma(x)$  denote the set of neighbors of x in  $G_{old}$ 

#### Common neighbors:

$$score(x, y) = |\Gamma(x) \cap \Gamma(y)|$$

A adjacency matrix  $A_{x,y}^2$ : Number of different paths of length 2

#### Jaccard coefficient:

$$score(x,y) = \frac{|\Gamma(x) \cap \Gamma(y)|}{|\Gamma(x) \cup \Gamma(y)|}$$

The probability that both x and y have a feature for a randomly selected feature that either x or y has

### Adamic/Adar

$$score(x, y) = \sum_{z \in \Gamma(x) \cap \Gamma(y)} \frac{1}{\log |\Gamma(z)|}$$

- $\checkmark$  Assigns large weights to common neighbors z of x and y which themselves have few neighbors (weight rare features more heavily)
- Neighbors who are linked with 2 nodes are assigned weight = 1/log(2) = 1.4
- Neighbors who are linked with 5 nodes are assigned weight = 1/log(5) = 0.62

#### Preferential attachment

Based on the premise that the probability that a new edge has node x as its endpoint is proportional to  $|\Gamma(x)|$ , i.e., nodes like to form ties with 'popular' nodes

$$score(x, y) = |\Gamma(x)||\Gamma(y)|$$

✓ Researchers found empirical evidence to suggest that co-authorship is correlated with the product of the neighborhood sizes

This depends on the degrees of the nodes not on their neighbors per se

- 1. Overlap
- 2. Jaccard
- 3. Adamic/Adar
- 4. Preferential attachment

## Methods for Link Prediction: Shortest Path

For x, 
$$y \in V \times V - E_{old}$$
,  
score(x, y) = (negated) length of shortest path between  
x and y

✓ If there are more than n pairs of nodes tied for the shortest path length, order them at random.

Not just the shortest, but *all* paths between two nodes

## $Katz_{\beta}$ measure

$$\mathsf{score}(x,y) := \sum_{\ell=1}^\infty \beta^\ell \cdot |\mathsf{paths}_{x,y}^{\langle \ell \rangle}|$$

Sum over all paths of length /

6 > 0 (< 1) is a parameter of the predictor, exponentially damped to count short paths more heavily

✓ Small  $\theta$  predictions much like common neighbors  $\theta = 0$ , degree, maximal  $\theta$ , eigenvalue

## $Katz_{\beta}$ measure

$$\mathrm{score}(x,y) := \sum^{\infty} \beta^{\ell} \cdot |\mathrm{paths}_{x,y}^{\langle \ell \rangle}|$$

$$\sum_{l=1}^{\infty} \beta^l \cdot |\mathsf{paths}_{xy}^{\langle l \rangle}| = \beta A_{xy} + \beta^2 (A^2)_{xy} + \beta^3 (A^3)_{xy} + \cdots$$

Closed form: 
$$(I - \beta A)^{-1} - I$$

- *Unweighted* version, in which path<sub>x,y</sub><sup>(1)</sup> = **1**, if x and y have collaborated, **0** otherwise
- Weighted version, in which  $path_{x,y}^{(1)} = \text{#times} x$  and y have collaborated

Consider a *random walk* on  $G_{old}$  that starts at x and iteratively moves to a neighbor of x chosen uniformly at random from  $\Gamma(x)$ .

The Hitting Time  $H_{x,y}$  from x to y is the expected number of steps it takes for the random walk starting at x to reach y.

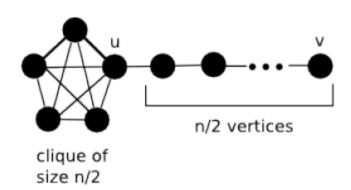
$$score(x, y) = -H_{x,y}$$

The Commute Time  $C_{x,y}$  from x to y is the expected number of steps to travel from x to y and from y to x

$$score(x, y) = -(H_{x,y} + H_{y,x})$$

Not symmetric, can be shown

$$h_{vu} = \Theta(n^2)$$
$$h_{uv} = \Theta(n^3)$$



Example: hit time in a line



Can also consider stationary-normed versions:

(to counteract the fact that  $H_{x,y}$  is rather small when y is a node with a large stationary probability)

score(x, y) = 
$$-H_{x,y} \pi_y$$
  
score(x, y) =  $-(H_{x,y} \pi_y + H_{y,x} \pi_x)$ 

The hitting time and commute time measures are sensitive to parts of the graph far away from x and y -> periodically reset the walk

Random walk on  $G_{old}$  that starts at x and has a probability  $\alpha$  of returning to x at each step

Rooted Page Rank: Starts from x, with probability (1 - a) moves to a random neighbor and with probability a returns to x

score(x, y) = stationary probability of y in a rooted PageRank

### SimRank

Two objects are *similar*, if they are *related to similar objects* 

Two objects x and y are similar, if they are related to objects a and b respectively and a and b are themselves similar

Base case: similarity(x, x) = 1

$$\mathrm{similarity}(x,y) := \gamma \cdot \frac{\sum_{a \in \Gamma(x)} \sum_{b \in \Gamma(y)} \mathrm{similarity}(a,b)}{|\Gamma(x)| \cdot |\Gamma(y)|}$$

Average similarity between neighbors of x and neighbors of y

$$score(x, y) = similarity(x, y)$$

Introduced for directed graphs, I(x): in-neighbors of x

Average similarity between in-neighbors of  $\alpha$  and in-neighbors of b C a constant between 0 and 1  $n^2$  equations

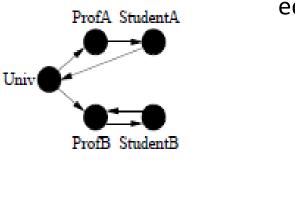
#### Iterative computation

 $s_0(x, y) = 1$  if x = y and 0 otherwise  $s_{k+1}$  based on the  $s_k$  values of its (in-neighbors) computed at iteration k

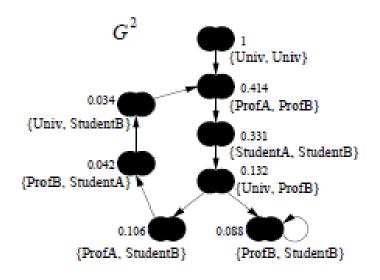
ProfA StudentA

Graph G<sup>2</sup>:

A node for each pair of nodes  $(x, y) \rightarrow (a, b)$ , if  $x \rightarrow a$  and  $y \rightarrow b$  Scores *flow* from a node to its neighbors C gives the rate of *decay* as similarity flows across edges (C = 0.8) in the example)

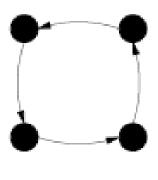


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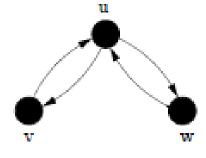


Self-loops
Prune by considering only nodes within a a radius

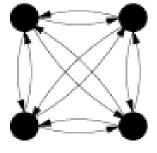
Expected Meeting Distance (EMD): how soon two random surfers are expected to meet at the same node if they started at nodes x and y and randomly walked (in lock step) the graph backwards



 $=\infty$ 



 $m(u, v) = m(u,w) = \infty$ , m(v, w) = 1v and w are much more similar than u is to v or w.



= 3, a lower similarity than between v and w but higher than between u and v (or u and w).

Let us consider G<sup>2</sup>

A node (a, b) as a state of the tour in G: if a moves to c, b moves to d in G, then (a, b) moves to (c, d) in  $G^2$ 

A tour in  $G^2$  of length n represents a pair of tours in G where each has length n

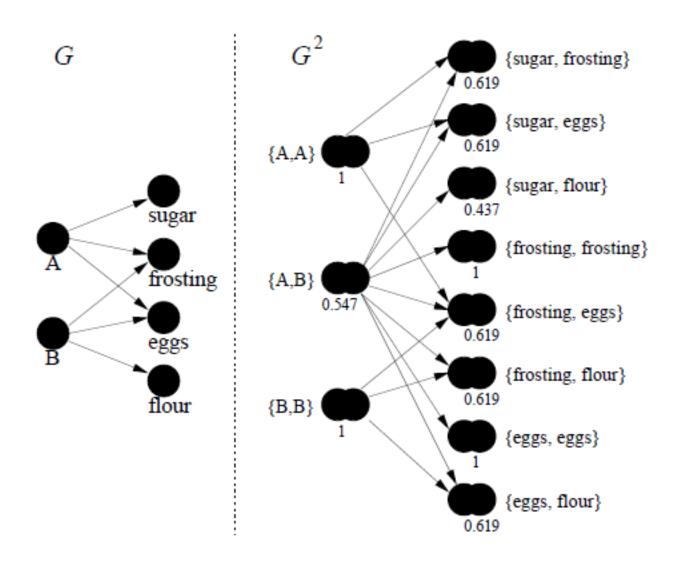
What are the states in G<sup>2</sup> that correspond to "meeting" points?

