

Poster at EuroSys '08, Glasgow, UK **Okeanos - Reliable Archival Storage for Heterogeneous** Stream Data

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Low-Rate Stream Results Introduction Journaling Inefficiencies • A general-purpose stream archival facility • Operation sequence in ordered journaling: • Final location traffic: could serve as a building block for a variety of At each journal commit interval, before the Ordered journaling iournal record is updated (metadata only), data applications, e.g. random access is flushed to the final location network packet monitoring > more data than needed At each pdflush wake-up interval, both data and urban traffic control . metadata are flushed to their final locations Data journaling General monitoring case: >only needed data logged within journal record: messages received from massive numbers of · metadata written synchronously /sec Final Location :: Bit Rate = 1Kbps sensors (Bytes x 10.⁶) /₁ sequential writes → efficiency reception at potentially different rates -*-Ext3 small amount of data → efficiency data should be stably stored on disk within final location: • Existing systems' inadequacies: metadata written asynchronously Traditional systems (such as relational DBs): · data written synchronously 0.5 0.4 0.3 0.2 0.1 0.5 not engineered to efficiently store continuous continuous disk traffic → inefficient for small writes stream data automatically generated from sensors in real time Operation sequence in data journaling: Modern stream storage servers: At each journal commit interval, the journal · basically designed to store stream files of limited Disk 1 record updated (both data and metadata) size for repetitive playback 0 50 100 150 200 250 300 400 500 Number of Streams 650 · inefficient to constantly accumulate continuous At each pdflush wake-up interval, both data and stream data for archival purposes metadata are flushed to their final locations Journal traffic: • Monitoring sensors may generate: within journal record Ordered journaling high-resolution video and audio streams at · both data and metadata written synchronously > only metadata logged large rates sequential writes → efficiency intermittent variations in environmental large amount of data → inefficient for large Data journaling conditions at much lower rates volumes of data ≻sequential writes ≻more data than needed within final location Summarizing received data should: · both data and metadata written asynchronously External Journal :: Bit Rate = 1Kbps 9 2.5 be stably stored on the storage facility deferred writes (write-coalescing) → efficient for 6 (Bytes x 10.)/ not compromise the sequential playback - Ext2 - Data journaling small writes performance 2 Journaling Filesystems Ordered Journaling 1.5 Preserve filesystem consistency across system Metadata **Disk Throughput** crashes at minimal recovery time 1 improvement of operation reliability Data Journaling 0.5 Metadata * + Data Two alternative journaling modes can be used: Journal Record metadata-only logging (ordered journaling) 0 50 100 150 200 250 300 400 500 Number of Streams 650 write-ahead logging of file contents (data

In Progress

- For each individual stream the system should automatically:
 - identify the most appropriate journaling approach
 - adjust its behavior according to the varying features of the stream over time
- Data journaling inefficient for low-rate streams Possible solution: differential logging

References

- [1] P. J. Desnoyers and P. Shenoy, Hyperion: High Volume Stream Archival for Retrospective Querying, USENIX Annual Technical Conference, June 2007
- [2] V. Prabhakaran et al., Analysis and Evolution of Journaling File Systems. USENIX Annual Technical Conference, April 2005, pp. 105-120.

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- journaling)

Data journaling negatively affects disk

throughput in sequential write workloads doubly stores data at both journal record and the final location in the file system structures

According to aggregate workload characteristics:

- Ordered journaling \rightarrow efficient for sequential access
- Data journaling \rightarrow efficient for random access

Stream Rate Characteristics

- Efficient and reliable storage of multiple concurrent streams
 - aggregate workload \rightarrow random-access behavior appends corresponding to individual streams \rightarrow perfectly sequential

Data received from low-rate streams:

- significant overhead for the immediate movement from memory to the final disk location
- disk penalized with small writes
- flushing data to the final location can be deferred to a later more convenient time
- Data received from higher-rate streams: larger amount of data
 - can be moved directly to their final destination on disk without compromising the efficiency of the storage device



High-Rate Stream Results

O Total disk traffic:

- random-access behavior
- ordered journaling better than data journaling





Okeanos - Reliable archival storage for heterogeneous stream data¹

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The prevalence of continuous monitoring processes for system management purposes and general physical site safety make stream processing applications highly relevant in modern computing infrastructures. Recently proposed stream management engines demonstrate the feasibility of flexibly applying time-series operators on massive numbers of streams in real time as their data arrive to the system. In principle, dropping prices in computer hardware should also make possible the storage of high-resolution or numerous streams for entire months or years.