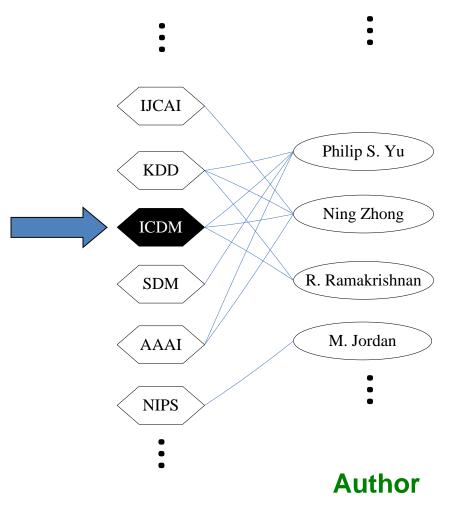
What are the states in G<sup>2</sup> that correspond to "meeting" points? Singleton nodes (common neighbors)

The EMD m(a, b) is just the expected distance (hitting time) in  $G^2$  between (a, b) and any singleton node

The sum is taken over all walks that start from (a, b) and end at a singleton node

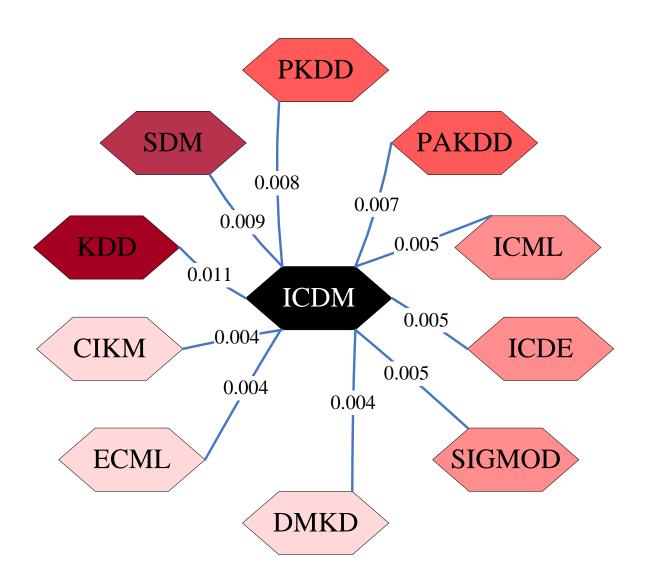
## SimRank for bipartite graphs





Q: What is most related conference to ICDM?

Conference



# Methods for Link Prediction: based on paths

- 1. Shortest paths
- 2. Katz
- 3. Hitting and commute time
- 4. Rooted page rank
- 5. SimRank

### Methods for Link Prediction: other

### Low rank approximations

M adjacency matrix , represent M with a lower rank matrix  $M_k$ 

Apply SVD (singular value decomposition)

The rank-k matrix that best approximates M

## Singular Value Decomposition

$$A = U \sum_{[n \times r]} V^{\mathsf{T}} = \begin{bmatrix} \vec{\mathsf{u}}_1 & \vec{\mathsf{u}}_2 & \cdots & \vec{\mathsf{u}}_r \end{bmatrix} \begin{bmatrix} \sigma_1 & & & & \\ & \sigma_2 & & \\ & & \ddots & & \\ & & & \sigma_r \end{bmatrix} \begin{bmatrix} \vec{\mathsf{v}}_1 \\ \vec{\mathsf{v}}_2 \\ \vdots \\ \vec{\mathsf{v}}_r \end{bmatrix}$$

- r : rank of matrix A
- $\sigma_1 \ge \sigma_2 \ge ... \ge \sigma_r$ : singular values (square roots of eig-vals  $AA^T$ ,  $A^TA$ )
- $\vec{\mathbf{u}}_1, \vec{\mathbf{u}}_2, \cdots, \vec{\mathbf{u}}$  left singular vectors (eig-vectors of  $\mathbf{A}\mathbf{A}^{\mathsf{T}}$ )
- $\vec{V}_1, \vec{V}_2, \cdots$ ; ight singular vectors (eig-vectors of  $\vec{A}^T \vec{A}$ )

• 
$$A = \sigma_1 \vec{u}_1 \vec{v}_1^T + \sigma_2 \vec{u}_2 \vec{v}_2^T + \dots + \sigma_r \vec{u}_r \vec{v}_r^T$$

#### Methods for Link Prediction: other

### **Unseen Bigrams**

Unseen bigrams: pairs of word that co-occur in a test corpus, but not in the corresponding training corpus

Not just score(x, y) but score(z, y) for nodes z that are similar to  $x ext{ --- } S_v^{(\delta)}$  the  $\delta$  nodes  $most\ related\ to\ x$ 

$$\mathsf{score}^*_{unweighted}(x,y) \ := \ \left| \{z : z \in \Gamma(y) \cap S_x^{\langle \delta \rangle} \} \right|$$
 
$$\mathsf{score}^*_{weighted}(x,y) \ := \ \sum_{z \in \Gamma(u) \cap S_x^{\langle \delta \rangle}} \mathsf{score}(x,z)$$

# Methods for Link Prediction: High-level approaches

### Clustering

- Compute score(x, y) for all edges in E<sub>old</sub>
- Delete the (1-p) fraction of the edges whose score is the lowest, for some parameter p
- Recompute score(x, y) for all pairs in the subgraph

#### How to Evaluate the Prediction (outline)

Each link predictor p outputs a ranked list  $L_p$  of pairs in  $V \times V - E_{old}$ : predicted new collaborations in decreasing order of confidence

In this paper, focus on Core, thus

$$E*_{new} = E_{new} \cap (Core \times Core) = |E*_{new}|$$

Evaluation method: Size of the intersection of

- the first n edge predictions from  $L_p$  that are in Core  $\times$  Core, and
- the set E\*<sub>new</sub>

\* How many of the (relevant) top-n predictions are correct (precision?)

## Evaluation: baseline

#### **Baseline: random predictor**

Randomly select pairs of authors who did not collaborate in the training interval

Probability that a random prediction is correct:

$$\frac{|E_{new}|}{\binom{|\mathsf{Core}|}{2} - |E_{old}|}$$

In the datasets, from 0.15% (cond-mat) to 0.48% (astro-ph)

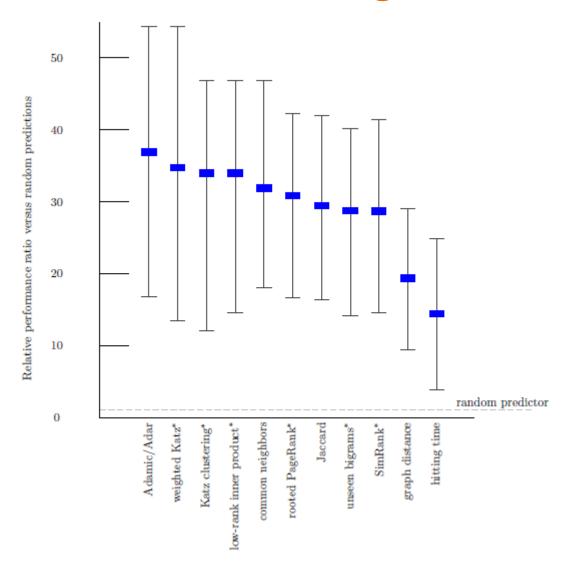
# Evaluation: Factor improvement over random

predictor	astro-ph	cond-mat	gr-qc	hep-ph	hep-th
probability that a random prediction is correct	0.475%	0.147%	0.341%	0.207%	0.153%
graph distance (all distance-two pairs)	9.4	25.1	21.3	12.0	29.0
common neighbors	18.0	40.8	27.1	26.9	46.9
preferential attachment	4.7	6.0	7.5	15.2	7.4
Adamic/Adar	16.8	54.4	30.1	33.2	50.2
Jaccard	16.4	42.0	19.8	27.6	41.5
SimRank $\gamma = 0.8$	14.5	39.0	22.7	26.0	41.5
hitting time	6.4	23.7	24.9	3.8	13.3
hitting time—normed by stationary distribution	5.3	23.7	11.0	11.3	21.2
commute time	5.2	15.4	33.0	17.0	23.2
commute time—normed by stationary distribution	5.3	16.0	11.0	11.3	16.2
rooted PageRank $\alpha = 0.01$	10.8	27.8	33.0	18.7	29.1
$\alpha = 0.05$	13.8	39.6	35.2	24.5	41.1
$\alpha = 0.15$	16.6	40.8	27.1	27.5	42.3
$\alpha = 0.30$	17.1	42.0	24.9	29.8	46.5
$\alpha = 0.50$	16.8	40.8	24.2	30.6	46.5
Katz (weighted) $\beta = 0.05$	3.0	21.3	19.8	2.4	12.9
$\beta = 0.005$	13.4	54.4	30.1	24.0	51.9
$\beta = 0.0005$	14.5	53.8	30.1	32.5	51.5
Katz (unweighted) $\beta = 0.05$	10.9	41.4	37.4	18.7	47.7
$\beta = 0.005$	16.8	41.4	37.4	24.1	49.4
$\beta = 0.0005$	16.7	41.4	37.4	24.8	49.4

# Evaluation: Factor improvement over random

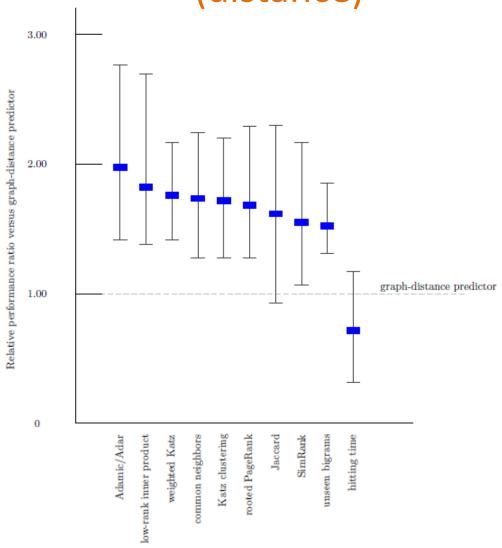
predictor		astro-ph	cond-mat	gr-qc	hep-ph	hep-th
probability that a random	0.475%	0.147%	0.341%	0.207%	0.153%	
graph distance (all distance	e-two pairs)	9.4	25.1	21.3	12.0	29.0
common neighbors		18.0	40.8	27.1	26.9	46.9
Low-rank approximation:	rank = 1024	15.2	53.8	29.3	34.8	49.8
Inner product	rank = 256	14.6	46.7	29.3	32.3	46.9
	rank = 64	13.0	44.4	27.1	30.7	47.3
	rank = 16	10.0	21.3	31.5	27.8	35.3
	rank = 4	8.8	15.4	42.5	19.5	22.8
	rank = 1	6.9	5.9	44.7	17.6	14.5
Low-rank approximation:	rank = 1024	8.2	16.6	6.6	18.5	21.6
Matrix entry	rank = 256	15.4	36.1	8.1	26.2	37.4
	rank = 64	13.7	46.1	16.9	28.1	40.7
	rank = 16	9.1	21.3	26.4	23.1	34.0
	rank = 4	8.8	15.4	39.6	20.0	22.4
	rank = 1	6.9	5.9	44.7	17.6	14.5
Low-rank approximation:	rank = 1024	11.4	27.2	30.1	27.0	32.0
Katz ( $\beta = 0.005$ )	rank = 256	15.4	42.0	11.0	34.2	38.6
	rank = 64	13.1	45.0	19.1	32.2	41.1
	rank = 16	9.2	21.3	27.1	24.8	34.9
	rank = 4	7.0	15.4	41.1	19.7	22.8
	rank = 1	0.4	5.9	44.7	17.6	14.5
unseen bigrams	common neighbors, $\delta = 8$	13.5	36.7	30.1	15.6	46.9
(weighted)	common neighbors, $\delta = 16$	13.4	39.6	38.9	18.5	48.6
	Katz ( $\beta = 0.005$ ), $\delta = 8$	16.8	37.9	24.9	24.1	51.1
	Katz ( $\beta = 0.005$ ), $\delta = 16$	16.5	39.6	35.2	24.7	50.6
unseen bigrams	common neighbors, $\delta = 8$	14.1	40.2	27.9	22.2	39.4
(unweighted)	common neighbors, $\delta = 16$	15.3	39.0	42.5	22.0	42.3
	Katz ( $\beta = 0.005$ ), $\delta = 8$	13.1	36.7	32.3	21.6	37.8
	Katz ( $\beta = 0.005$ ), $\delta = 16$	10.3	29.6	41.8	12.2	37.8
clustering:	ho = 0.10	7.4	37.3	46.9	32.9	37.8
Katz $(\beta_1 = 0.001, \beta_2 = 0.1)$	,	12.0	46.1	46.9	21.0	44.0
	$\rho = 0.20$	4.6	34.3	19.8	21.2	35.7
	$\rho = 0.25$	3.3	27.2	20.5	19.4	17.4

### Evaluation: Average relevance performance

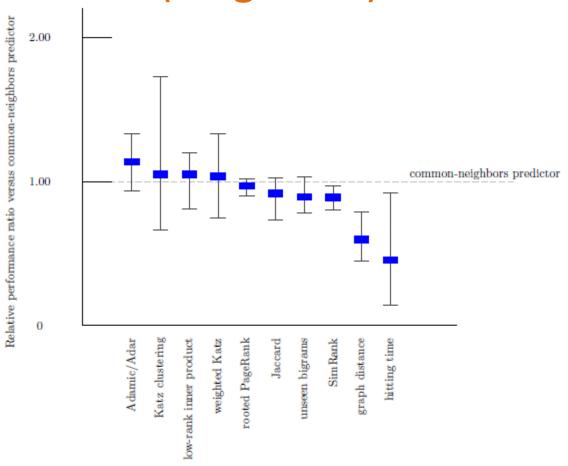


- average ratio over the five datasets of the given predictor's performance versus a baseline predictor's performance.
- the error bars indicate the minimum and maximum of this ratio over the five datasets.
- the parameters for the starred predictors are: (1) for weighted Katz,  $\beta$ = 0.005; (2) for Katz clustering,  $\beta$ 1 = 0.001;  $\rho$  = 0.15;  $\beta$ 2 = 0.1; (3) for low-rank inner product, rank = 256; (4) for rooted Pagerank,  $\alpha$  = 0.15; (5) for unseen bigrams, unweighted, common neighbors with  $\delta$  = 8; and (6) for SimRank, C ( $\gamma$ ) = 0.8.

# Evaluation: Average relevance performance (distance)



# Evaluation: Average relevance performance (neighbors)



## Evaluation: prediction overlap

	Adamic/Adar	Katz clustering	common neighbors	hitting time	Jaccard's coefficient	weighted Katz	low-rank inner product	rooted Pagerank	SimRank	unseen bigrams
Adamic/Adar	1150	638	520	193	442	1011	905	528	372	486
Katz clustering		1150	411	182	285	630	623	347	245	389
common neighbors			1150	135	506	494	467	305	332	489
hitting time				1150	87	191	192	247	130	156
Jaccard's coefficient					1150	414	382	504	845	458
weighted Katz						1150	1013	488	344	474
low-rank inner product							1150	453	320	448
rooted Pagerank								1150	678	461
SimRank									1150	423
unseen bigrams										1150

❖ How much similar are the predictions made by the different methods?

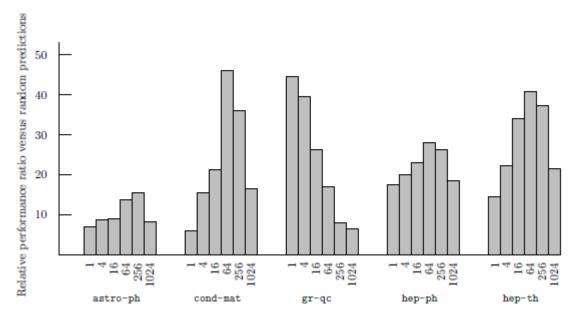
Why?

correct

	Adamic/Adar	Katz clustering	common neighbors	hitting time	Jaccard's coefficient	weighted Katz	low-rank inner product	rooted Pagerank	SimRank	unseen bigrams
Adamic/Adar	92	65	53	22	43	87	72	44	36	49
Katz clustering		78	41	20	29	66	60	31	22	37
common neighbors			69	13	43	52	43	27	26	40
hitting time				40	8	22	19	17	9	15
Jaccard's coefficient					71	41	32	39	51	43
weighted Katz						92	75	44	32	51
low-rank inner product							79	39	26	46
rooted Pagerank								69	48	39
SimRank									66	34 68
unseen bigrams										68

#### **Evaluation**: datasets

\* How much does the performance of the different methods depends on the dataset?