Prior research has made the case that traditional systems (such as relational databases or general-purpose file systems) are not engineered to efficiently store continuous stream data that are automatically generated from sensors in real time [1]. Similarly, modern stream storage servers are basically designed to store stream files of limited size for repetitive playback rather than constantly accumulating continuous stream data for archival purposes. In the general monitoring case, we are interested in receiving messages from massive numbers of sensors at potentially different rates and reliably storing their data on disk files before acknowledging their reception as successful. Some sensors may generate high-resolution video and audio streams at large rates while others may send intermittent variations in environmental conditions at much lower rates. Across all these heterogeneous cases, we need the received data to be stably stored on the same storage facility without compromising the sequential playback performance required for effective visualization or statistics-gathering processing.

In our vision, a general-purpose stream archival facility could serve as a building block for a variety of applications in the entire range from network packet monitoring to urban traffic control with the appropriate indexing functionality built separately at a higher level when needed. In order to improve their operation reliability, general-purpose file systems apply journaling techniques to preserve metadata consistency across system crashes at minimal recovery time. Some of them additionally use writeahead logging of file contents (or data journaling) in order to prevent small writes from penalizing disk access performance. Nevertheless, data journaling negatively affects disk throughput in sequential write workloads due to doubly storing data at both their temporary journal record and also their final location in the file system structures (Figure 1). Alternatively, disk writing of data to its final location before updating the corresponding metadata can provide consistency guarantees similar to those of data journaling without the extra overhead associated with the latter. Comparisons across different journaling methods with general-purpose file server traffic have shown that either

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ordered data writing or data journaling may lead to better performance depending on whether the aggregate workload is sequential or random-access [2]. In our undergoing research, we focus on the efficient and reliable storage of multiple concurrent streams whose aggregate workload demonstrates random-access behavior even though appends corresponding to individual streams may be perfectly sequential.



Figure 1. Total disk traffic when writing multiple high-rate streams of 10Mbps each on Ext3/Linux 2.6.18. The disk throughput of data journaling is twice that of ordered journaling and flattens out faster.



Figure 2. Journal and final location traffic generated from lowrate streams of 1Kbps. Although, in both journaling modes the disk is penalized with small writes, in contrast to the randomaccess final location writes, log records are written sequentially, which makes differential logging a possible solution to data journaling inefficiency.

Data journaling flushes data to stable log storage and easily restores it after a crash presuming that copying to the final location can be deferred to a later more convenient time. This is useful for low-rate streams that individually would incur significant overhead for their immediate movement from memory to their final disk location (Figure 2). Instead, the larger amounts of data received from higherrate streams can be moved directly to their final destination on disk without compromising the efficient operation of the storage device. Ideally, the system should automatically identify the most appropriate journaling approach for each individual stream and adjust its behavior according to the varying features of the stream over time.

In conclusion, we motivate the necessity for building systems facilities for the archival storage of heterogeneous streams with different rate and content characteristics. Ideally, such facilities should be able to reliable store massive numbers of streams while requiring minimal recovery time across crashes and supporting lowoverhead stream playback for the needs of visualization and statistical processing.

[1] P. J. Desnoyers and P. Shenoy, *Hyperion: High Volume Stream Archival for Retrospective Querying*, USENIX Annual Technical Conference, June 2007.

[2] V. Prabhakaran et al., *Analysis and Evolution of Journaling File Systems*. USENIX Annual Technical Conference, April 2005, pp. 105-120.

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