- (rank) On 4 of the 5 datasets best at an intermediate rank
  On qr-qc, best at rank 1, does it have a "simpler" structure"?
- On hep-ph, preferential attachment the best
- Why is astro-ph "difficult"?

The culture of physicists and physics collaboration

### Evaluation: small world

The shortest path even in unrelated disciplines is often very short

Basic classifier on graph distances does not work

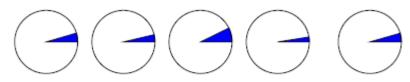
#### Evaluation: restricting to distance three

Many pairs of authors separated by a graph distance of 2 will not collaborate and Many pairs who collaborate are at distance greater than 2

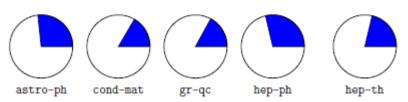
Disregard all distance 2 pairs (do not just "close" triangles)

	astro-ph	cond-mat	gr-qc	hep-ph	hep-th
# pairs at distance two	33862	5145	935	37687	7545
# new collaborations at distance two	1533	190	68	945	335
# new collaborations	5751	1150	400	3294	1576

Proportion of distance-two pairs that form an edge:



Proportion of new edges that are between distance-two pairs:



$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						
preferential attachment $\gamma=0.8$ $3.2$ $2.6$ $8.5$ $4.7$ $1.4$ $1.6$ $3.5$ $4.7$ $1.4$ $1.6$ $3.5$ $4.7$ $4.4$ $1.1$ $1.1$ $1.6$ $1.$	predictor	astro-ph	cond-mat			hep-th
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	graph distance (all distance-three pairs)					8.6
initing time hitting time hitting time horizontal part of the problem of the pro	preferential attachment	3.2	2.6	8.6	4.7	1.4
initing time—normed by stationary distribution  20 25 0.0 25 6.6 commute time—commed by stationary distribution  38 5.9 21.1 5.9 6.6 commute time—normed by stationary distribution  26 0.8 1.1 4.8 4.7 cotocle PageRank $\alpha = 0.05$ 5.3 13.5 21.1 8.7 16.6 commute time—normed by stationary distribution $\alpha = 0.05$ 6.3 13.5 21.1 8.7 16.6 commute time—according to the property of the	SimRank $\gamma =$	0.8 5.9	14.3	10.6	7.6	21.9
$\begin{array}{c} \text{commute time} \\ co$	hitting time					
$\begin{array}{c} \text{commute time-normed by stationary distribution} \\ \text{rototed PageRank} \\ & \alpha = -0.01 \\ & \alpha = -0.15 \\ & \alpha = -0.05 \\ & \alpha = -0.005 \\ & \beta = -0.0005 \\ & \beta $	hitting time—normed by stationary distribution	n 2.0	2.5	0.0	2.5	6.6
$\begin{array}{c} \cos \cot \text{PageRank} & \alpha = 0.01 \\ \alpha = 0.05 \\ \alpha = 0.05 \\ \alpha = 0.15 \\ \alpha = 0.05 \\ \alpha = 0.05 \\ \alpha = 0.05 \\ \alpha = 0.03 \\ \alpha = 0.05 \\ \beta = 0.000 \\ \beta = 0.05 \\ \beta = 0.000 \\ \beta = 0.05 \\ \beta = 0.0005 \\ \beta = $	commute time	3.8	5.9	21.1	5.9	6.6
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Katz (weighted) $\beta = 0.05$ 1.5 5.9 11.6 2.3 2.7 $\beta = 0.005$ 5.5 14.3 28.5 4.2 12.6 $\beta = 0.0005$ 6.2 13.5 27.5 4.2 12.6 $\beta = 0.0005$ 6.2 13.5 27.5 4.2 12.6 $\beta = 0.005$ 6.2 3.1 12.7 30.6 9.0 12.6 $\beta = 0.005$ 9.1 11.8 30.6 5.1 17.9 $\beta = 0.005$ 9.2 11.8 30.6 5.1 17.9 1.0 12.6 $\beta = 0.005$ 9.2 11.8 30.6 5.1 17.9 1.0 $\beta = 0.005$ 9.2 11.8 30.6 5.1 17.9 1.0 $\beta = 0.005$ 9.2 11.8 30.6 5.1 17.9 1.0 $\beta = 0.005$ 9.2 11.8 30.6 5.1 17.9 1.0 $\beta = 0.005$ 11.8 32.7 2.0 4.0 6.0 $\beta = 0.005$ 11.8 32.7 2.0 4.0 4.0 $\beta = 0.005$ 11.8 32.7 4.2 8.5 19.9 19.0 $\beta = 0.005$ 11.8 32.7 4.2 8.5 19.9 19.0 $\beta = 0.005$ 11.8 32.1 40.0 $\beta = 0.005$ 11.8 42.2 39.0 $\beta = 0.005$ 11.8 42.2 39						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						
Katz (unweighted) $\beta=0.05$ 2.3 12.7 30.6 9.0 12.6 $\beta=0.005$ 9.1 11.8 30.6 5.1 17.9 $\beta=0.0005$ 9.2 11.8 30.6 5.1 17.9 $\beta=0.0005$ 9.2 11.8 30.6 5.1 17.9 $\beta=0.0005$ 9.2 11.8 30.6 5.1 17.9 Low-rank approximation: rank = 1024 2.3 2.5 9.5 3.9.9 10.6 rank = 10.6 5.3 6.7 6.3 6.8 15.3 rank = 16 5.3 6.7 6.3 6.8 15.3 rank = 16 5.3 6.7 6.3 6.8 15.3 rank = 4 5.1 6.7 32.7 2.0 4.7 rank = 1 6.1 2.5 32.7 4.2 8.0 Low-rank approximation: rank = 1024 4.1 6.7 6.3 5.9 13.3 rank = 266 3.8 8.4 3.2 8.5 19.9 rank = 64 2.9 11.8 2.1 4.0 10.0 rank = 16 4.8 8.4 4.2 5.9 16.6 rank = 16 4.4 8.4 4.2 5.9 16.6 rank = 16 4.2 5.3 2.7 4.2 8.0 16.0 rank = 16 4.2 5.3 2.7 4.2 8.0 16.0 rank = 16 4.2 5.3 2.7 4.2 8.0 16.0 rank = 16 4.2 5.3 2.7 4.2 8.0 16.0 rank = 16 4.2 5.3 2.7 4.2 8.0 16.0 rank = 16 4.5 8.0 5.7 5.0 4.7 rank = 16 1.2 5.5 3.7 4.2 8.0 16.0 rank = 16 5.0 8.4 5.3 5.9 13.3 16.0 rank = 16 5.0 8.4 5.3 5.9 15.9 rank = 46 2.8 11.8 2.1 4.2 10.6 rank = 16 5.0 8.4 5.3 5.9 15.9 rank = 4 5.2 6.7 28.5 5.0 4.7 rank = 16 5.0 8.4 5.3 5.9 15.9 rank = 4 5.2 6.7 28.5 2.0 4.7 rank = 10 3.2 5.3 2.7 4.2 8.0 18.0 rank = 10 3.2 5.3 2.7 4.2 8.0 18.0 rank = 10 3.2 5.3 2.7 4.2 8.0 18.0 rank = 10 3.2 5.3 2.7 4.2 8.0 18.0 rank = 10 3.2 5.3 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.5 2.7 4.2 8.0 18.0 rank = 10 3.2 5.2 5.7 1.1 3.7 4.5 2.2 3.0 18.0 rank = 10 3.2 5.0 5.9 7.0 7.0 7.0 7.0 7.0 7.0 7.0 7.0 7.0 7.0						
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β = 0.0008   9.2   11.8   30.6   5.1   17.9   Low-rank approximation: rank = 1024   2.3   2.5   9.5   4.0   6.0   Inner product rank = 256   4.8   5.9   5.3   9.9   10.5   rank = 64   3.8   12.7   5.3   7.1   11.3   rank = 16   5.3   6.7   6.3   6.8   15.3   rank = 16   5.3   6.7   32.7   2.0   4.7   rank = 1   6.1   2.5   32.7   2.0   4.7   rank = 1   6.1   2.5   32.7   2.0   4.7   rank = 1   6.1   2.5   32.7   2.0   4.7   rank = 16   4.4   4.6   6.7   32.7   2.0   4.7   rank = 16   4.4   8.4   4.2   5.9   13.3   Matrix entry rank = 266   3.8   8.4   3.2   8.5   rank = 64   2.9   11.8   2.1   4.0   10.0   rank = 16   4.4   8.4   4.2   5.9   16.6   rank = 4   4.9   6.7   27.5   2.0   4.7   rank = 1   6.1   2.5   32.7   2.0   4.7   rank = 1   6.1   2.5   32.7   2.0   4.7   rank = 16   4.8   8.4   3.2   8.5   Matrix entry rank = 166   4.8   8.4   4.2   5.9   16.6   rank = 4   4.9   6.7   27.5   2.0   4.7   rank = 16   5.8   8.4   3.2   8.5   Matrix entry rank = 166   4.8   8.1   2.1   4.2   10.6   rank = 4   4.9   6.7   27.5   2.0   4.7   rank = 16   5.8   8.4   3.2   8.5   rank = 46   2.8   11.8   2.1   4.2   10.6   rank = 46   2.8   11.8   2.1   4.2   10.6   rank = 4   5.2   6.7   28.5   5.9   15.9   rank = 4   5.2   6.7   28.5   2.0   4.7   rank = 1   0.3   2.5   32.7   4.2   8.0   rank = 4   5.2   6.7   28.5   2.0   4.7   rank = 1   0.3   2.5   32.7   4.2   8.0   rank = 4   5.2   6.7   28.5   2.0   4.7   rank = 4   5.2   6.7   28.5   5.9   15.9   rank = 4   5.2   6.7   28.5   5.0   4.7   rank = 4   5.2   6.7   28.5   5.0   4.7   rank = 4   5.2   6.7   28.5   5.0   rank = 4   5.2   6.7   28.5   5.0   4.7   rank = 4   5.2   6.7   28.5   5.0   rank = 4   5.2   6.7   28.5   rank = 64   6.8   6.7   11.8   rank = 64   6.8   6.7						
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$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$\beta = 0.00$			30.6	5.1	17.9
rank = 64 a. 3.8 a. 12.7 b. 3.3 c. 7.1 a. 11.3 a. a. 12.7 c. 3.3 c. 7.1 a. 11.3 a.	Low-rank approximation: rank = 1		2.5	9.5		6.0
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	Inner product rank = :					10.6
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						
Low-rank approximation:   rank = 1024   4.1   6.7   6.3   5.9   13.3     Matrix entry   rank = 64   2.9   11.8   2.1   4.0   10.0     rank = 16   4.4   8.4   4.2   5.9   16.6     rank = 16   6.1   2.5   32.7   4.2   8.0     Low-rank approximation:   rank = 1024   4.3   6.7   28.5   5.9   13.3     rank = 20   6.8   4.3   8.2   8.5   5.9   13.3     rank = 16   5.0   8.4   5.3   5.9   15.9     rank = 64   2.8   11.8   2.1   4.2   10.6     rank = 16   5.0   8.4   5.3   5.9   15.9     rank = 4   5.2   6.7   28.5   2.0   4.7     rank = 4   0.3   2.5   32.7   4.2   8.0     rank = 4   0.2   0.7   28.5   2.0   4.7     rank = 4   0.3   2.5   32.7   4.2   8.0     rank = 5   5.8   6.7   14.8   4.2   23.9     weighted)   common neighbors, δ = 8   5.8   6.7   14.8   4.2   23.9     rank = 1   0.005), δ = 8   5.2   10.1   22.2   2.8     rank = 1   0.005, δ = 8   5.4   5.1   13.7   4.5     rank = 1   7.6   6.6   6.1   2.0   3.7   15.3     rank = 1   7.6   6.6   6.3   8.4   25.3   4.8   21.9     rank = 1   7.6   6.5   6.5   5.1   7.9   2.2   2.0     rank = 1   7.6   6.6   6.7   2.2   2.0     rank = 1   7.6   6.8   6.7   2.2   2.0     rank = 1   7.6   6.7   2.2   2.0     rank = 1   7.6   6.8   2.2   2.7   7.6   6.6     rank = 1   7.4   2.2   2.3   5.9   7.4   4.5   8.0     rank = 1   7.4   2.3   5.9   7.4   4.5   8.0     rank = 1   7.4   2.3   5.9   7.4   4.5   8.0     rank = 1   7.4   2.3   5.9   7.4   4.5   8.0     rank = 1   7.6   7.8   7.4   5.8   8.0     rank = 1   7.8   7.8   7.4   5.8   8.0     rank = 1   7.8   7.8   7.4   8.6   7.8   7.4   8.6     rank = 1   7.8   7.8   7.8   7.4   8.6						
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rank = 64						
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$\begin{array}{c} \text{(unweighted)} & \text{common neighbors, } \delta = 16 \\ \hline \text{Katz } (\beta = 0.005),  \delta = 8 \\ \hline \text{Katz } (\beta = 0.005),  \delta = 16 \\ \hline \text{(Satz } (\beta = 0.005),  \delta = 16 \\ \hline \text{(Satz } (\beta = 0.005),  \delta = 16 \\ \hline \text{(Satz } (\beta = 0.005),  \delta = 16 \\ \hline \text{(Satz } (\beta = 0.001),  \delta = 10 \\ \hline \text{(Satz } (\beta = 0.001,  \beta_2 = 0.1) \\ \hline \rho = 0.15 \\ \hline \rho = 0.20 \\ \hline \end{pmatrix} \begin{array}{c} 3.2 \\ \hline \text{(Satz } (\beta = 0.001,  \beta_2 = 0.1) \\ \hline \text{(Satz } (\beta = 0.0$						
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clustering: $\rho = 0.10$ 3.2 4.2 31.7 7.1 8.6 Katz $(\beta_1 = 0.001, \beta_2 = 0.1)$ $\rho = 0.15$ 4.6 4.2 32.7 7.6 6.6 $\rho = 0.20$ 2.3 5.9 7.4 4.5 8.0						
Katz $(\beta_1 = 0.001, \beta_2 = 0.1)$ $\rho = 0.15$ 4.6 4.2 32.7 7.6 6.6 $\rho = 0.20$ 2.3 5.9 7.4 4.5 8.0	0 71		-			
$\rho = 0.20$ 2.3 5.9 7.4 4.5 8.0						
$\rho = 0.25$ 2.0 11.8 6.3 6.8 5.3						
	$\rho = 0$	0.25 2.0	11.8	6.3	6.8	5.3

#### Evaluation: the breadth of data

#### Three additional datasets

- 1. Proceedings of STOC and FOCS
- 2. Papers for Citeseer
- 3. All five of the arXiv sections

STOC/FOCS	arXiv sections	combined arXiv sections	Citeseer
6.1	18.0-46.9	71.2	147.0

Common neighbors vs Random

✓ Suggests that is easier to predict links within communities

#### **Extensions**

- ❖ Improve performance. Even the best (Katz clustering on gr-qc) correct on only about 16% of its prediction
- Improve efficiency on very large networks (approximation of distances)
- Treat more recent collaborations as more important
- Additional information (paper titles, author institutions, etc)

To some extent latently present in the graph

#### **Outline**

- Estimating a score for each edge (seminal work of Liben-Nowell&Kleinberg
  - Neighbors measures, Distance measures, Other methods
  - Evaluation
- Classification approach
- Twitter

## **Using Supervised Learning**

Given a collection of records (*training set* )

Each record contains a set of attributes (features) + the class attribute.

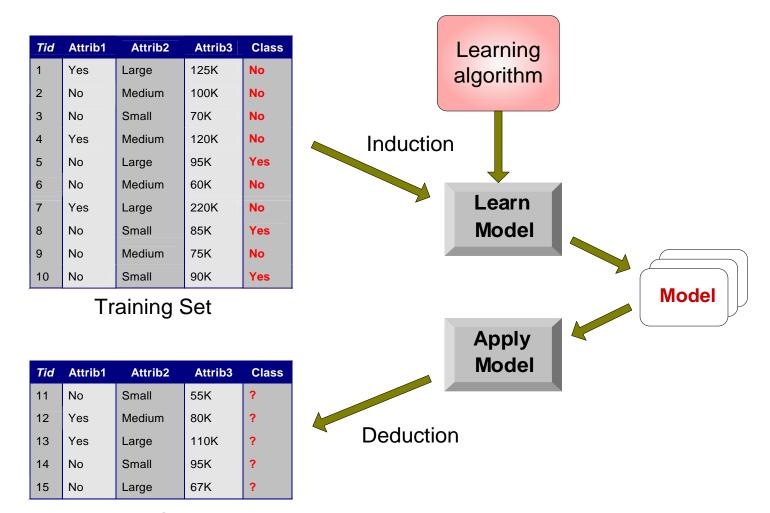
Find a *model* for the class attribute as a function of the values of other attributes.

Goal: previously unseen records should be assigned a class as accurately as possible.

A test set is used to determine the accuracy of the model.

Usually, the given data set is divided into training and test sets, with training set used to build the model and test set used to validate it.

# Illustrating the Classification Task



Test Set

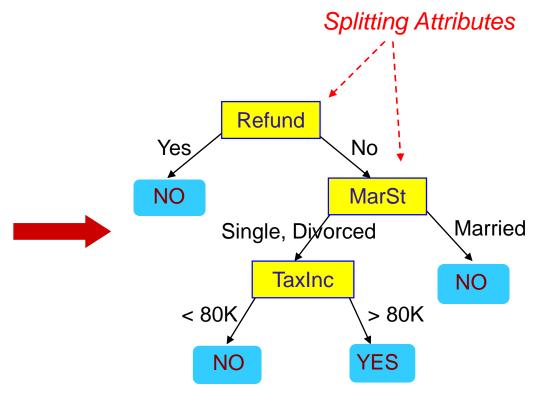
## Classification Techniques

- Decision Tree based Methods
- Rule-based Methods
- Memory based reasoning
- Neural Networks
- Naïve Bayes and Bayesian Belief Networks
- Support Vector Machines
- Logistic Regression

## Example of a Decision Tree

categorical continuous

Tid	Refund	Marital Status	Taxable Income	Cheat
1	Yes	Single	125K	No
2	No	Married	100K	No
3	No	Single	70K	No
4	Yes	Married	120K	No
5	No	Divorced	95K	Yes
6	No	Married	60K	No
7	Yes	Divorced	220K	No
8	No	Single	85K	Yes
9	No	Married	75K	No
10	No	Single	90K	Yes



**Training Data** 

Model: Decision Tree

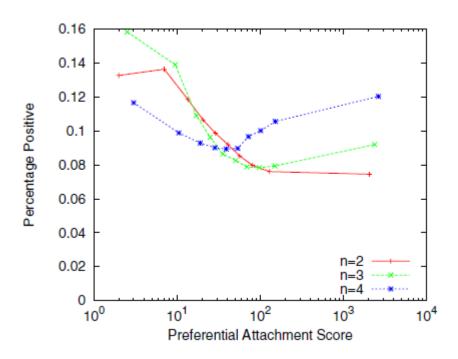
#### Classification for Link Prediction

Class? Features (predictors)?

Name	Parameters	HPLP	HPLP+
In-Degree(i)	-	<b>√</b>	✓
In-Volume(i)	-	✓	✓
In-Degree(j)	-	✓	✓
In-Volume(j)	-	✓	✓
Out-Degree(i)	-	✓	✓
Out-Volume(i)	-	✓	✓
Out-Degree(j)	-	✓	✓
Out-Volume(j)	-	✓	✓
Common $Nbrs(i,j)$	-	✓	✓
Max. Flow $(i,j)$	l = 5	✓	✓
Shortest $Paths(i,j)$	l = 5	✓	✓
PropFlow(i,j)	l = 5	✓	✓
Adamic/Adar(i,j)	-		✓
Jaccard's $Coef(i,j)$	-		✓
Katz(i,j)	$l = 5,  \beta = 0.005$		✓
Pref Attach $(i,j)$	-		✓

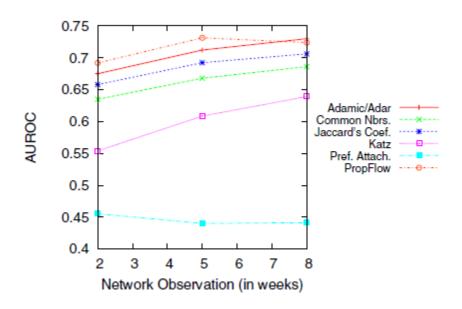
PropFlow: random walks, stops at I or when cycle

## Using Supervised Learning: why?



- Even training on a single feature may outperform ranking (restriction to n-neighborhoods)
- Dependencies between features

## How to split the graph to get train data



- tx length of computing features ty length of determining the class attribute
- Large tx => better quality of features as the network reaches saturation
- Increasing ty increases positives

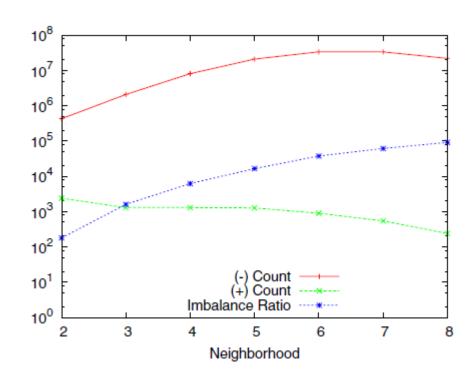
#### **Imbalance**

Sparse networks: |E| = k |V| for constant k << |V|</p>

The class imbalance ratio for link prediction in a sparse network is  $\Omega(|V|/1)$  when at most |V| nodes are added

Missing links is  $|V|^2$ Positives V

Treat each neighborhood as a separate problem



### Metrics for Performance Evaluation

#### **Confusion Matrix:**

	PREDICTED CLASS						
		Class=Yes	Class=No				
A O.T. I.A.I	Class=Yes	TP	FN				
ACTUAL CLASS	Class=No	FP	TN				

$$Accuracy = \frac{TP + TN}{TP + TN + FP + FN}$$

#### **ROC Curve**

TPR (sensitivity)=TP/(TP+FN) (percentage of positive classified as such)

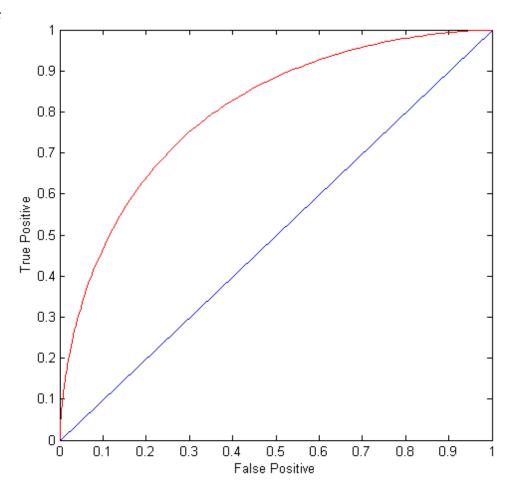
FPR = FP/(TN+FP) (percentage of negative
 classified as positive)

- (0,0): declare everything to be negative class
- (1,1): declare everything to be positive class
- (0,1): ideal

Diagonal line: Random guessing

Below diagonal line: prediction is

opposite of the true class



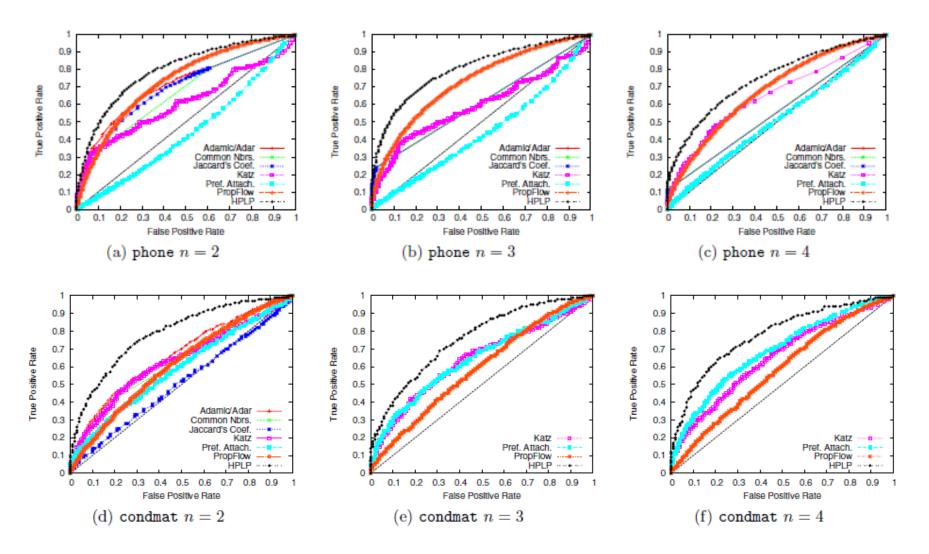
**AUC**: area under the ROC

### Results

Ensemble of classifiers: Random Forest

Random forest: Ensemble classifier that constructs a multitude of decision trees at training time and output the class that is the mode (most frequent) of the classes (classification) or mean prediction (regression) of the individual trees.

## Results



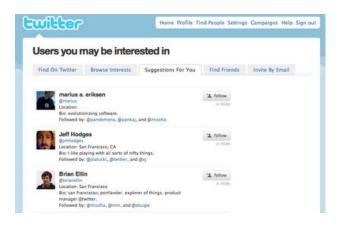
#### **Outline**

- Estimating a score for each edge (seminal work of Liben-Nowell&Kleinberg
  - Neighbors measures, distance measures, other methods
  - Evaluation
- Classification approach
  - Brief background on classification
  - Issues
- The who to follow service at Twitter
  - Some practical considerations
  - Overall architecture of a real social network
  - SALSA (yet another link analysis algorithm
  - Some evaluation issues

### Introduction

Wtf ("Who to Follow"): the Twitter user recommendation service

- Twitter: 200 million users, 400 million tweets every day (as of early 2013)
   <a href="http://www.internetlivestats.com/twitter-statistics/">http://www.internetlivestats.com/twitter-statistics/</a>
- Twitter needs to help existing and new users to discover connections to sustain and grow
- Also used for search relevance, discovery, promoted products, etc.



## Introduction

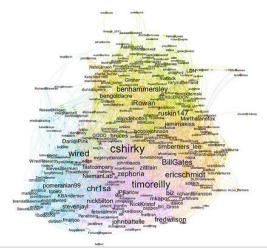
#### Difference between:

- Interested in
- Similar to

Is it a "social" network as Facebook?

# The Twitter graph

- Node: user (directed) edge: follows
- Statistics (August 2012)
  - over 20 billion edges (only active users)
  - power law distributions of in-degrees and out-degrees.
  - over 1000 with more than 1 million followers,
  - 25 users with more than 10 million followers.



http://blog.ouseful.info/2011/07/07/visualising-twitter-friend-connections-using-gephi-an-example-using-wireduk-friends-network/

## The Twitter graph: storage

- Stored in a graph database called FlockDB which uses MySQL as the underlying storage engine
- Sharding and replication by a framework called Gizzard
- Both custom solutions developed internally but open sourced
- FlockDB holds the "ground truth" for the state of the graph
- Optimized for low-latency, high-throughput reads and writes, and efficient intersection of adjacency lists (needed to deliver @-replies, or messages targeted to a specific user received mutual followers of the sender and recipient)
  - hundreds of thousands of reads per second and tens of thousands of writes per second.

# The Twitter graph: analysis

- Instead of simple get/put queries, many graph algorithms involve large sequential scans over many vertices followed by self-joins (for example, to materialize egocentric follower neighborhoods)
- not time sensitive unlike graph manipulations tied directly to user actions (adding a follower) which have tight latency bounds.

**OLAP** (online analytical processing) vs. **OLTP** (online transaction processing)

analytical workloads that depend on sequential scans vs. short, primarily seek-based workloads that provide an interactive service

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**OLAP** (online analytical processing) vs. **OLTP** (online transaction processing)

analytical workloads that depend on sequential scans vs. short, primarily seek-based workloads that provide a user-facing service

## History of WTF

3 engineers, project started in spring 2010, product delivered in summer 2010

Basic assumption: the whole graph fits into memory of a single server

## Design Decisions: To Hadoop or not?

#### Case study: MapReduce implementation of PageRank

- Each iteration a MapReduce job
- Serialize the graph as adjacency lists for each vertex, along with the current PageRank value.
- Mappers process all vertices in parallel: for each vertex on the adjacency list, the mapper emits an intermediate key-value pair: (destination vertex, partial PageRank)
- Gather all key-value pairs with the same destination vertex, and each Reducer sums up the partial PageRank contributions
- Convergence requires dozens of iterations. A control program sets up the MapReduce job, waits for it to complete, and checks for convergence by reading in the updated PageRank vector and comparing it with the previous.
- This basic structure can be *applied to a large class* of "message-passing" graph (e.g., breadth-first search)

## Design Decisions: To Hadoop or not?

#### Shortcomings of MapReduce implementation of PageRank

- MapReduce jobs have relatively high startup costs (in Hadoop, on a large, busy cluster, can be tens of seconds), this places a lower bound on iteration time.
- Scale-free graphs, whose edge distributions follow power laws, create stragglers in the reduce phase. (e.g., the reducer assigned to google.com)
   Combiners and other local aggregation techniques help
- Must shuffle the graph structure (i.e., adjacency lists) from the mappers to the reducers at each iteration. Since in most cases the graph structure is static, wasted effort (sorting, network traffic, etc.).
- The *PageRank vector is serialized to HDFS*, along with the graph structure, at each iteration. Excellent fault tolerance, but at the cost of performance.

# Design Decisions: To Hadoop or not?

#### Besides Hadoop:

Improvements: HaLoop Twister, and PrIter

#### Alternatives:

- Google's Pregel implements the Bulk Synchronous Parallel model: computations at graph vertices that dispatch "messages" to other vertices. Processing proceeds in supersteps with synchronization barriers between each.
- GraphLab and its distributed variant: computations either through an update function which defines local computations (on a single vertex) or through a sync mechanism which defines global aggregation in the context of different consistency models.

# Design Decisions: To Hadoop or not?

#### Decided to build their own system

#### Hadoop reconsidered:

- new architecture completely on Hadoop
- in Pig a high-level dataflow language for large, semi-structured datasets compiled into physical plans executed on Hadoop
- Pig Latin primitives for projection, selection, group, join, etc.

#### Why not some other graph processing system?

For compatibility, e.g., to use existing analytics hooks for job scheduling, dependency management, etc.

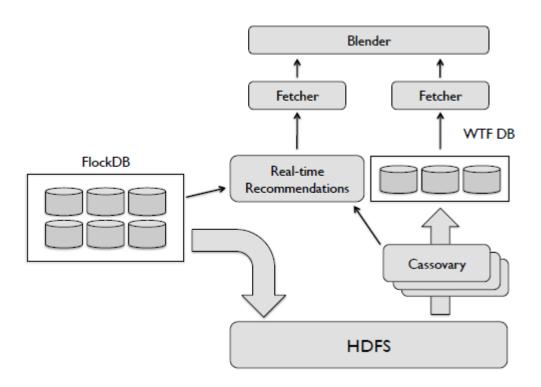
# Design Decisions: How much memory?

#### in-memory processing on a single server

#### Why?

- 1. The alternative (a partitioned, distributed graph processing engine) significantly more complex and difficult to build,
- It was feasible (72GB -> 144GB, 5 bytes per edge (no metadata);
   24-36 months lead time)
- In memory not uncommon (google indexes + Facebook, Twitter many cache servers
- A single machine
  - Graph distribution still hard (hash partitioning, minimize the number of edges that cross-partition (two stage, over partition #clusters>>#servers, still skew problems), use replication to provide n-hop guarantee (all nneighbors in a singe site)
  - Avoids extra protocols (e.g., replication for fault-tolerance)

# **Overall Architecture**



# Overall Architecture: Cassovary

#### In memory graph processing engine, written in Scala

- Once loaded into memory, graph is immutable
- Fault tolerance provided by replication, i.e., running many instances of Cassovary, each holding a complete copy of the graph in memory.
- Access to the graph via vertex-based queries such as retrieving the set of outgoing edges for a vertex and sampling a random outgoing edge.
- Multi-threaded: each query is handled by a separate thread.
- Graph stored as optimized adjacency lists: the adjacency lists of all vertices in large shared arrays plus indexes (start, length) into these shared arrays
- No compression
- Random walks implemented using the Monte-Carlo method,
  - the walk is carried out from a vertex by repeatedly choosing a random outgoing edge and updating a visit counter.
  - Slower than a standard matrix-based implementation, but low runtime memory overhead

# Overall Architecture: Flow

- 1. Daily snapshots of the Twitter graph imported from FlockDB into the Hadoop data warehouse
- 2. The entire graph loaded into memory onto the Cassovary servers, each holds a complete copy of the graph in memory.
- Constantly generate recommendations for users consuming from a distributed queue containing all Twitter users sorted by a "last refresh" timestamp (~500 ms per thread to generate ~100 recommendations for a user)
- 4. Output from the Cassovary servers inserted into a sharded MySQL database, called, WTF DB.
- 5. Once recommendations have been computed for a user, the user is enqueued again with an updated timestamp. Active users who consume (or are estimated to soon exhaust) all their recommendations are requeued with much higher priority; typically, these users receive new suggestions within an hour.

# Overall Architecture: Flow

Graph loaded once a day, what about new users?

#### Link prediction for new users

- Challenging due to sparsity: their egocentric networks small and not well connected to the rest of the graph (cold start problem)
- Important for social media services: user retention strongly affected by ability to find a community with which to engage.
- Any system intervention is only effective within a relatively short time window. (if users are unable to find value in a service, they are unlikely to return)
- 1. new users are given high priority in the Cassovary queue,
- 2. a completely independent set of algorithms for real-time recommendations, specifically targeting new users.

# Algorithms: Circle of trust

Circle of trust: the result of an egocentric random walk (similar to personalized (rooted) PageRank)

- Computed in an online fashion (from scratch each time) given a set of parameters (# of random walk steps, reset probability, pruning settings to discard low probability vertices, parameters to control sampling of outgoing edges at vertices with large out-degrees, etc.)
- Used in a variety of Twitter products, e.g., in search and discovery, content from users in one's circle of trust upweighted

# Algorithms

- Asymmetric nature of the follow relationship (other social networks e.g., Facebook or LinkedIn require the consent of both participating members)
- Directed edge case is similar to the user-item recommendations problem where the "item" is also a user.

SALSA (Stochastic Approach for Link-Structure Analysis) a variation of HITS

As in HITS

hubs authorities

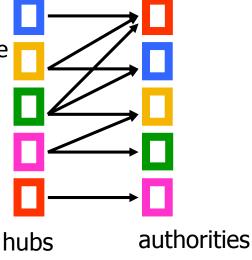
#### HITS

- Good hubs point to good authorities
- Good authorities are pointed by good hubs

hub weight = sum of the authority weights of the
authorities pointed to by the hub

$$h_i = \sum_{j:i\to j} a_j$$

authority weight = sum of the hub weights that point to this authority.  $a_i = \sum_i h_i$ 



#### Random walks to rank hubs and authorities

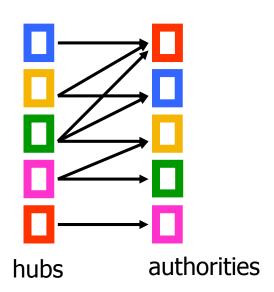
- Two different random walks (Markov chains): a chain of hubs and a chain of authorities
- Each walk traverses nodes only in one side by traversing two links in each step h->a-h, a->h->a

Transition matrices of each Markov chain:

H and A

W: the adjacency of the directed graph  $W_r$ : divide each entry by the sum of its row  $W_c$ : divide each entry by the sum of its column

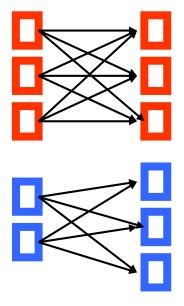
$$H = W_r W_c^T$$
$$A = W_c^T W_r$$



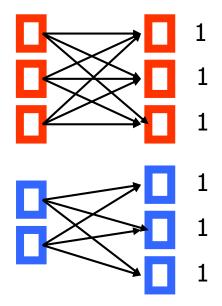
Proportional to the degree

- Reduces the problem of HITS with tightly knit communities (TKC effect)
- Better for single-topic communities
- More efficient implementation

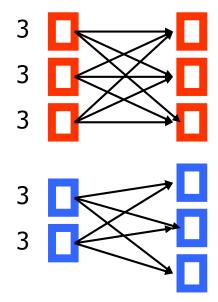
- The HITS algorithm favors the most dense community of hubs and authorities
  - Tightly Knit Community (TKC) effect



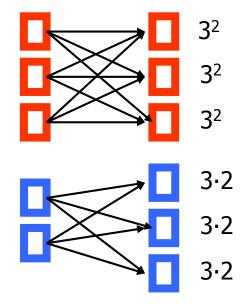
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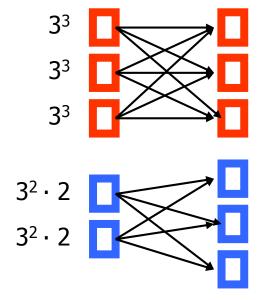
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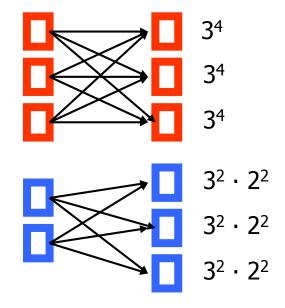
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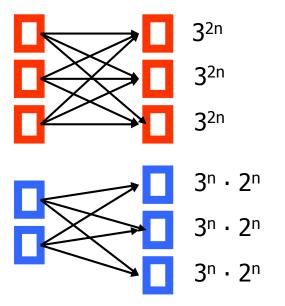


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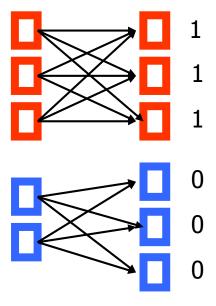
- The HITS algorithm favors the most dense community of hubs and authorities
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weight of node p is proportional to the number of (BF)<sup>n</sup> paths that leave node p

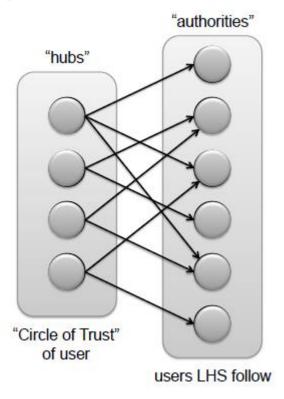


after n iterations

- The HITS algorithm favors the most dense community of hubs and authorities
  - Tightly Knit Community (TKC) effect



after normalization with the max element as  $n \rightarrow \infty$ 



Hubs: 500 top-ranked nodes from the user's circle of trust

Authorities: users that the hubs follow

Hub vertices: user similarity (based on homophily, also useful)

Authority vertices: "interested in" user recommendations.

#### How it works

SALSA mimics the recursive nature of the problem:

- A user u is likely to follow those who are followed by users that are similar to u.
- A user is similar to u if the user follow the same (or similar) users.
- I. SALSA provides *similar users* to u on the LHS and *similar followings* of those on the RHS.
- II. The random walk ensures equitable distribution of scores in both directions
- III. Similar users are selected from the circle of trust of the user through personalized PageRank.

# **Evaluation**

- Offline experiments on retrospective data
- Online A/B testing on live traffic

#### Various parameters may interfere:

- How the results are rendered (e.g., explanations)
- Platform (mobile, etc.)
- New vs old users

# **Evaluation:** metrics

#### Follow-through rate (FTR) (precision)

- Does not capture recall
- Does not capture lifecycle effects (newer users more receptive, etc.)
- Does not measure the quality of the recommendations:
   all follow edges are not equal

#### Engagement per impression (EPI):

After a recommendation is accepted, the amount of engagement by the user on that recommendation in a specified time interval called the observation interval.

## **Extensions**

- Add metadata to vertices (e.g., user profile information) and edges (e.g., edge weights, timestamp, etc.)
- Consider interaction graphs (e.g., graphs defined in terms of retweets, favorites, replies, etc.)

## **Extensions**

#### Two phase algorithm

- Candidate generation: produce a list of promising recommendations for each user, using any algorithm
- Rescoring: apply a machine-learned model to the candidates, binary classification problem (logistic regression)

First phase: recall + diversity

Second phase: precision + maintain diversity

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